

Configurations of Web Services

Marco Barbosa¹ Luís Barbosa¹

¹Departamento de Informática
Universidade do Minho
Braga - Portugal

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Two views on Componentware

Introduction

- The OO legacy
 - components as (collections of) classes/objects
 - method invocation as the kernel of component composition
 - resort to middleware intervention to loosen tight-coupling
- The Coordination Paradigm View
 - temporal/spatial decoupling to support a looser inter-component dependency
 - amenable to external control
 - requires anonymous communication

Aims

Aims

- Discuss an orchestration model combining REO-like connectors with behaviourally annotated interfaces
- Configuration = Components + Connectors + Glue code
- Tentative application to model configurations of web-services

Interface

Definition

A web-service S interface is specified by

- a *port signature*, $\text{sig}(S)$ over \mathbb{D} , given by a port name and a polarity annotation (either in(put) or out(put))
- a *use pattern*, $\text{use}(S)$, given by a process term over port names.

Generic Process Algebra

cf. "Process Algebra à la Bird-Merteens" [Bar01, RBB06]

- Processes are inhabitants of a final coalgebra;
- Combinators defined by coinductive extension;
- Interaction discipline: θ

Interaction Structure

- In CCS, $a\theta\bar{a} = \tau$
- In CSP, $a\theta a = a$ for all action $a \in Act$.
- In architectural configurations, ...
 - allow different interaction disciplines to coexist.

Use Patterns and Interaction

- Let \mathcal{P} be a set of port identifiers and S a (the specification of) a web service. Its use pattern $use(S)$ is given by a process expression over $Act = \mathcal{P}(\mathcal{P})$, given by

$$P ::= \mathbf{0} \mid \alpha.P \mid P + P \mid P \otimes P \mid P \parallel P \mid P; P \mid P \mid P \mid \sigma P \mid \text{fix } (x = P)$$

- choosing $Act = \mathcal{P}(\mathcal{P})$ allows for the synchronous activation of several ports in a single computational step

Use Patterns and Interaction

- All interaction between web services is mediated by a specific connector
- Therefore, if two web services are active their joint behaviour will allow the realization of both use patterns either simultaneously or in an independent way:
- The joint behaviour of a collection $\{S_i \mid i \in n\}$ of ws is

$$use(S_1) \mid \dots \mid use(S_n)$$

where the interaction discipline is fixed by $\theta = \cup$.

Examples

$$use(S_1) = \text{fix } (x = a.x + b.x)$$

$$use(S_2) = \text{fix } (x' = cd.x'), \text{ where, } cd \stackrel{\text{abv}}{=} \{c, d\}$$

$$use(S_1) \mid use(S_2) = \text{fix } (x = acd.x + bcd.x + a.x + b.x + cd.x)$$

Services are *coordinated* via Connectors

- What are connectors and how do they compose?
- How do web services' interfaces and connectors interact in a configuration?

Connectors

A connector \mathbb{C} is defined through:

- a relation $\text{data}.\llbracket\mathbb{C}\rrbracket : \mathbb{D}^m \longleftarrow \mathbb{D}^n$ which records the flow of data;
- a process expression $\text{port}.\llbracket\mathbb{C}\rrbracket$ which gives the pattern of port activation.

Connectors

Basic Connectors

$$\text{data.}[\![\bullet \xrightarrow{\quad} \bullet]\!] = \text{Id}_{\mathbb{D}}, \quad \text{port.}[\![\bullet \xrightarrow{\quad} \bullet]\!] = \text{fix } (x = ab.x)$$

$$\text{data.}[\![\bullet \xrightarrow{\diamond} \bullet]\!] \subseteq \text{Id}_{\mathbb{D}}, \quad \text{port.}[\![\bullet \xrightarrow{\diamond} \bullet]\!] = \text{fix } (x = ab.x + a.x)$$

$$\text{data.}[\![\bullet \xrightarrow{\nabla} \bullet]\!] = \mathbb{D} \times \mathbb{D}, \quad \text{port.}[\![\bullet \xrightarrow{\nabla} \bullet]\!] = \text{fix } (x = ab.x)$$

$$\text{data.}[\![\bullet \xrightarrow{\nabla} \bullet]\!] = \mathbb{D} \times \mathbb{D}, \quad \text{port.}[\![\bullet \xrightarrow{\nabla} \bullet]\!] = \text{fix } (x = a.x + b.x)$$

$$\text{data.}[\![\bullet \xrightarrow{\square} \bullet]\!] = \text{Id}_{\mathbb{D}}, \quad \text{port.}[\![\bullet \xrightarrow{\square} \bullet]\!] = \text{fix } (x = a.b.x)$$

Connector Combinators

Aggregation

This combinator places its arguments side-by-side, with no direct interaction between them:

$\text{port.}[[C_1 \boxtimes C_2]] = \text{port.}[[C_1]] \mid \text{port.}[[C_2]]$ with $\theta = \cup$

Combinators

Hook

Acts as a *feedback* mechanism.

On the *data* side:

Suppose data. $\llbracket C \rrbracket = R : \mathbb{D}^n \longleftarrow \mathbb{D}^m$. Then,

$$R \swarrow_i^j : \mathbb{D}^{n-1} \longleftarrow \mathbb{D}^{m-1}$$

$$t = t_m, \dots, t_{i+i}, t_{i-i}, \dots, t_0, \text{ and } t' = t'_n, \dots, t'_{j+i}, t'_{j-i}, t'_0$$

$$t(R \swarrow_i^j) t' \text{ iff}$$

$$\exists x. (t_n, \dots, t_{i+i}, x, t_{i-i}, \dots, t_0) R (t_m, \dots, t_{j+i}, x, t_{j-i}, \dots, t_0)$$

Combinators

Hook

On the behavioural side:

$\text{port.}[\![C \uparrow_i^j]\!]_i$ is obtained from $\text{port.}[\![C]\!]_i$, by deleting references to ports i and j .

To be well-formed it is required that i and j appear in different factors of some form of parallel composition (\parallel , \otimes , or $|$).

Combinators

Join

Plugs ports with same polarity

- Right Join: $(\mathbb{C}^i > z)$ (non deterministic merger)
- Left Join: $(z <^i \mathbb{C})$ (broadcaster)

At the behavioural level, both operators act as *port renamers*

$$\text{port.}[(\mathbb{C}^i > n)] = \text{port.}[(n <^i \mathbb{C})] = \{n \leftarrow i, n \leftarrow j\} \text{port.}[\mathbb{C}]$$

Configurations

Configuration Structure

A configuration involving a collection $S = \{S_i \mid i \in n\}$ of web services is a tuple

$$\langle U, \mathbb{C}, \sigma \rangle$$

where

- $U = use(S_1) \mid use(S_2) \mid \dots \mid use(S_n)$ is the (joint) use pattern for S
- \mathbb{C} is a connector
- σ a mapping of ports in S to ports in \mathbb{C}

Configuration Behaviour

The behaviour $bh(\Gamma)$ of a configuration $\Gamma = \langle U, \mathbb{C}, \sigma \rangle$ is given by

$$bh(\Gamma) = \sigma U \otimes \text{port.}[\mathbb{C}]$$

where θ underlying the \otimes connective is given by

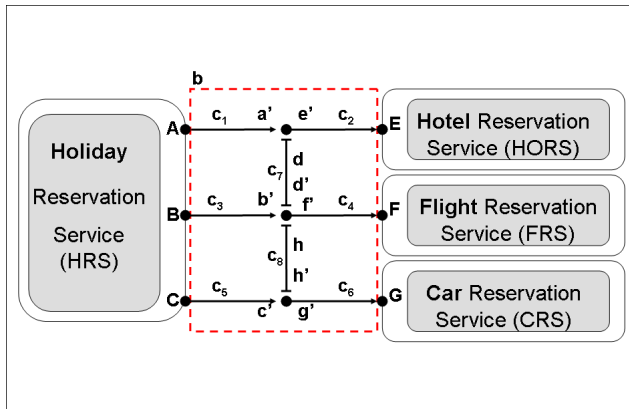
$$c \theta c' = \begin{cases} c \cap (c' \cup \text{free}) & \Leftarrow c' \subseteq c \\ \emptyset & \Leftarrow \text{otherwise} \end{cases}$$

and free denotes the set of unplugged ports in U , *i.e.*, not in the domain of mapping σ .

Configuration Behaviour: intuitions

- Interaction is achieved by the simultaneous activation of identically named ports
- There is no interaction if the connector intends to activate ports which are not linked to the ones offered by the web services' side.
- The dual situation is allowed, *i.e.*, if the web services' side offers activation of all ports plugged to the ones offered by the connectors' side, their intersection is the resulting interaction.
- Moreover activation of unplugged web services' ports is always possible.

Holiday Reservation (cf. [DA04])



Holiday Reservation

Configuration

$HR = \langle WHR, SB, \sigma_{HS} \rangle$, where

$WHR = use(HRS) \mid use(HORS) \mid use(FRS) \mid use(CRS)$

$\sigma_{HS} = \{a \leftarrow A, b \leftarrow B, c \leftarrow C, d \leftarrow D, e \leftarrow E, f \leftarrow F, g \leftarrow G\}$

Usage

$use(HRS) = \text{fix } (x = a.x + b.x + c.x + abc.x)$

$use(HORS) = \text{fix } (x = e.x)$

$use(FRS) = \text{fix } (x = f.x)$

$use(CRS) = \text{fix } (x = g.x)$

Holiday Reservation

Connector

$\text{port.}[[c_1]] = \text{fix } (x = aa'.x), \text{ port.}[[c_2]] = \text{fix } (x = e'e.x),$
 $\text{port.}[[c_3]] = \text{fix } (x = bb'.x), \text{ port.}[[c_4]] = \text{fix } (x = f'f.x),$

...

$Cn_1 = \text{port.}[[(n <_d^{e'} (c_2 \boxtimes c_7))]] = \text{fix } (x = \text{end}'.x)$

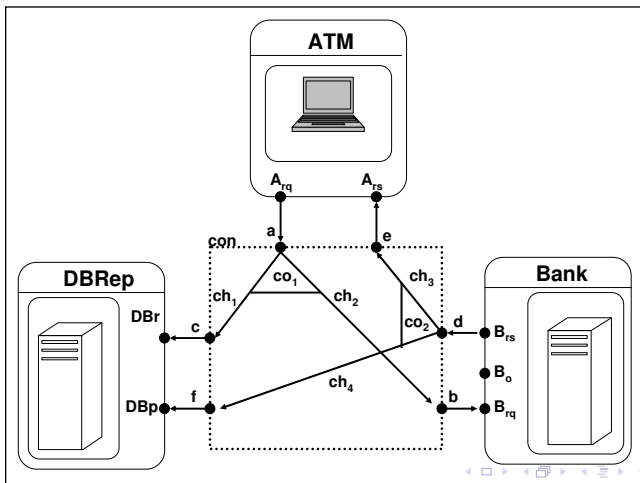
$Cn_2 = \text{port.}[[(c_1 \boxtimes Cn_1) \uparrow_a^n]] = \text{fix } (x = aed'.x)$

...

$\text{port.}[[b]] = \text{fix } (x = abcefg.x),$ and finally

$bh(HR) = \text{fix } (x = abcefg.x)$

Bank System



Bank System

Configuration

$BS = \langle WBS, DBC, \sigma_{BS} \rangle$, where

$WBS = use(ATM) \mid use(Bank) \mid use(DBRep)$

$\sigma_{HS} = \{a \leftarrow A_{rq}, e \leftarrow A_{rs}, c \leftarrow DB_r, f \leftarrow DB_p, d \leftarrow B_{rs}, b \leftarrow B_{rq}\}$

Bank System

Use Patterns

$$use(ATM) = \text{fix } (x = a.e.x)$$
$$use(Bank) = \text{fix } (y = b.d.y)$$
$$use(DBRep) = \text{fix } (z = c.z + f.z)$$

Bank System

Configuration

$$\begin{aligned} \text{port.}[[DBC]] &= \text{port.}[[(co_1 \boxtimes co_2)]] = \\ &\quad \text{fix } (x = abc.x + def.x + abcdef.x) \\ bh(BS) &= \text{fix } (x = abc.def.x) \end{aligned}$$

Conclusions and Future Work

Conclusions

- Formal model for behavioural interfaces and configurations
- Exogenous coordination (cf, REO model)
- Role of *generic process algebra* [Bar01,RBB06]
(cf, coexistence of different interaction disciplines)

Conclusion and Future Work

Future Work

- Expressing *workflow patterns*
 - **Pattern 2 (Parallel split)**

$$use(WS_2) = P_1 \mid \dots \mid P_n$$

- **Pattern 3 (Synchronization)**

$$use(WS_3) = (a_1.a_2.\dots.a_n.\S \otimes b_1.b_2.\dots.b_n.\S); P$$

- Is this framework suitable for expressing the semantics of orchestration languages?
- if so, how easily can properties be proved?