

Lecture 8 - The theory of pure quantum processes

Summary.

- (1) From linear maps to quantum maps.
- (2) The theory of pure quantum maps.

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Introduction.

The process theory of linear maps is enough to represent typical primitives in quantum computation as processes. For example, classical logic gates are made into linear maps as illustrated below¹. Logic gates are functions, but given a basis, one can always convert the correspondence $\mathbf{x} \mapsto \mathbf{y}$ to linear mapping between states labeled by \mathbf{x} and \mathbf{y} . Each map contributes to a term (expressed as an outer product): summing over all of them recovers the linear map.

$$\boxed{f} = \sum_{(a_1 \dots a_m \mapsto b_1 \dots b_n) \in F} \begin{array}{c} \downarrow b_1 \quad \dots \quad \downarrow b_n \\ \uparrow a_1 \quad \dots \quad \uparrow a_m \end{array}$$

For example,

$$\boxed{\text{XOR}} = \begin{array}{c} \downarrow 0 \\ \uparrow 0 \quad \uparrow 0 \end{array} + \begin{array}{c} \downarrow 1 \\ \uparrow 0 \quad \uparrow 1 \end{array} + \begin{array}{c} \downarrow 1 \\ \uparrow 1 \quad \uparrow 0 \end{array} + \begin{array}{c} \downarrow 0 \\ \uparrow 1 \quad \uparrow 1 \end{array}$$

Encoding such gates as linear maps is key to the circuit model of quantum computation, namely when restricting to gates that yield unitary linear maps, such as NOT. Typical quantum gates, such as CNOT, are expressed similarly

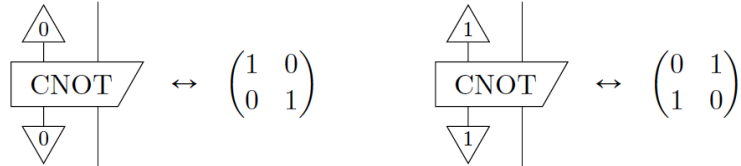
$$\boxed{\text{CNOT}} := \begin{array}{c} \downarrow 0 \quad \downarrow 0 \\ \uparrow 0 \quad \uparrow 0 \end{array} + \begin{array}{c} \downarrow 0 \quad \downarrow 1 \\ \uparrow 0 \quad \uparrow 1 \end{array} + \begin{array}{c} \downarrow 1 \quad \downarrow 1 \\ \uparrow 1 \quad \uparrow 0 \end{array} + \begin{array}{c} \downarrow 1 \quad \downarrow 0 \\ \uparrow 1 \quad \uparrow 1 \end{array}$$

¹Pictures are taken from Coecke and Kissinger book, *Picturing Quantum processes*, CUP, 2017.

whose matrix representation

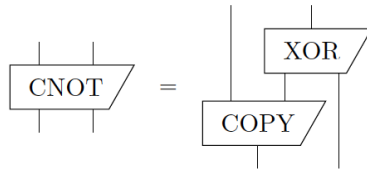
$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

is properly explained by the fact that the first basis element selects which transformation should be applied to the second, as depicted below.

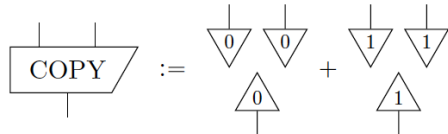


Exercise 1

Show that



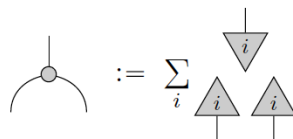
where



The XOR gate is, up to a scalar, the adjoint of COPY over the Hadamard basis

$$\downarrow_0 := \frac{1}{\sqrt{2}} \left(\downarrow_0 + \downarrow_1 \right) \quad \downarrow_1 := \frac{1}{\sqrt{2}} \left(\downarrow_0 - \downarrow_1 \right)$$

which in the diagrammatic language ZX-calculus² is depicted as



²The name comes from the Z and X bases, i.e. the computational and the Hadamard bases.

Actually,

$$\begin{array}{c} \triangle k \\ | \\ \circ \\ / \quad \backslash \\ \triangle i \quad \triangle j \end{array} = \left(\begin{array}{c} \triangle k \\ | \\ \nabla 0 \\ | \\ \triangle 0 \quad \triangle 0 \\ / \quad \backslash \\ \triangle i \quad \triangle j \end{array} + \begin{array}{c} \triangle k \\ | \\ \nabla 1 \\ | \\ \triangle 1 \quad \triangle 1 \\ / \quad \backslash \\ \triangle i \quad \triangle j \end{array} \right) = \frac{1}{\sqrt{2}} \cdot \frac{1}{2} (1 + (-1)^{i+j+k})$$

because

$$\begin{array}{c} \triangle 0 \\ | \\ \nabla i \end{array} = \begin{array}{c} \triangle i \\ | \\ \nabla 0 \end{array} = \frac{1}{\sqrt{2}} \quad \begin{array}{c} \triangle 1 \\ | \\ \nabla i \end{array} = \begin{array}{c} \triangle i \\ | \\ \nabla 1 \end{array} = (-1)^i \frac{1}{\sqrt{2}}$$

Expression

$$\frac{1}{\sqrt{2}} \cdot \frac{1}{2} (1 + (-1)^{i+j+k})$$

is equal to $\frac{1}{\sqrt{2}}$ if $i + j + k$ is even, and 0 otherwise. But $i + j + k$ is even when $i \text{ XOR } j = k$. Thus, ignoring factor $\frac{1}{\sqrt{2}}$, this corresponds to XOR.

Similarly, the Hadamard gate, which maps the Hadamard basis into the computational one is the self-conjugated, self-adjoint linear map

$$\boxed{H} := \sum_i \begin{array}{c} \nabla i \\ | \\ \triangle i \end{array}$$

Exercise 2

Show that

$$\begin{aligned}
 \begin{array}{c} | \quad | \\ \nabla B_0 \end{array} &:= \frac{1}{\sqrt{2}} \left(\begin{array}{c} | \quad | \\ \nabla 0 \quad \nabla 0 \end{array} + \begin{array}{c} | \quad | \\ \nabla 1 \quad \nabla 1 \end{array} \right) \\
 \begin{array}{c} | \quad | \\ \nabla B_1 \end{array} &:= \frac{1}{\sqrt{2}} \left(\begin{array}{c} | \quad | \\ \nabla 0 \quad \nabla 1 \end{array} + \begin{array}{c} | \quad | \\ \nabla 1 \quad \nabla 0 \end{array} \right) \\
 \begin{array}{c} | \quad | \\ \nabla B_2 \end{array} &:= \frac{1}{\sqrt{2}} \left(\begin{array}{c} | \quad | \\ \nabla 0 \quad \nabla 0 \end{array} - \begin{array}{c} | \quad | \\ \nabla 1 \quad \nabla 1 \end{array} \right) \\
 \begin{array}{c} | \quad | \\ \nabla B_3 \end{array} &:= \frac{1}{\sqrt{2}} \left(\begin{array}{c} | \quad | \\ \nabla 0 \quad \nabla 1 \end{array} - \begin{array}{c} | \quad | \\ \nabla 1 \quad \nabla 0 \end{array} \right)
 \end{aligned}$$

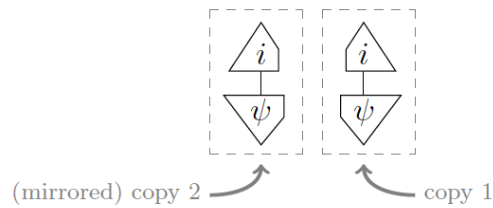
forms a basis for $\mathcal{C}^2 \otimes \mathcal{C}^2$ which does not correspond to a product of two bases for \mathcal{C}^2 .

Linear maps, however, are not full adequate to represent quantum processes. For example, quantum processes are blind to global phases as they are not detected by quantum measurements. On the other hand, composing a state and an effect yields a scalar which

in the theory of linear maps is a complex number, and not a probability as expected by the generalised Born rule. Both these issues are solved by a curious, but straightforward procedure: process doubling.

Process doubling.

Multiplying a scalar by its conjugate, if its state and effect components are both normalised, yields a real number in the interval $[0, 1]$, i.e. a probability. This means that any normalised state along with any orthonormal basis yields a probability distribution, considering *doubled* inner products.

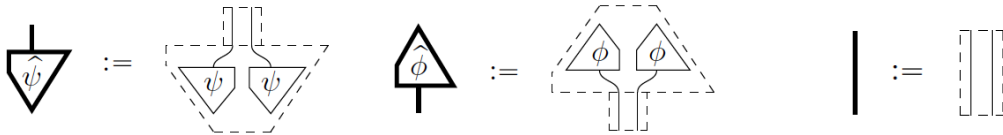


Exercise 3

Verify that

$$\sum_i \begin{array}{c} \triangleup i \\ \psi \\ \triangleleft \end{array} = 1$$

Doubling states and processes amounts to defining for each state ϕ of type \mathbf{A} its doubled version $\hat{\phi}$ of type $\hat{\mathbf{A}}$, which is basically two ‘hidden’ copies of \mathbf{A} . A new, bold notation is used to simplify the diagrams as follows.



Exercise 4

The doubling procedure is related to the passage from a pure state vector $|\phi\rangle$ to the corresponding density operator, which typically plays the role of a quantum state in the usual approaches to quantum computation. Explain why.

Another advantage of the doubling procedure is to eliminate global phases. Indeed, two states become the same state when doubled if and only if they are equal up to some number $e^{i\alpha}$ for $\alpha \in [0, 2\pi[$, i.e.

$$\begin{array}{c} \downarrow \\ \psi \end{array} \begin{array}{c} \downarrow \\ \psi \end{array} = \begin{array}{c} \downarrow \\ \phi \end{array} \begin{array}{c} \downarrow \\ \phi \end{array} \iff \begin{array}{c} \downarrow \\ \psi \end{array} = e^{i\alpha} \begin{array}{c} \downarrow \\ \phi \end{array}$$

The doubling procedure extends to processes in the obvious way

$$\text{double} \left(\begin{array}{c} \downarrow \\ f \end{array} \right) := \begin{array}{c} \downarrow \\ \hat{f} \end{array} = \begin{array}{c} \downarrow \\ \begin{array}{|c|c|} \hline f & f \\ \hline \end{array} \end{array}$$

Care is needed when dealing with processes with several inputs/outputs: pins must be connected in order and taking into consideration that the conjugate of a process is pictured as its mirror image — so, inputs and outputs are counted from right-to-left rather than left-to-right, which introduces a ‘twist’ in the wires connected to the conjugate process.

$$\begin{array}{c} n \quad 2 \quad 1 \\ \vdots \\ \downarrow \\ f \\ \vdots \\ m \quad 2 \quad 1 \end{array} \quad \begin{array}{c} 1 \quad 2 \quad n \\ \vdots \\ \downarrow \\ f \\ \vdots \\ 1 \quad 2 \quad m \end{array} \quad \begin{array}{c} \downarrow \\ \hat{f} \end{array} := \begin{array}{c} \downarrow \\ \begin{array}{|c|c|} \hline f & f \\ \hline \end{array} \end{array}$$

This procedure yields a theory of doubled processes which will be referred to in the sequel as *pure quantum maps*. As a process theory this is simply a subtheory of linear maps built as follows: its types are \hat{A} for all Hilbert spaces A , and its processes are $\hat{f} : \hat{A} \rightarrow \hat{B}$ for all linear maps $f : A \rightarrow B$. Being a subtheory of linear maps, pure quantum maps admits string diagrams. In particular,

$$\cup := \text{double} \left(\cup \right) = \begin{array}{c} \downarrow \\ \cup \end{array} = \begin{array}{c} \downarrow \\ \begin{array}{|c|c|} \hline \cup & \cup \\ \hline \end{array} \end{array}$$

and

$$\cap := \begin{array}{c} \downarrow \\ \cap \end{array}$$

Exercise 5

Verify the yanking laws in the theory of pure quantum maps.

Exercise 6

Show that doubling preserves sequential and parallel composition of processes.

Exercise 7

With caps and cups one builds a notion of transposition:

$$\begin{array}{c} | \\ \text{---} \\ \hat{f} \\ \text{---} \\ | \end{array} \mapsto \begin{array}{c} | \\ \text{---} \\ \hat{f} \\ \text{---} \\ | \end{array} := \begin{array}{c} | \\ \text{---} \\ \hat{f} \\ \text{---} \\ | \end{array}$$

Verify that it coincides with transposition in the original theory, i.e.

$$\text{double} \left(\begin{array}{c} | \\ \text{---} \\ f \\ \text{---} \\ | \end{array} \right) = \begin{array}{c} | \\ \text{---} \\ \hat{f} \\ \text{---} \\ | \end{array}$$

One may thus conclude that doubling preserves string diagrams. Thus, any of the calculations done for diagrams of linear maps lifts to pure quantum maps by doubling all of the diagrams. The converse is a bit more tricky: to go back to the theory of linear maps one needs to reintroduce global phases. Formally, for D, D' arbitrary diagrams in the theory of pure quantum maps,

$$\left(\exists e^{i\alpha} : \begin{array}{c} \dots \\ | \\ \text{---} \\ D \\ \text{---} \\ | \\ \dots \end{array} = e^{i\alpha} \begin{array}{c} \dots \\ | \\ \text{---} \\ D' \\ \text{---} \\ | \\ \dots \end{array} \right) \iff \begin{array}{c} \dots \\ | \\ \text{---} \\ \hat{D} \\ \text{---} \\ | \\ \dots \end{array} = \begin{array}{c} \dots \\ | \\ \text{---} \\ \hat{D}' \\ \text{---} \\ | \\ \dots \end{array}$$

To verify the right-to-left implication it suffices to show that for any linear maps f and g such that

$$\begin{array}{c} | \\ \text{---} \\ \hat{f} \\ \text{---} \\ | \end{array} = \begin{array}{c} | \\ \text{---} \\ \hat{g} \\ \text{---} \\ | \end{array}$$

there exists a α such that $f = e^{i\alpha}g$. Consider the following two scalars

$$\lambda := \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ | \\ \text{---} \\ \text{---} \\ \text{---} \end{array} \quad \mu := \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ | \\ \text{---} \\ \text{---} \\ \text{---} \end{array}$$

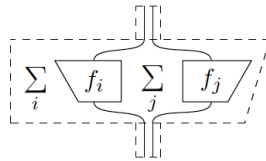
Exercise 8

Prove that \hat{f} is unitary iff f is.

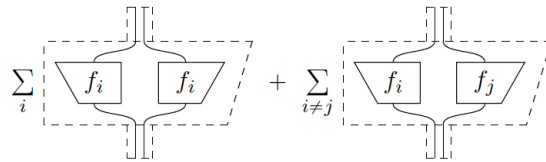
Exercise 9

Prove that \hat{f} is a projector iff there exists a projector g such that $\hat{f} = \hat{g}$.

Note, however, that sums in the doubled theory are not doubled sums. Actually, doubling a sum may not even lead to a pure quantum process. This is because doubling a sum involves two independent summations (i.e., over different indices):



which can be split into two sums for $i = j$ and $i \neq j$:



Thus,

$$\text{double} \left(\sum_i \begin{array}{|c|} \hline \diagdown \\ \hline f_i \\ \hline \diagup \\ \hline \end{array} \right) = \sum_i \begin{array}{|c|} \hline \diagdown \\ \hline \hat{f}_i \\ \hline \diagup \\ \hline \end{array} + \sum_{i \neq j} \begin{array}{|c|} \hline \diagdown \\ \hline f_i \\ \hline \diagup \\ \hline \end{array} \begin{array}{|c|} \hline \diagdown \\ \hline f_j \\ \hline \diagup \\ \hline \end{array}$$

In general this extra term will not be 0. For example, for $\lambda_1 = \lambda_2 = 1$, one gets

$$\text{double} \left(\sum_i \begin{array}{|c|} \hline \diagdown \\ \hline \lambda_i \\ \hline \diagup \\ \hline \end{array} \right) = \text{double}(1 + 1) = 4 \neq 2 = 1 + 1 = \sum_i \begin{array}{|c|} \hline \diagdown \\ \hline \lambda_i \\ \hline \diagup \\ \hline \end{array}$$

This is a fundamental observation: sums in the theory of linear maps capture quantum *superpositions*. In the theory of pure quantum maps on the other hand, sums capture uncertainty about which process actually happened. They correspond to *mixed* quantum states.

Another element that is not preserved by doubling concerns orthonormal bases. Actually, doubling a basis with more than one state does not yield a basis in the theory of pure

quantum maps. A counterexample is provided by the following states over a basis \mathcal{B} .

$$\begin{array}{c} \downarrow \\ \psi \end{array} := \sum_j \begin{array}{c} \downarrow \\ j \end{array} \quad \begin{array}{c} \downarrow \\ \phi \end{array} := \sum_j e^{i\alpha_j} \begin{array}{c} \downarrow \\ j \end{array}$$

where all parameters α_i are distinct. Since ϕ has at least two terms with non-equal coefficients $e^{i\alpha_j}$, the two states are not within a global phase of each other, which leads us to conclude that

$$\begin{array}{c} \downarrow \\ \hat{\psi} \end{array} \neq \begin{array}{c} \downarrow \\ \hat{\phi} \end{array}$$

However one can show that

$$\begin{array}{c} \downarrow \\ \hat{\psi} \\ \downarrow \\ \hat{\psi} \end{array} = \begin{array}{c} \downarrow \\ \hat{\phi} \\ \downarrow \\ \hat{\phi} \end{array}$$

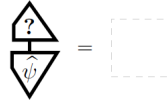
which implies that doubling \mathcal{B} does not yield a basis. It can be proved, however, that each basis for a type \mathcal{A} in linear maps can be extended to a (non-orthogonal) basis in linear maps for the type $\mathcal{A} \otimes \mathcal{A}$, consisting entirely of pure quantum maps. So, in particular, this new basis is also a basis in the theory of pure quantum maps for the type $\hat{\mathcal{A}}$.

Final remarks.

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- Despite some exceptions discussed above, almost all the structure of the theory of linear maps is preserved by doubling. This means that we can work with pure quantum maps as if they were just linear maps. Moreover, Born rule comes in a very straightforward way: computing the internal product of doubled states.
 - ... And we can do it with an extra advantage: no need to be concerned with global phases. Indeed, doubling gives us an alternative to the usual convention regarding quantum states as equivalence classes of states that are equal up to a global phase. Global phases are ignored as they cannot be discovered by measurement. But this fact is already clear in this setting, because the only way to obtain probabilities is through the Born rule, which 'only' makes sense in a doubling setting.
 - A final stage is still in order: going from pure quantum states to arbitrary such states. Therefore, we need to introduce *discarding*, to be able to ignore part of a system, thus capturing our lack of knowledge about its state. Such (impure) quantum maps can alternatively be described as probabilistic mixtures.

As discussed in the next lecture, this is achieved through the introduction of a *discarding* effect providing a test which succeeds with certainty but does not reveal anything about (and, of course, does not depend on) the (normalised) state that gets

discarded³.



References

- [1] B. Coecke and A. Kissinger. *Picturing Quantum Processes: A First Course in Quantum Theory and Diagrammatic Reasoning*. Cambridge University Press, 2017.

³The use of normalised states is essential here: suppose a discarding effect exists, thus yielding 1 when applied to an arbitrary state $\hat{\phi}$. Clearly the application to e.g. $5\hat{\phi}$ would return 5 rather than 1 ...