

Michael de Oliveira

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Bachelor Engineering Physics – University of Minho



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Master Engineering Physics with a specialization in Quantum computing – University of Minho



Thesis titled “On Quantum Bayesian Decision Making”



FUNDAÇÃO
CALOUSTE GULBENKIAN

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PhD in theoretical computer science – Quantum computing - International Iberian Nanotechnology Laboratory & Sorbonne University



Thesis titled “The Interplay between Quantum Foundations and Circuit Complexity”



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Thesis titled “The Interplay between Quantum Foundations and Circuit Complexity”



Upcoming- Research Scientist Foxconn 

Research on quantum complexity theory and quantum algorithms

My Master's thesis

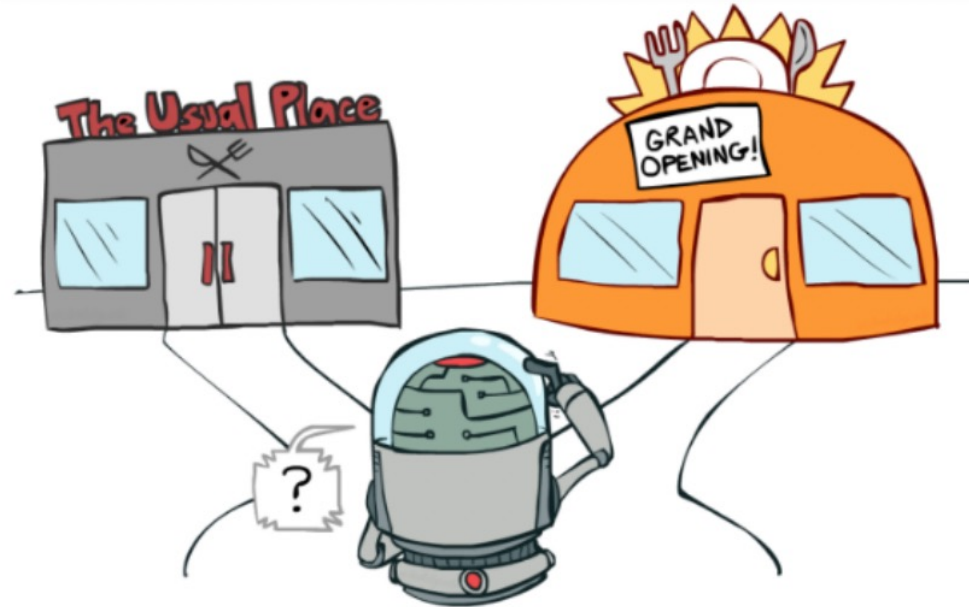
Automated decision making processs

We are presented every day with multiple decision to make.



Automated decision making processs

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However, we would like to delegate most of them to a kind of personal assistant who can make near-optimal decisions for us.

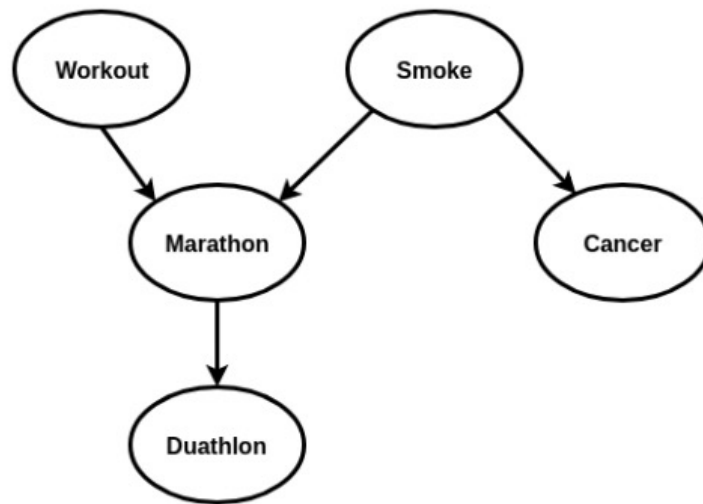
Automated decision making processs

First, we need a mathematical model that can describe probabilistic systems to handle uncertainty.

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Bayesian networks



Workout	$P(\text{Workout})$
True	0.7
False	0.3

Smoke	$P(\text{Smoke})$
True	0.6
False	0.4

Smoke	Workout	Marathon	$P(\text{Marathon} \text{Workout}, \text{Smoke})$
True	True	True	0.2
True	True	False	0.8
True	False	True	0.05
True	False	False	0.95
False	True	True	0.4
False	True	False	0.6
False	False	True	0.1
False	False	False	0.9

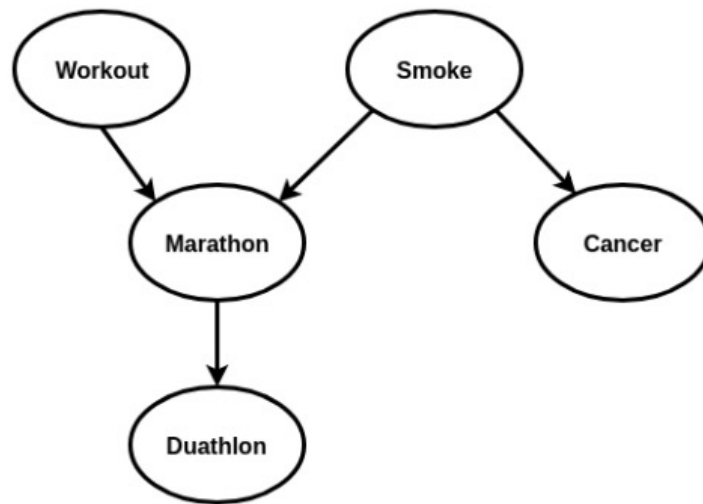
Smoke	Lung cancer	$P(\text{Lung cancer} \text{Smoke})$
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True	False	0.98
False	True	0.005
False	False	0.995

Marathon	Duathlon	$P(\text{Duathlon} \text{Marathon})$
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True	False	0.3
False	True	0
False	False	1

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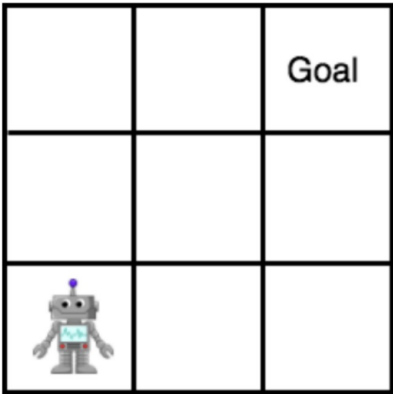
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Marathon	Duathlon	P(Duathlon Marathon)
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True	False	0.3
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$$P(x_1, \dots, x_n) = \prod_{i=1}^n P(x_i | \text{Parents}(X_i))$$

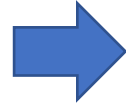
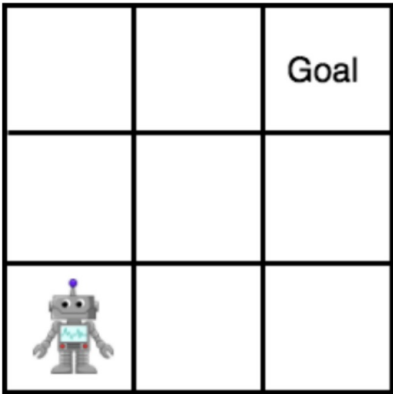
Automated decision making processs

Problem

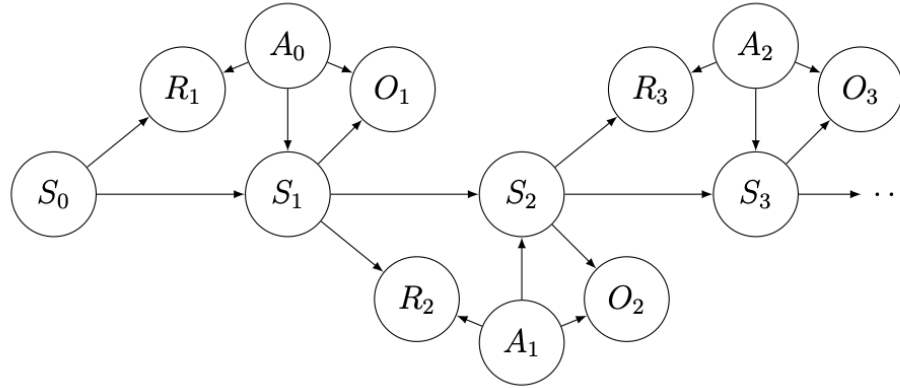


Automated decision making processes

Problem

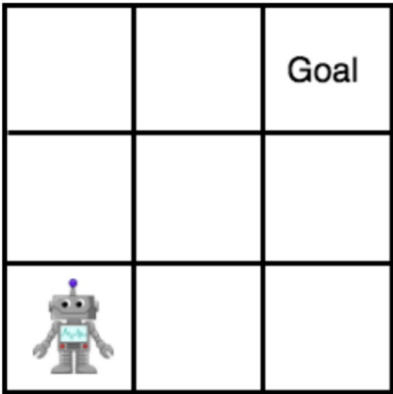


Model/abstraction

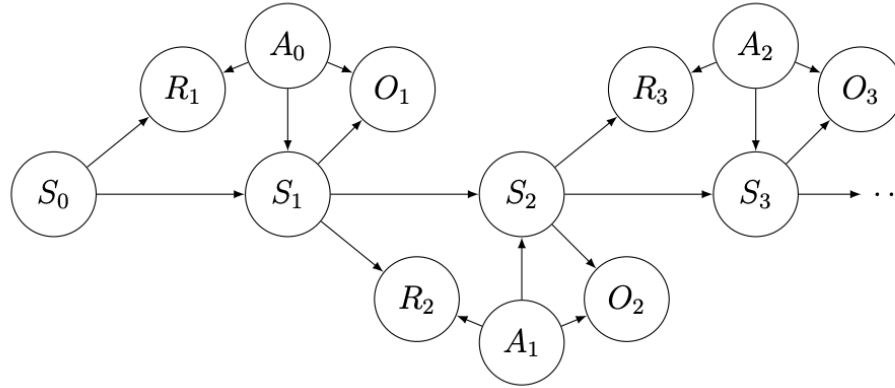


Automated decision making processes

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Model/abstraction



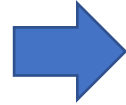
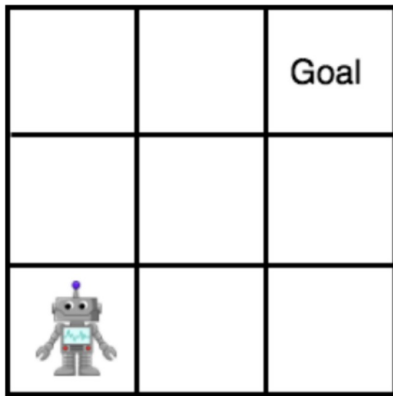
Decision making process



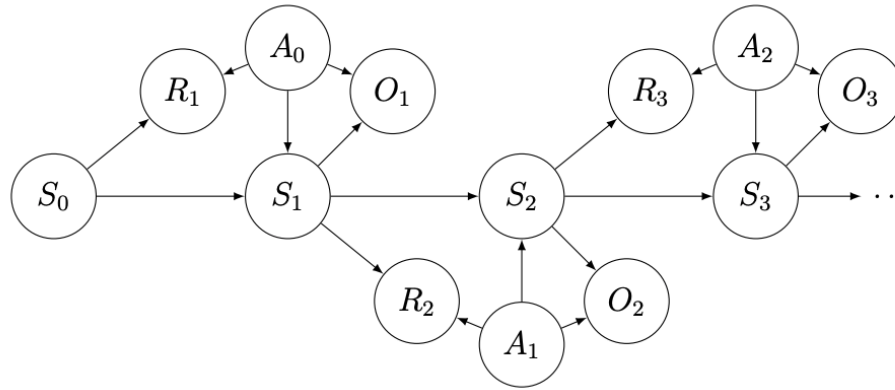
Compute A_1, A_2, \dots, A_n
s.t. $R_T = \sum_{i=1}^n R_i$ is maximal

Automated decision making processes

Problem



Model/abstraction



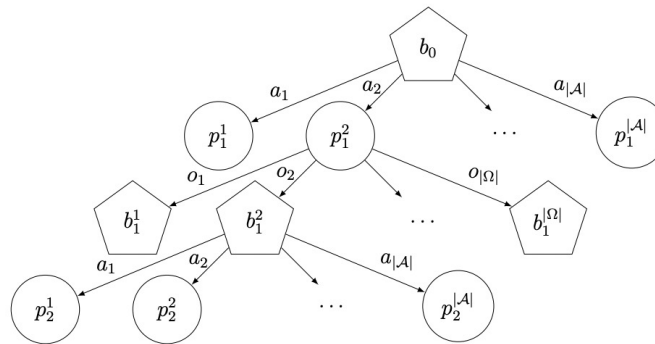
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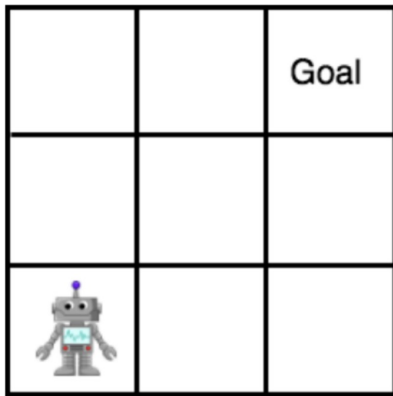


It exponentially harder to
interpolate over each
additional time step

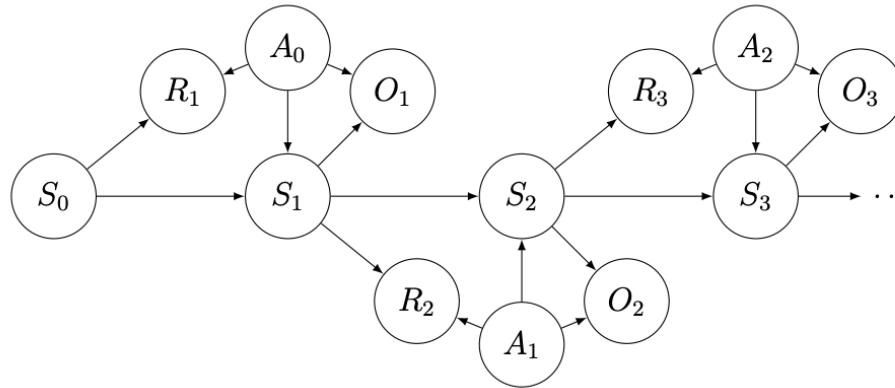


Automated decision making processes

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Model/abstraction

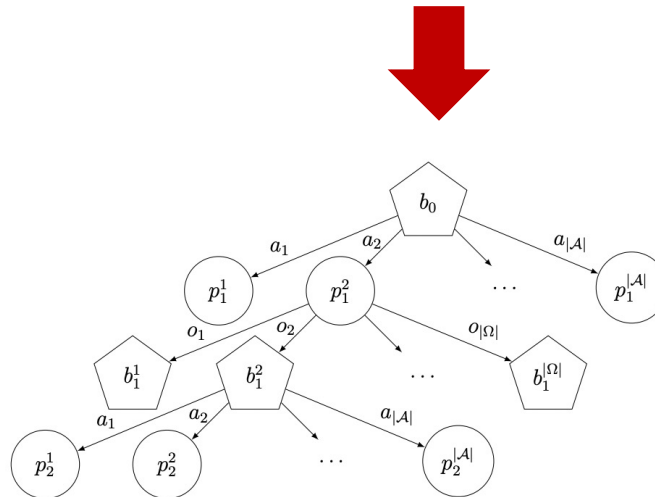


Decision making process



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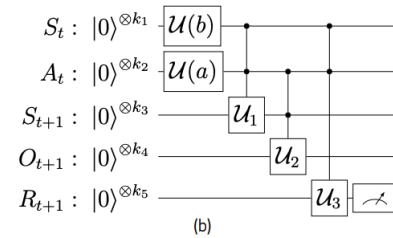
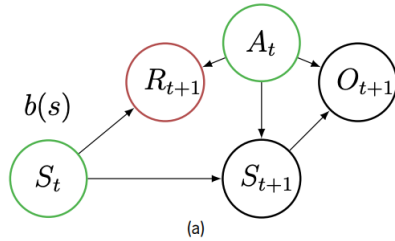
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Naïve quantum Bayesian decision making process

Quantum bayesian network states

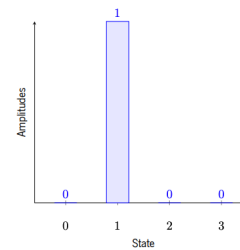
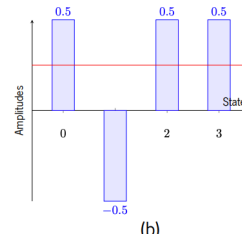
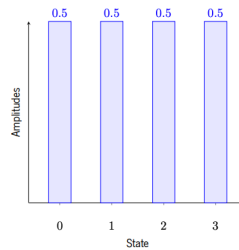
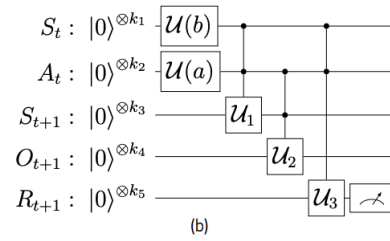
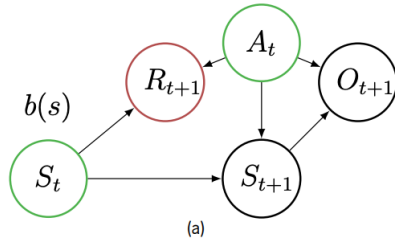


Naïve quantum Bayesian decision making process

Quantum bayesian
network states

+

Amplitude
amplification
algorithm

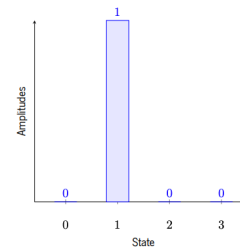
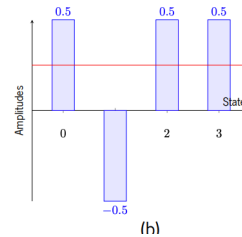
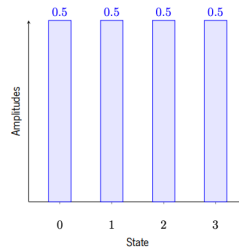
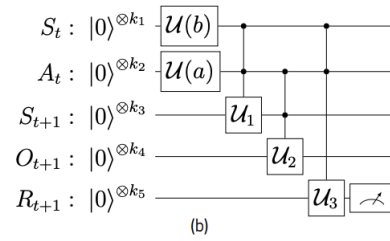
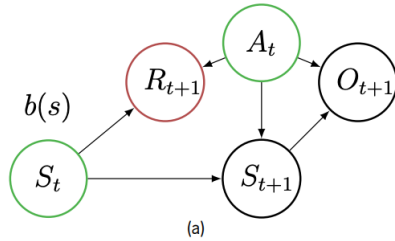


Naïve quantum Bayesian decision making process

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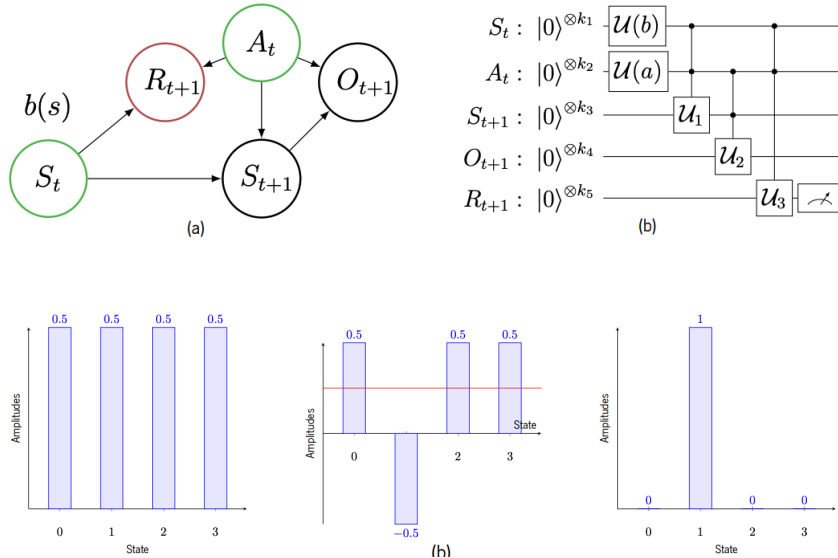
	Classical	Quantum
Complexity	$\mathcal{O}(NMP(e)^{-1})$	$\mathcal{O}\left(N2^M P(e)^{-\frac{1}{2}}\right)$

Naïve quantum Bayesian decision making process

Quantum bayesian
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	Classical	Quantum
Complexity	$\mathcal{O}(NMP(e)^{-1})$	$\mathcal{O}\left(N2^M P(e)^{-\frac{1}{2}}\right)$

$$EU(a|e) = \sum_r \underbrace{P(\text{Result} = r|a, e)}_{\text{Quantum}} * \underbrace{U(r)}_{\text{Classical}}$$

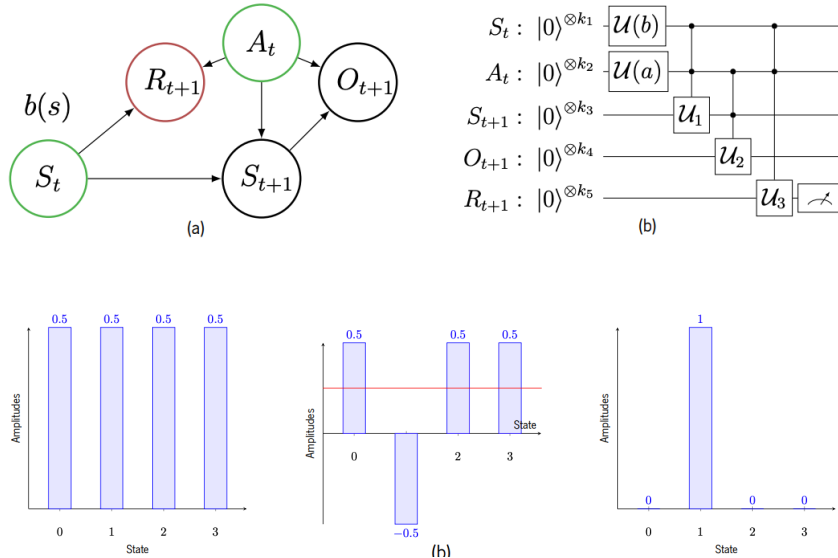
$$\text{action} = \text{argmax}_a EU(a|e)$$

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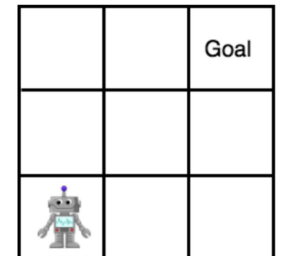
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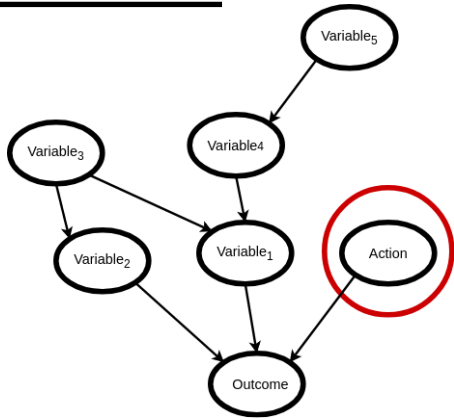


We obtain subquadratic quantum advantages for sparse Bayesian networks.



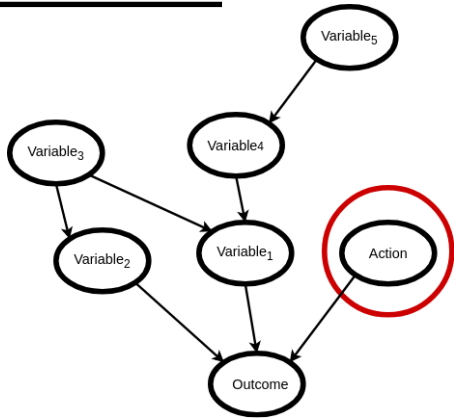
A new Quantum Bayesian Decision-Making

Condition



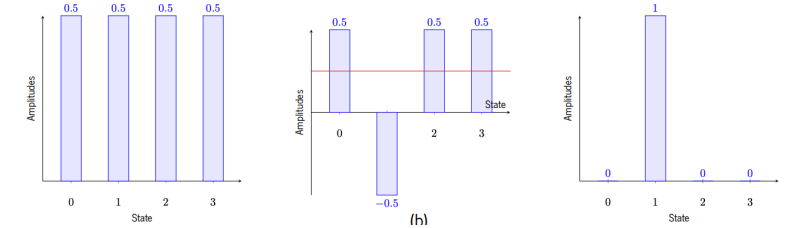
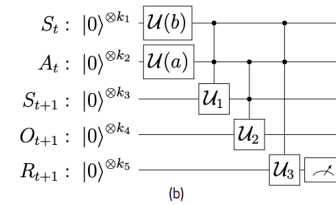
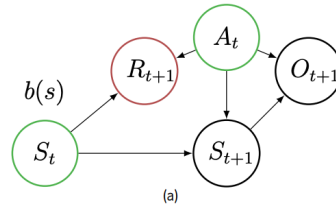
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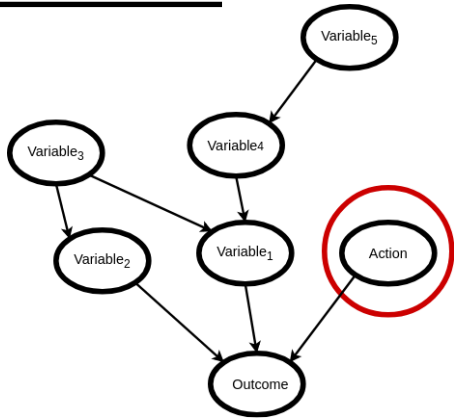
Solution

- We encode the BN but do not select an action
- We amplify the observation v . and the outcome v . based on their utility.



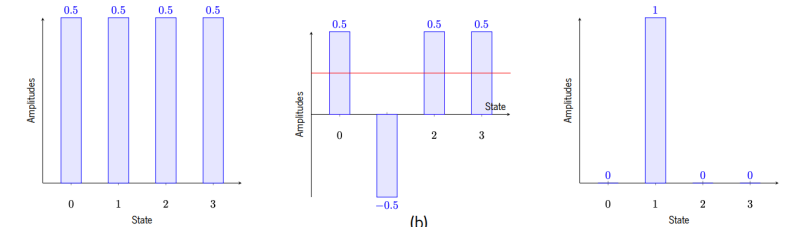
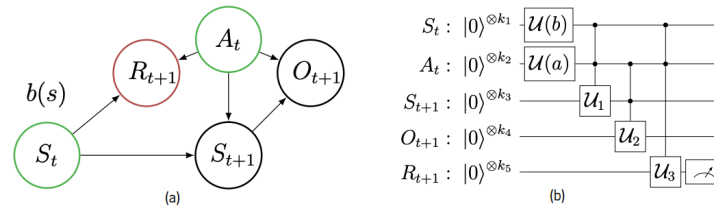
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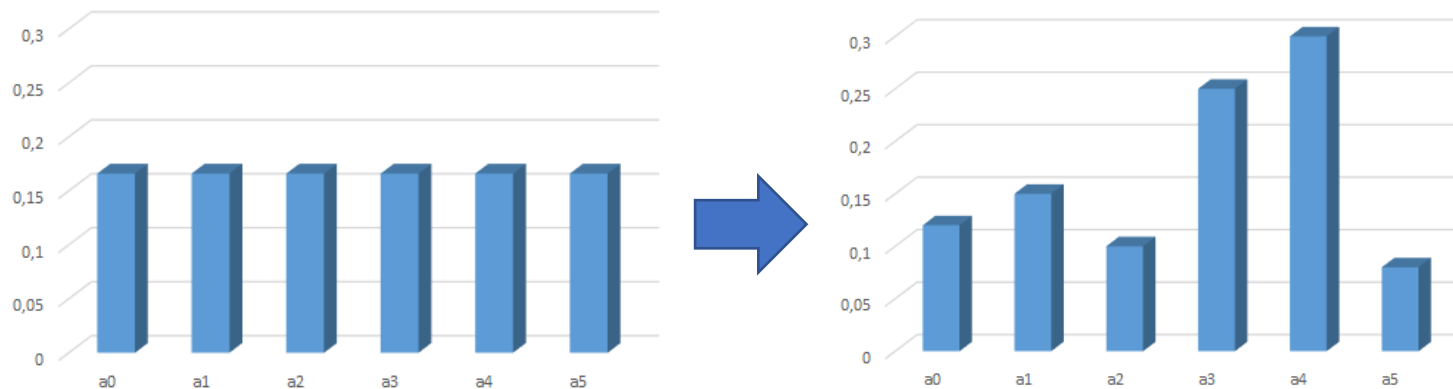


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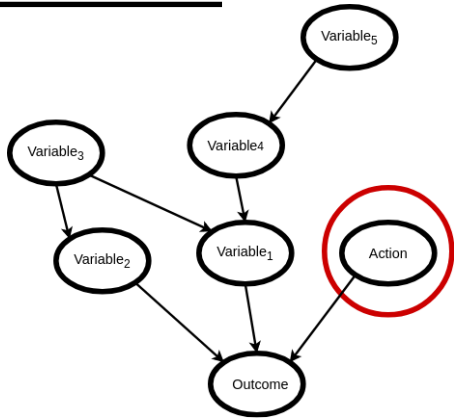


Consequence



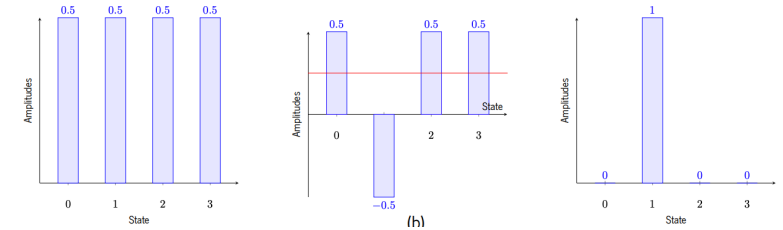
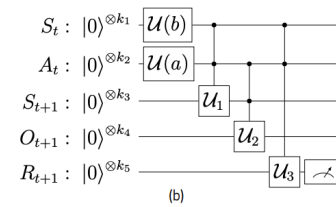
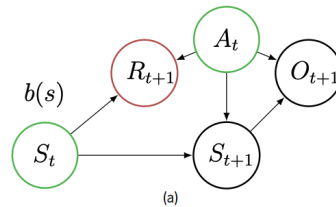
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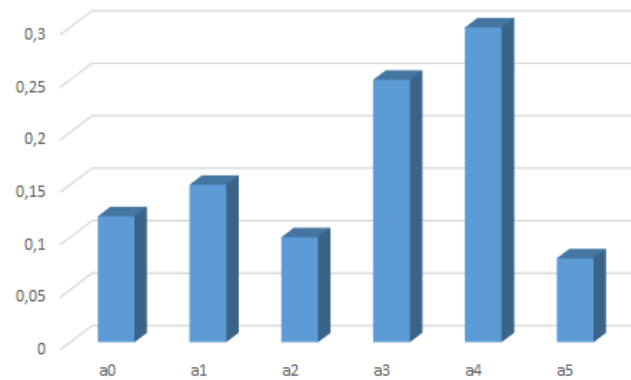
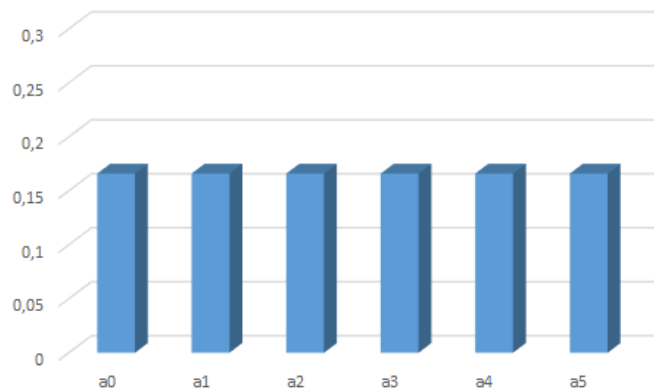


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Consequence



Result

For sparse Bayesian networks, we obtain a subquadratic quantum advantage but a strictly larger advantage than naïve amplitude amplification.

$$\frac{C}{Q_{naive}} = n^{c_1}, \quad \frac{C}{Q_{new}} = n^{c_2}$$

$$c_1 < c_2$$

On Quantum Bayesian Decision-Making

This work and its results gave origin to,

On Quantum Bayesian Decision-Making

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Presentation at an international
workshop (Santiago Chile)

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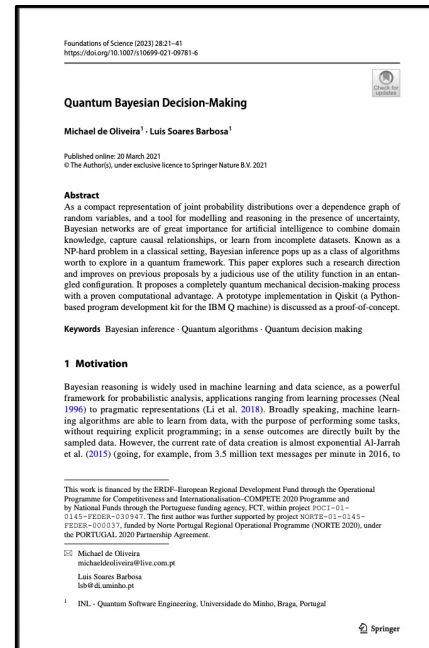
A paper in a scientific journal

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Content of my Gulbenkian research project and master's thesis

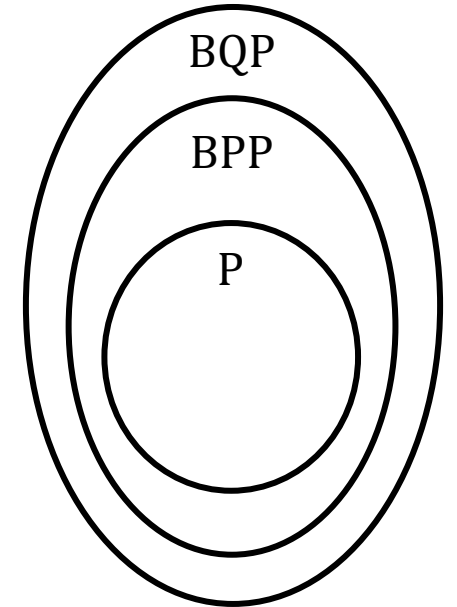
My Phd thesis and the quest for realizable
quantum computational advantages

Quantum advantage computational advantage

Main interest in quantum comes from the fact that these devices have **computation complexity theoretic implications**.

Quantum advantage computational advantage

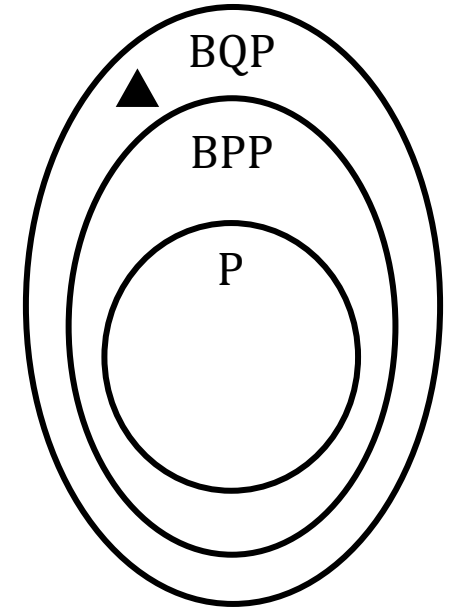
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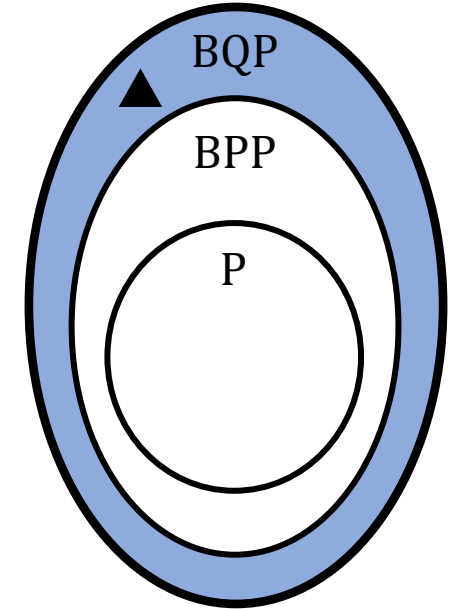
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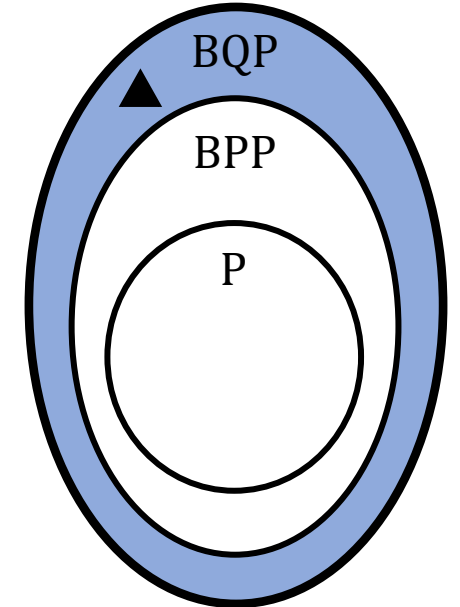
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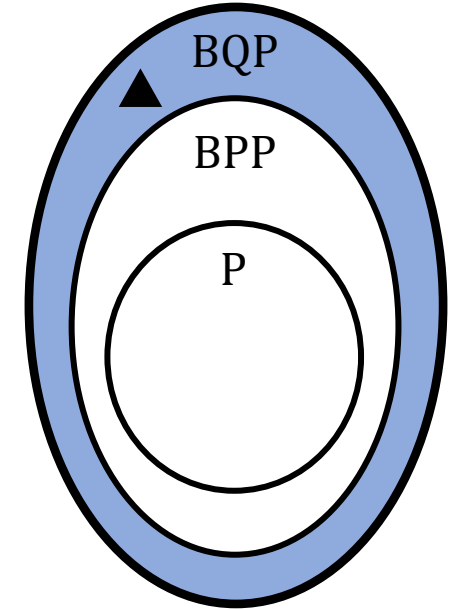
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 - Nevertheless, these quantum algorithms require **large-scale fault-tolerant quantum devices** (1700 qubits, 10^{36} Toffoli gates).



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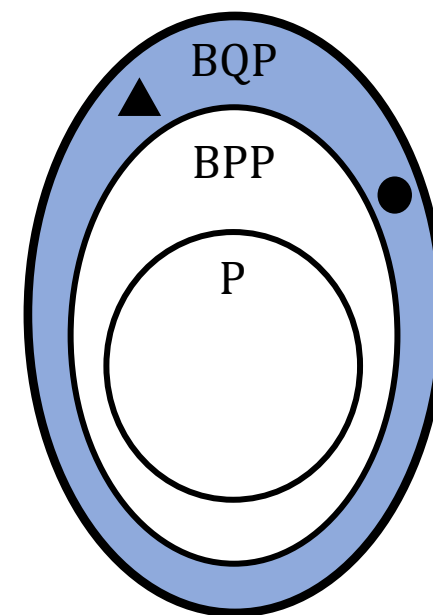
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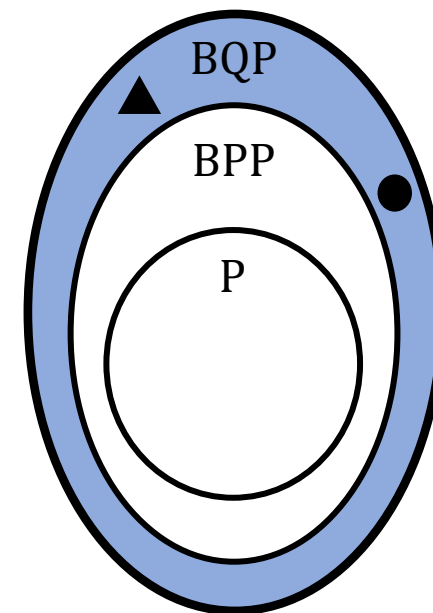
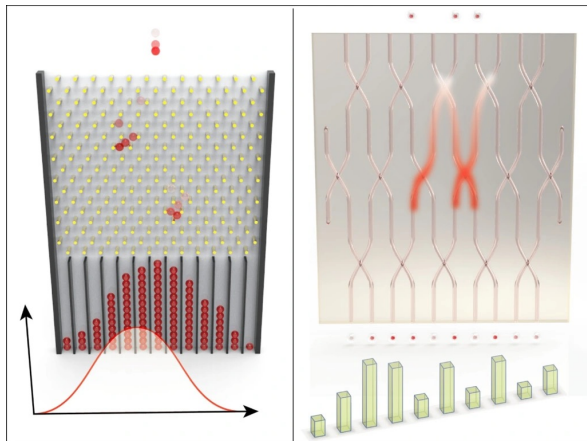
These algorithms (including Grover type search) are **not within reach in the next years ...**

Quantum advantage within NISQ or LISQ regime



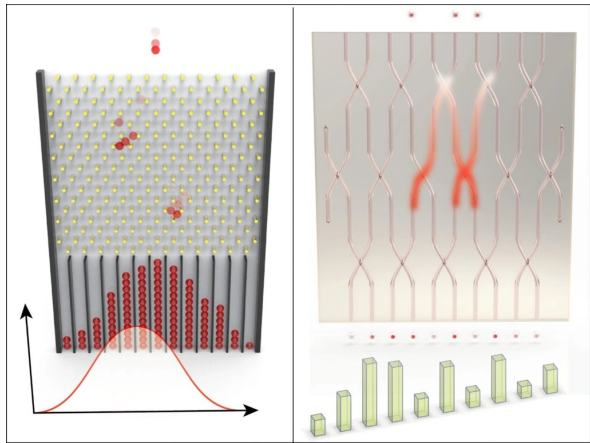
Quantum advantage within NISQ or LISQ regime

● Boson Sampling

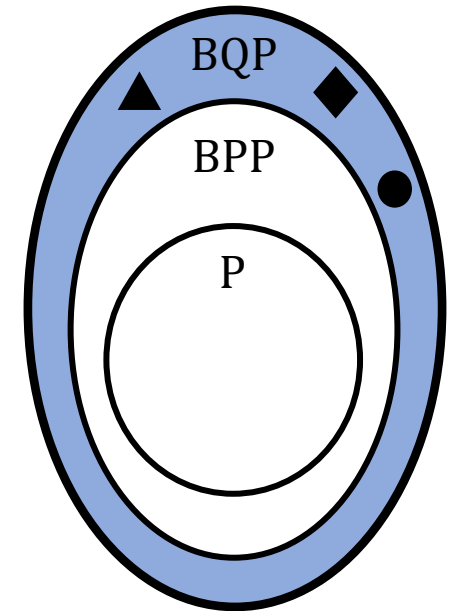
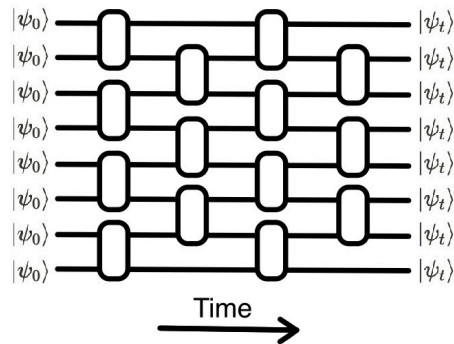


Quantum advantage within NISQ or LISQ regime

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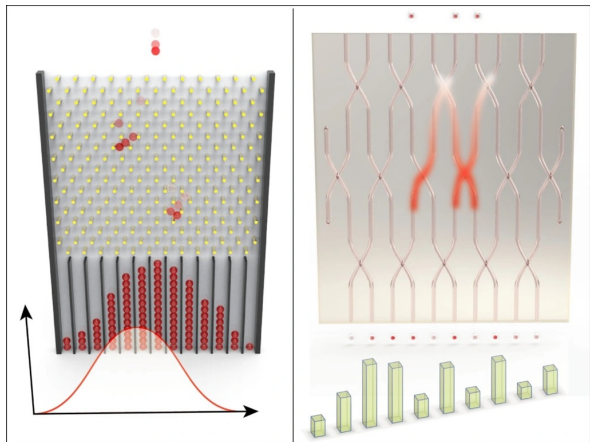


◆ Random circuit sampling

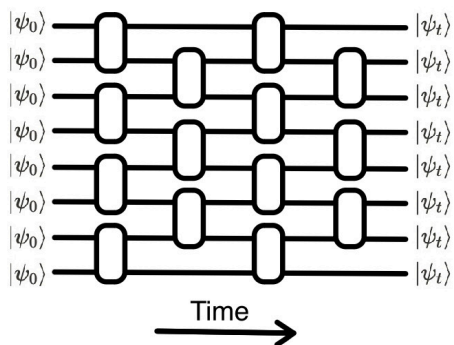


Quantum advantage within NISQ or LISQ regime

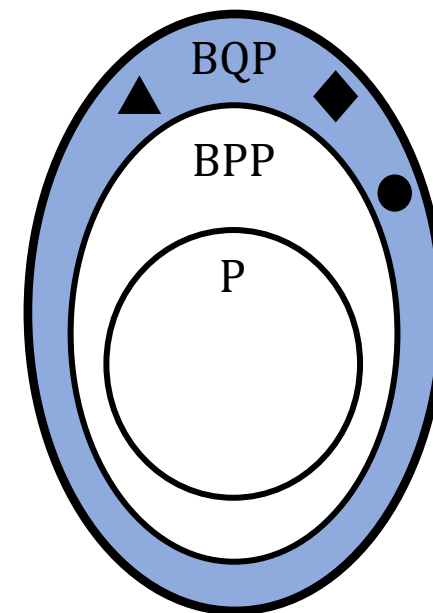
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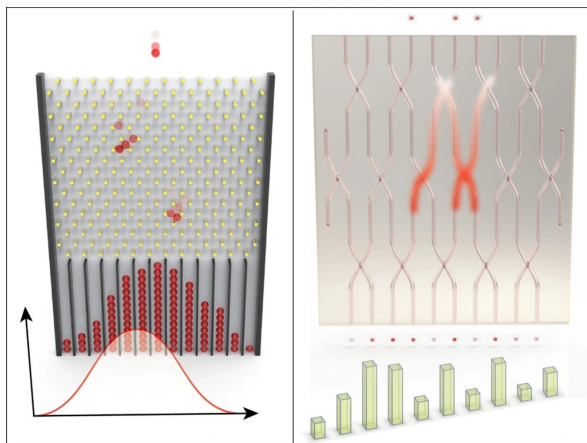


10^{25} years

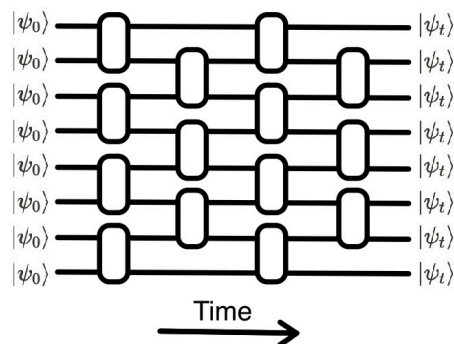


Quantum advantage within NISQ or LISQ regime

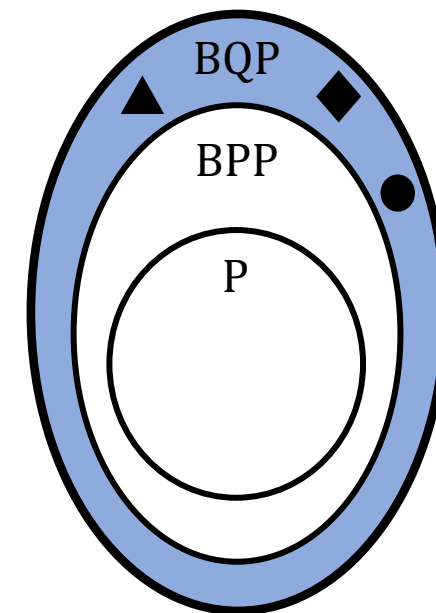
● Boson Sampling



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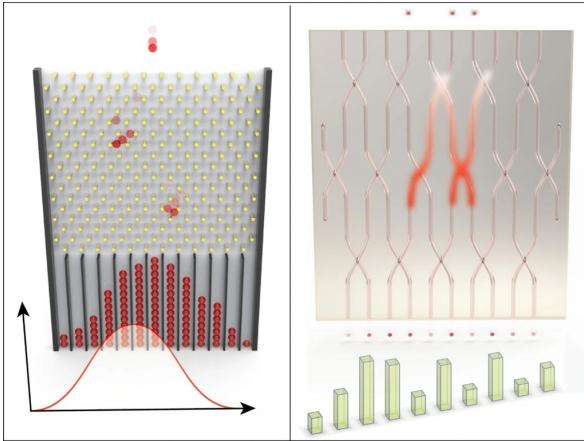
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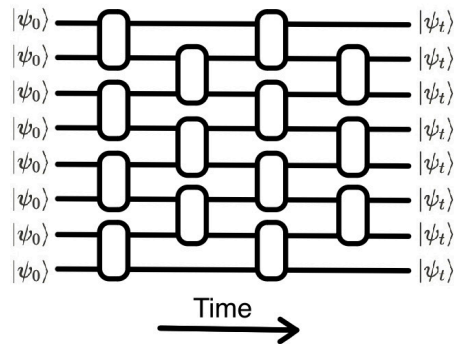
+ Shallow depth small quantum devices.

Quantum advantage within NISQ or LISQ regime

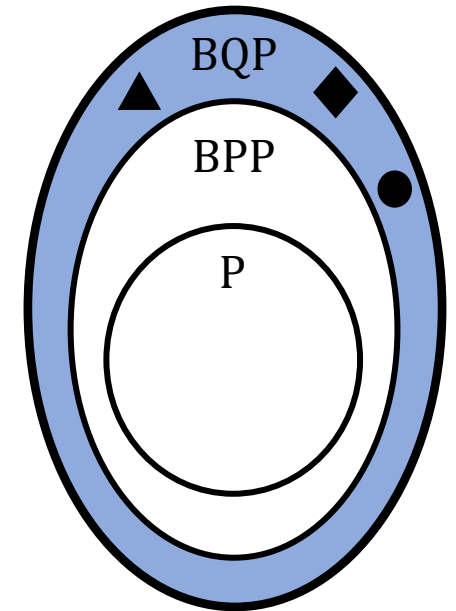
● Boson Sampling



◆ Random circuit sampling



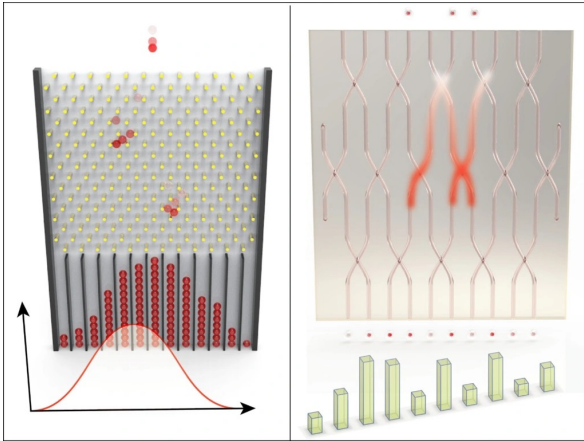
10^{25} years



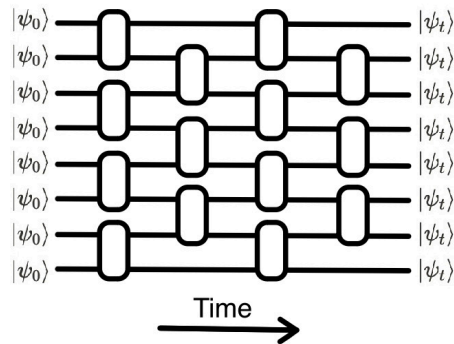
- + Shallow depth small quantum devices.
- Assuming the Polynomial Hierarchy does not collapse to the second level.

Quantum advantage within NISQ or LISQ regime

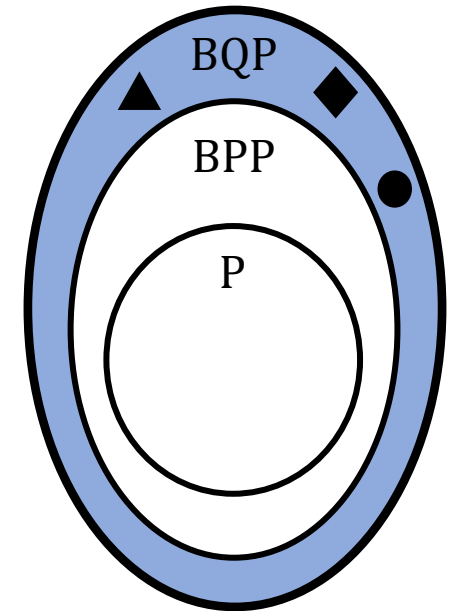
● Boson Sampling



◆ Random circuit sampling



10^{25} years



- + Shallow depth small quantum devices.
- Assuming the Polynomial Hierarchy does not collapse to the second level.
- It is hard to verify their correctness in the presence of noise.

Can we prove quantum advantage for any computational problem without relying on computational assumptions?

Unconditional quantum advantage within NISQ or LISQ regime

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Yes, if we fix time/depth to be constant!!!

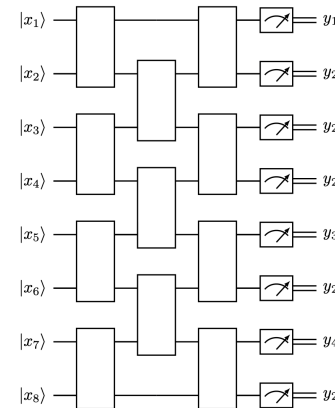
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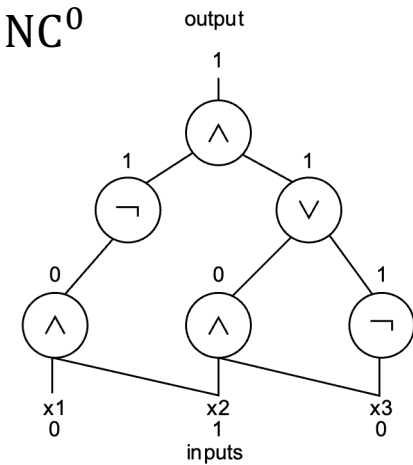
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QNC⁰



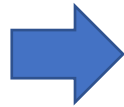
NC⁰



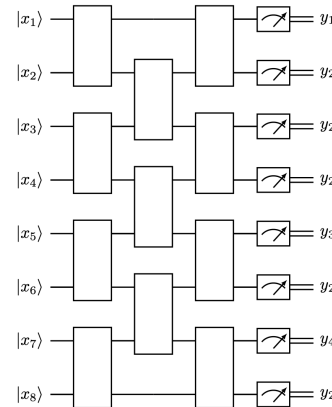
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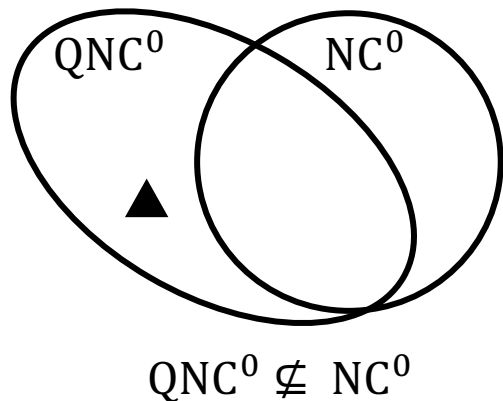
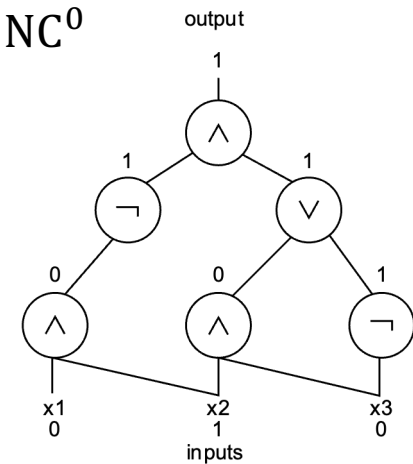


QNC^0



$>$

NC^0



Hidden Linear shift problem ▲

- + Shallow depth small quantum devices.
- + No assumptions.

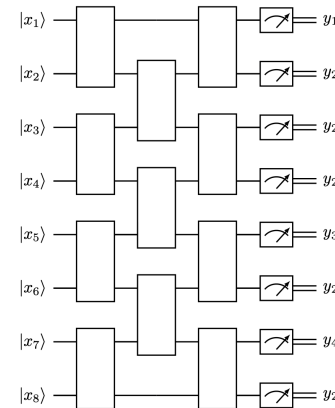
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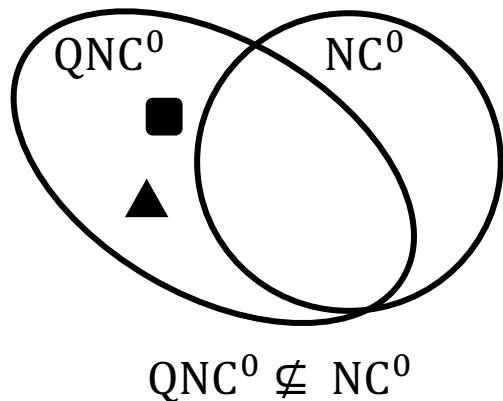
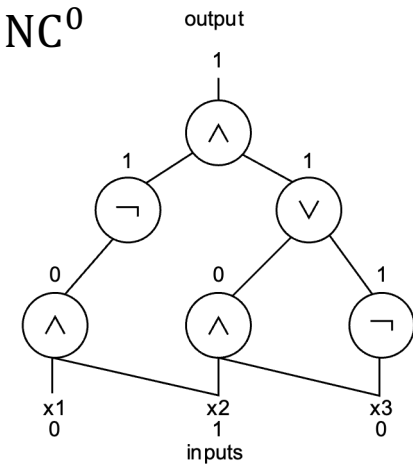
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1D Magic square problem ■

- ⋮
- + Error-robust.

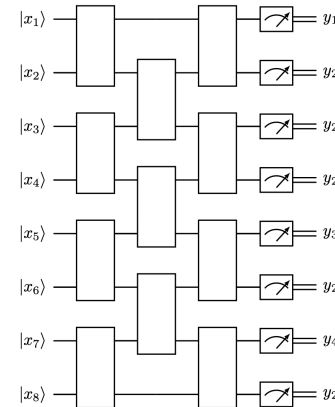
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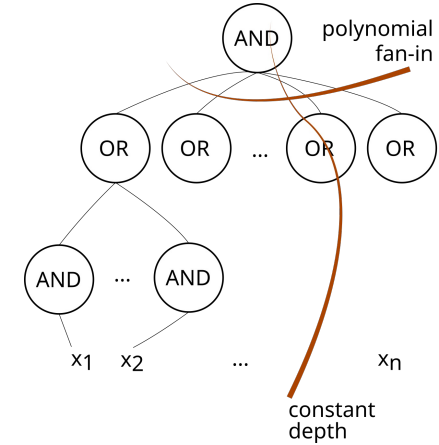


QNC^0



$>$

AC^0



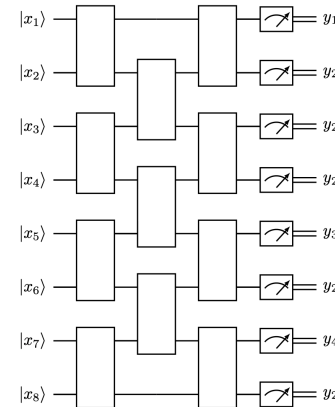
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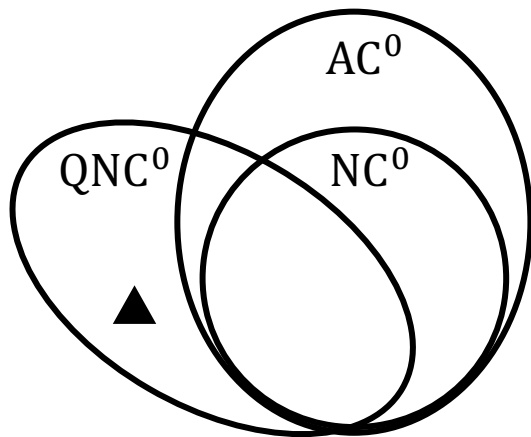
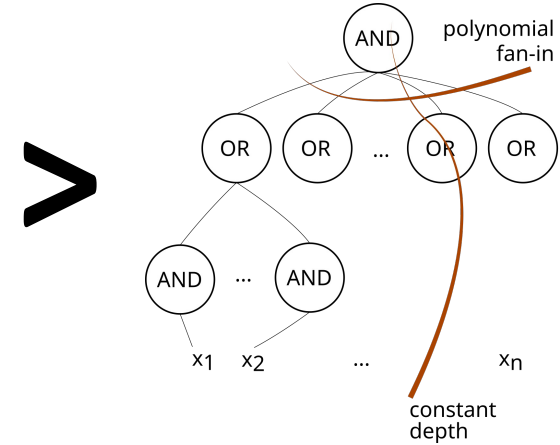
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▲ $\text{QNC}^0 \not\subseteq \text{AC}^0$

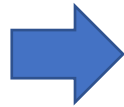
Parity halving problem ▲

+ Average-case hardness.

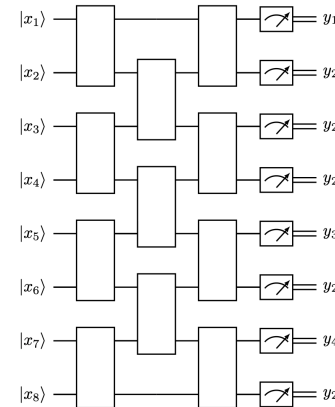
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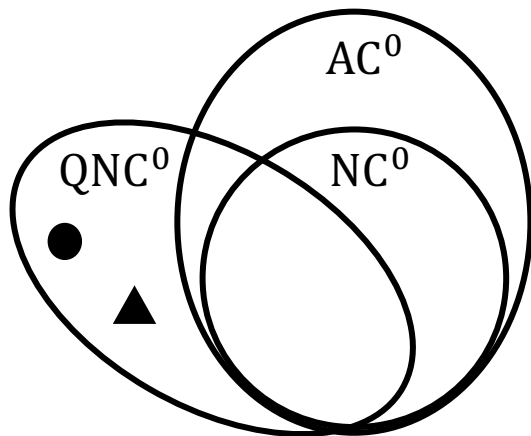
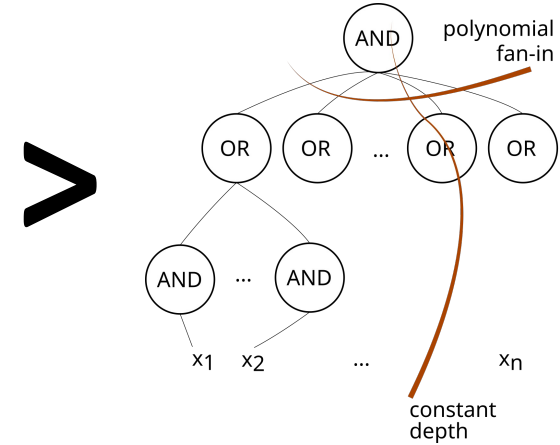
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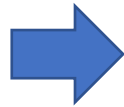
Single-qubit teleportation ●

+ Noise resilient.

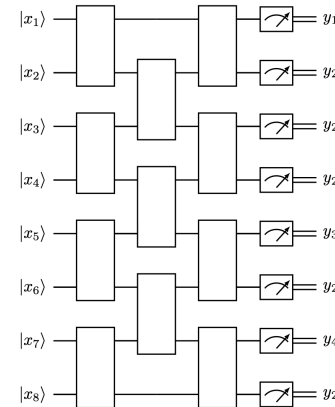
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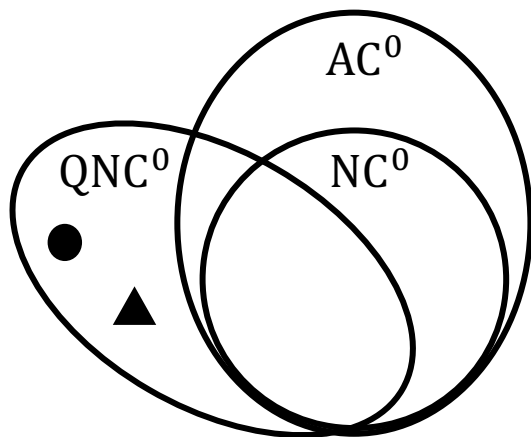
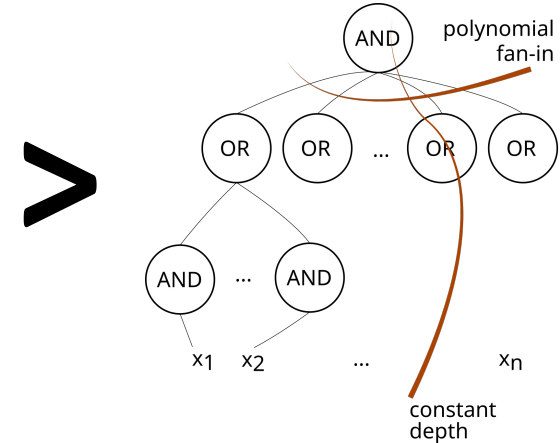
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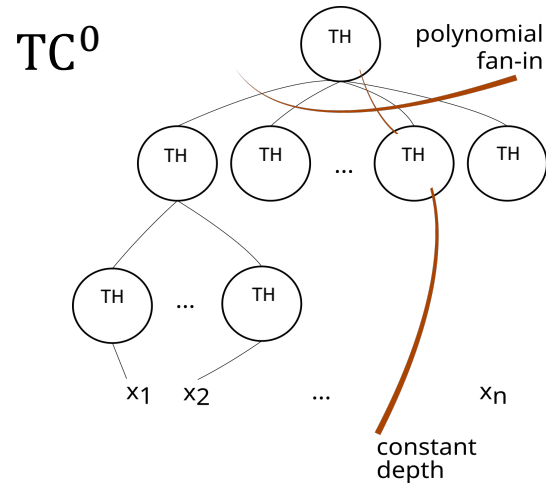
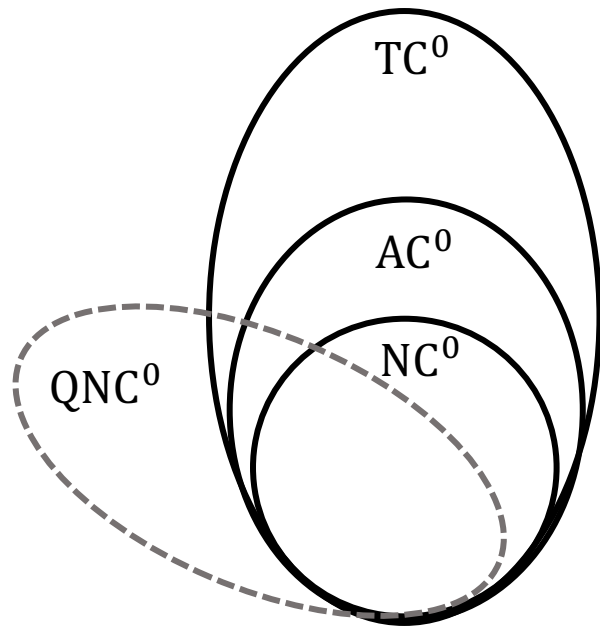
- + Average-case hardness.
- AC^0 circuit class is of rather small practical use.

Single-qubit teleportation ●

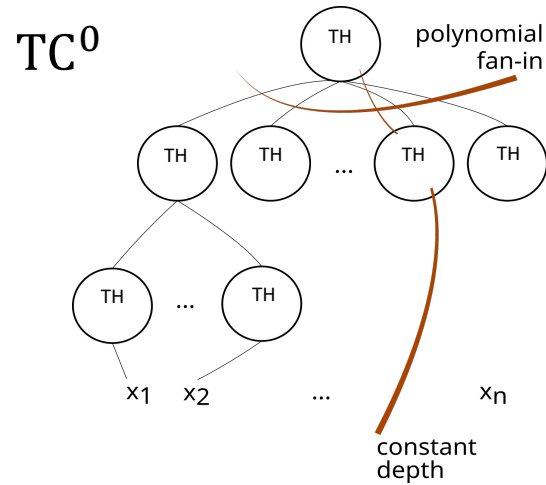
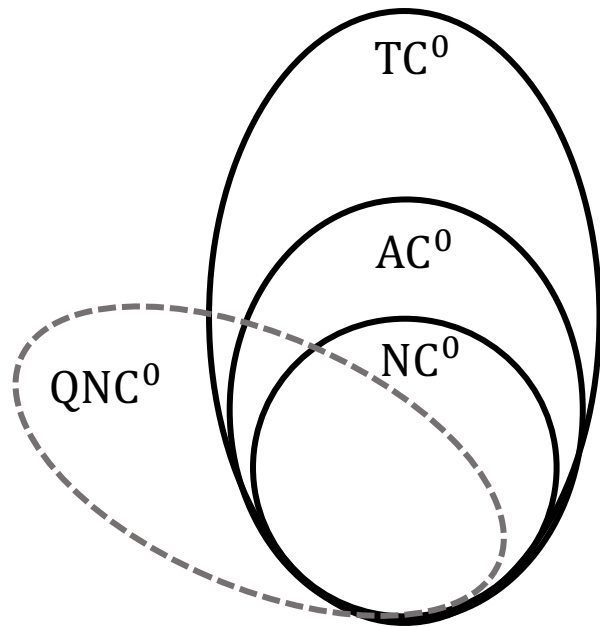
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How powerfull are constant-depth quantum circuits?

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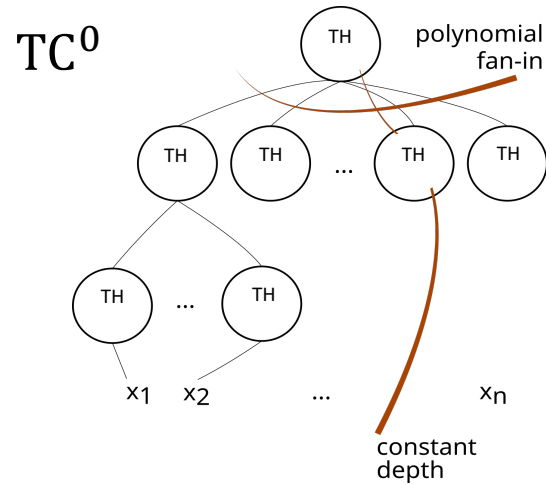
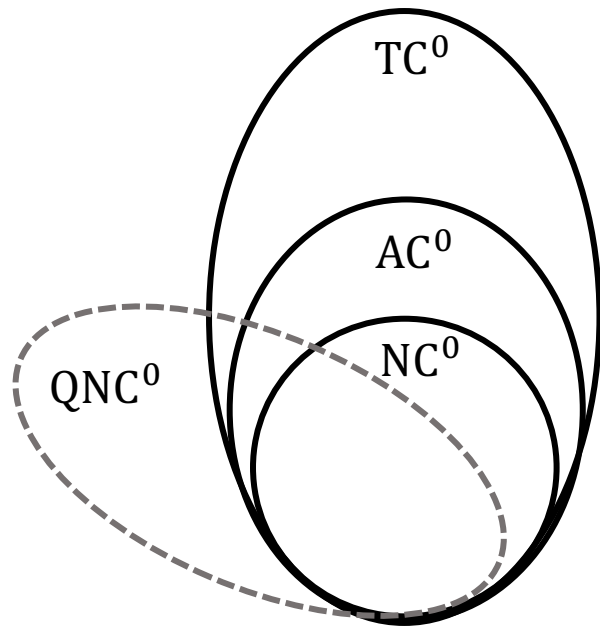


- TC^0 circuits are able a large set of algebraic problems

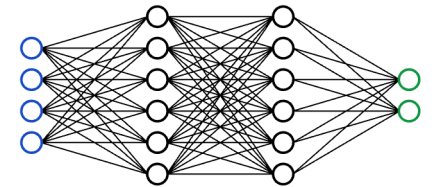
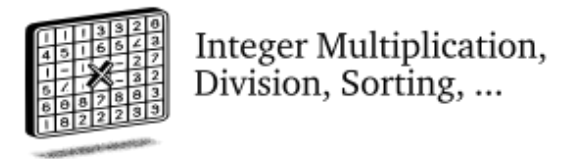


Integer Multiplication,
Division, Sorting, ...

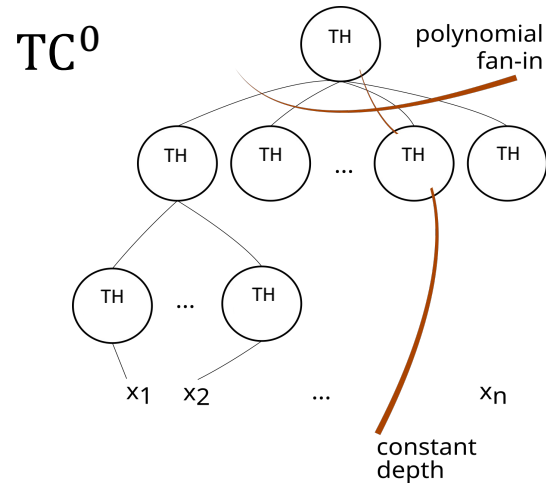
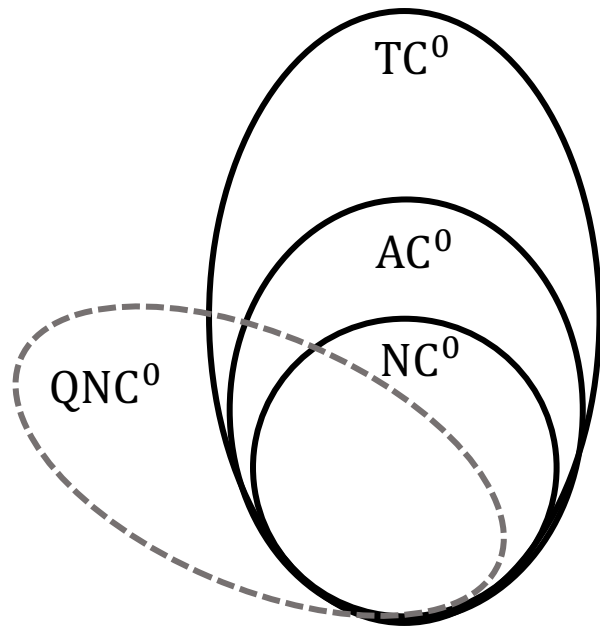
How powerfull are constant-depth quantum circuits?



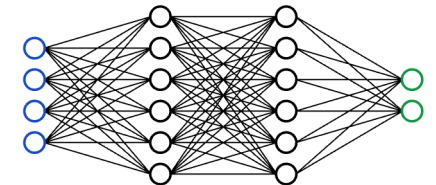
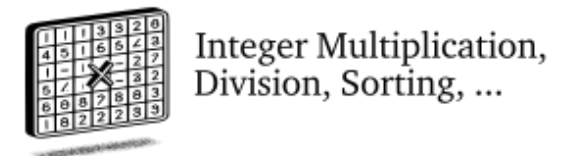
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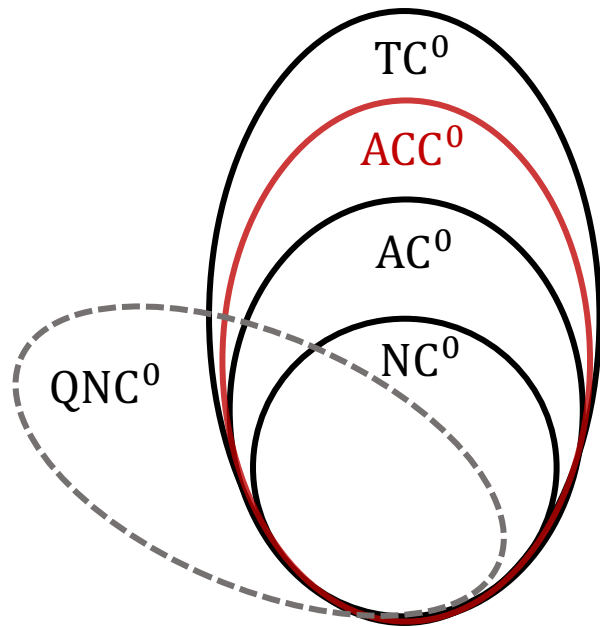


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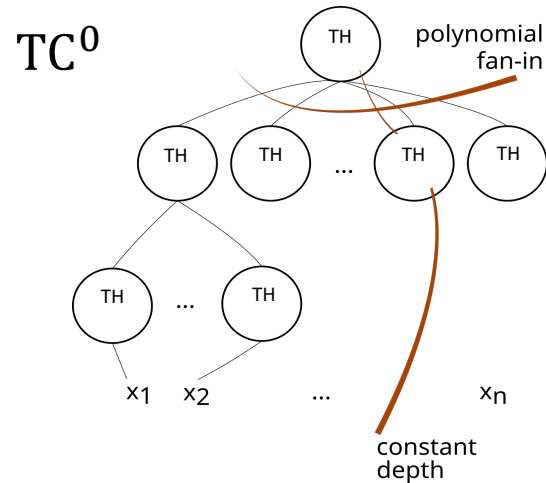


We do not have any strong lower bound techniques!

How powerfull are constant-depth quantum circuits?

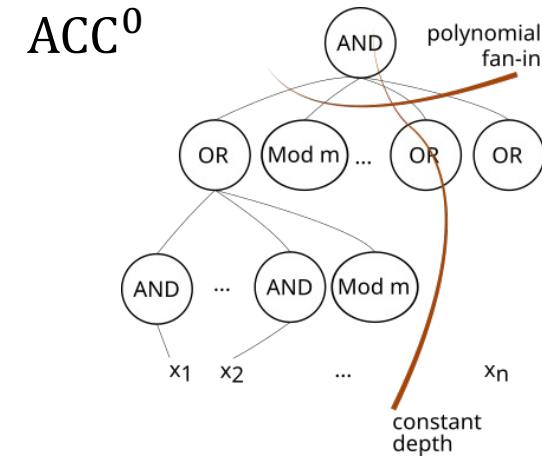


$E^{NP} \not\subseteq ACC^0$ [Williams16]

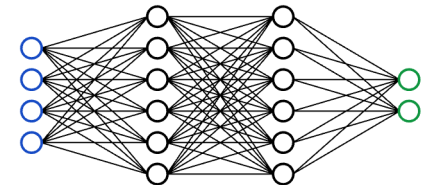


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- Alternatively can there exists a $\text{bPTF}^0[k]$ circuits parameterized by k can interpolate between AC^0 and beyond TC^0 .

How powerfull are constant-depth quantum circuits?

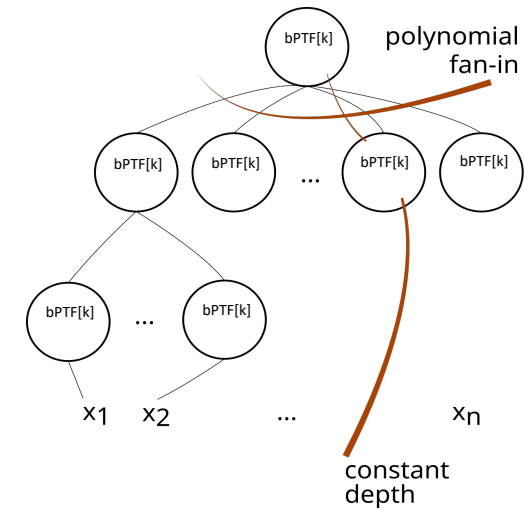
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Definition ($\text{bPTF}^0[k]$).

Constant-depth circuits with unbounded fan-in $\text{bPTF}[k]$ gates defined as follows,

$$f_{\text{or}}(x) = \begin{cases} P(x), & \sum_{i=0}^n x_i \leq k \\ 1, & \sum_{i=0}^n x_i > k \end{cases}$$

with $P: F_2^n \rightarrow F_2$ a polynomial over $F_2 = \{0,1\}$.



How powerfull are constant-depth quantum circuits?

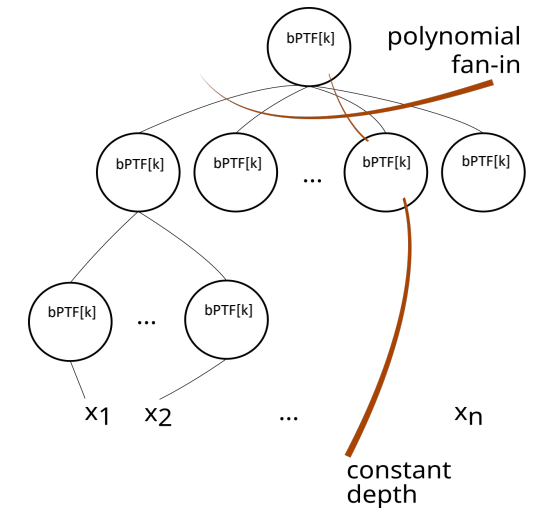
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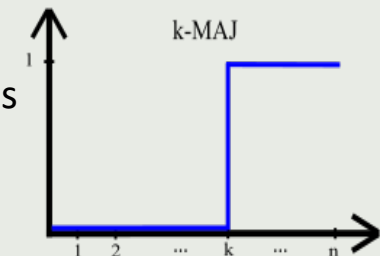
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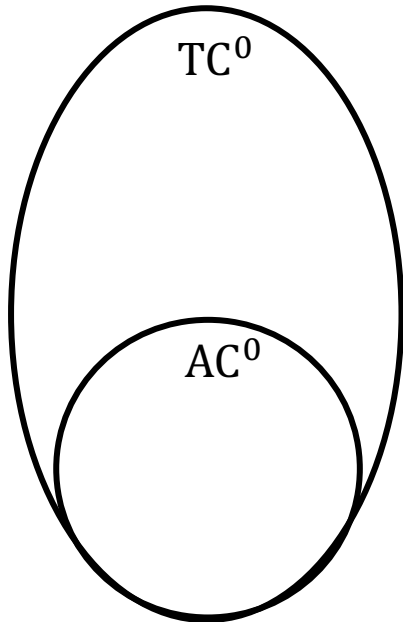
Example

$\text{bPTF}^0[k]$ includes k -parameterized TH gates



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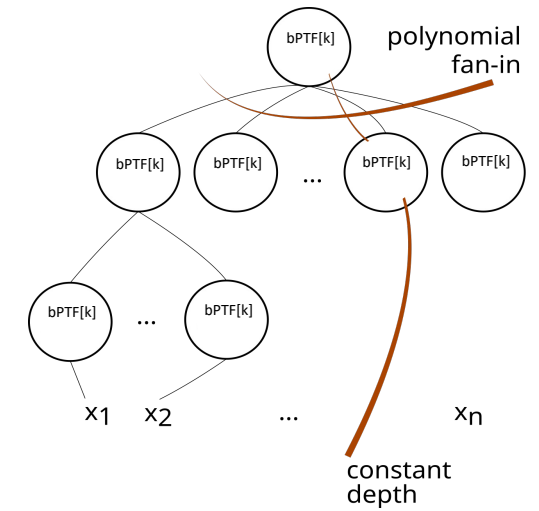


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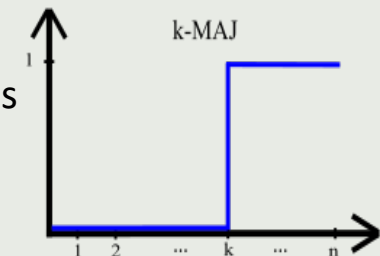
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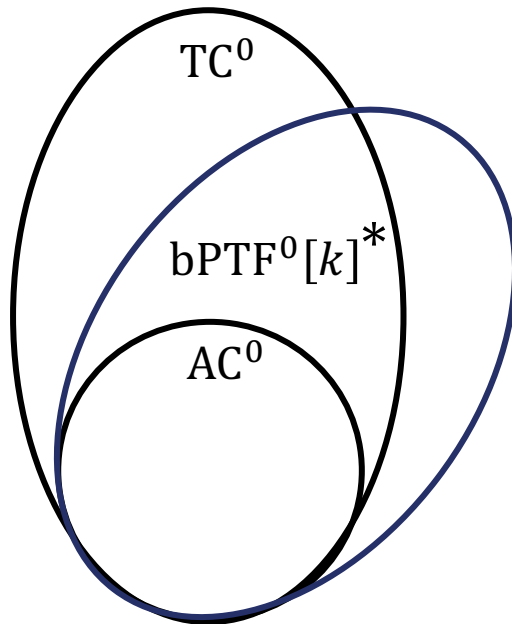
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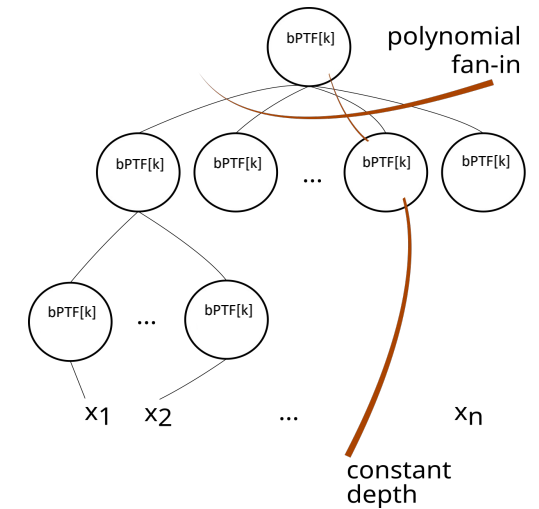


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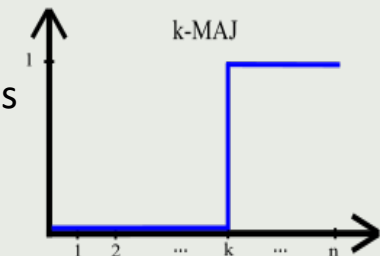
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$$\begin{aligned} \text{AC}^0 &= \text{bPTF}^0[1] \\ \text{AC}^0 &\subsetneq \text{bPTF}^0[k], \quad k = w(\log n) \end{aligned}$$

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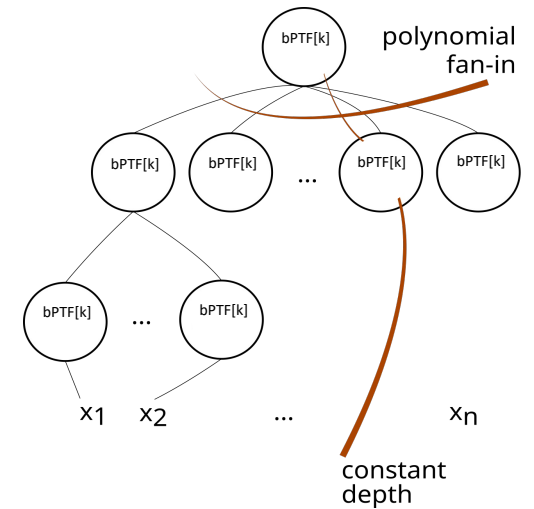
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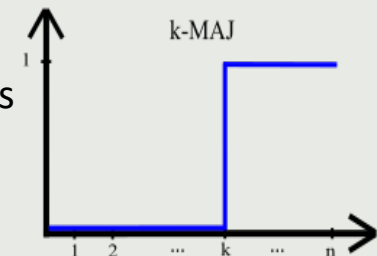
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Example

$\text{bPTF}^0[k]$ includes k -parameterized TH gates



All = $\text{bPTF}^0[n]$

TC^0

$\text{bPTF}^0[k]^*$

AC^0

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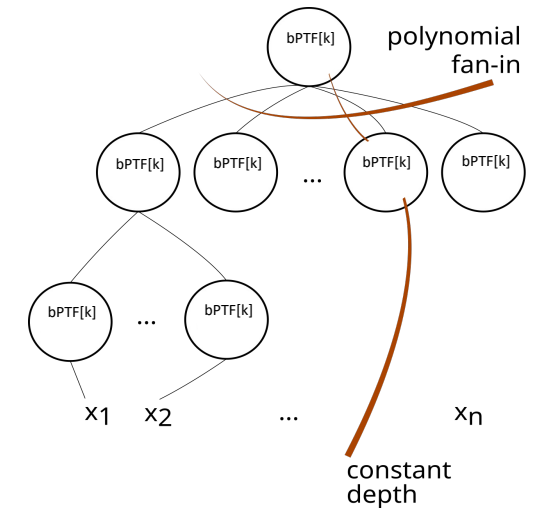
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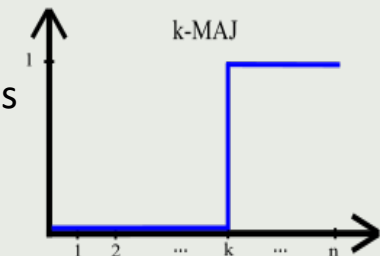
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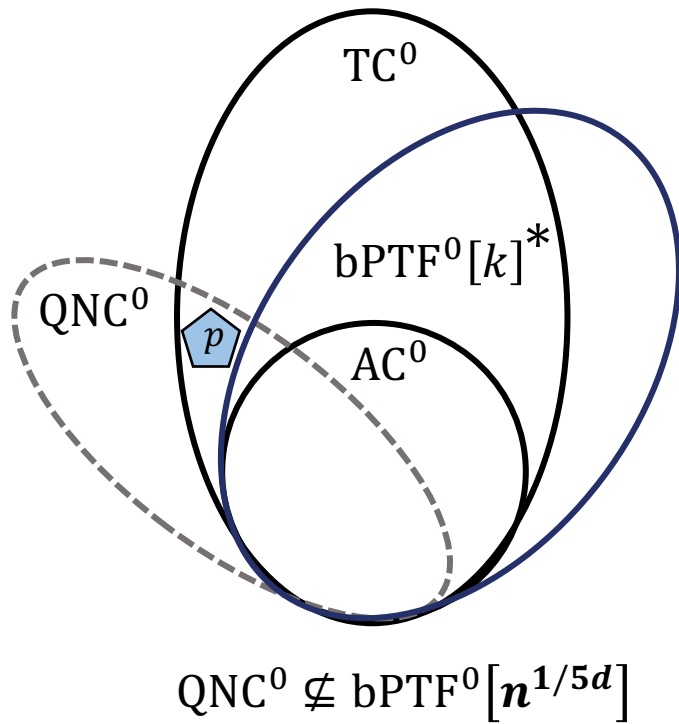
AC^0

QNC^0

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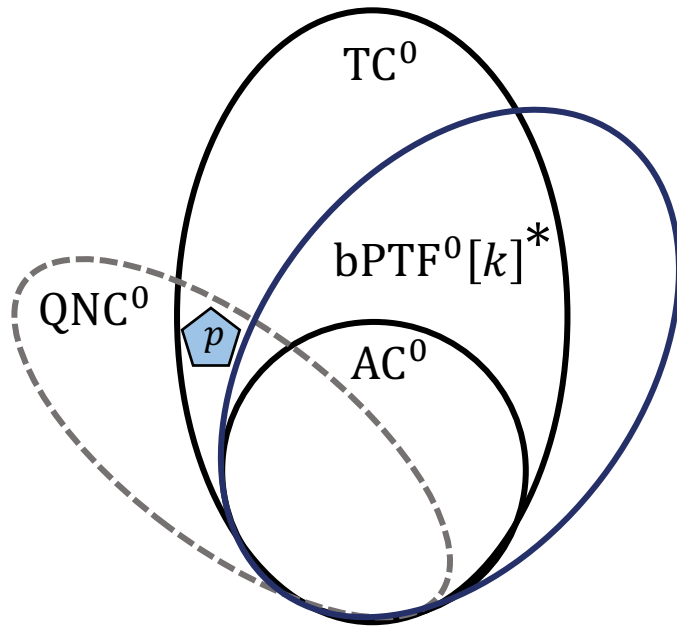
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New Quantum-Classical unconditional separations

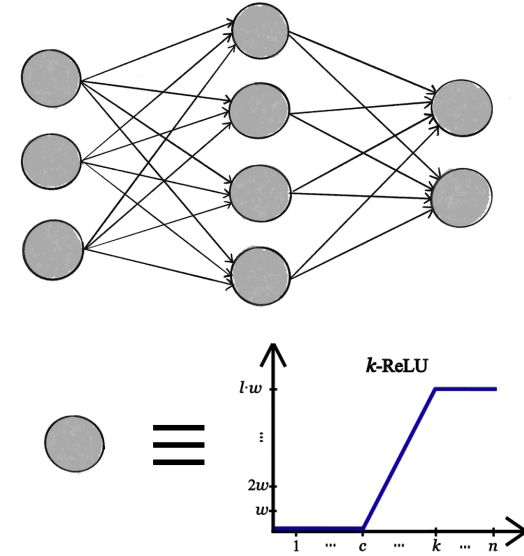
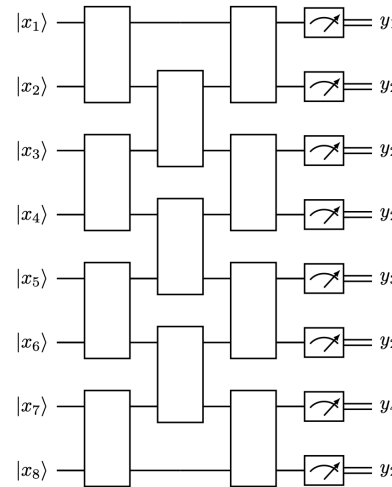


New Quantum-Classical unconditional separations

- We proved that **parallel quantum computation unconditionally outperforms** larger classical circuits classes, including biased Neural Networks.

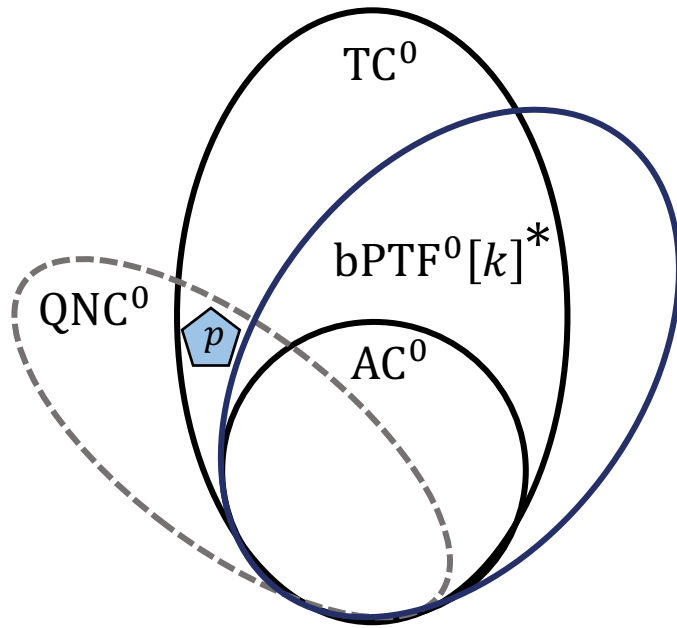


$$QNC^0 \not\subseteq bPTF^0[n^{1/5d}]$$

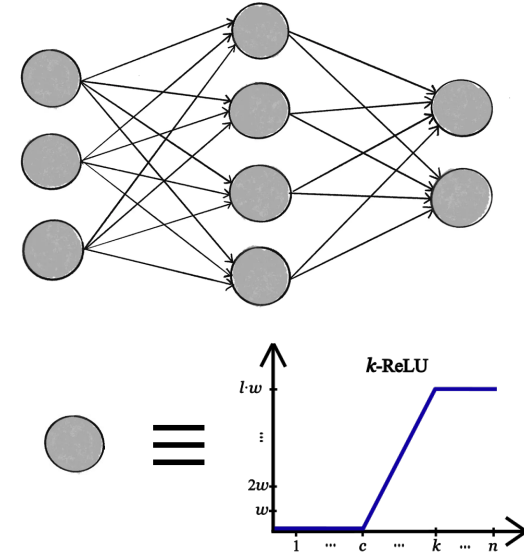
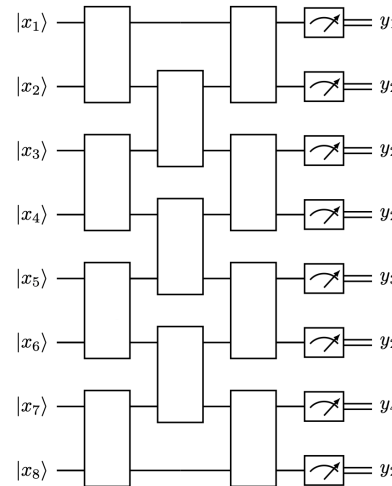


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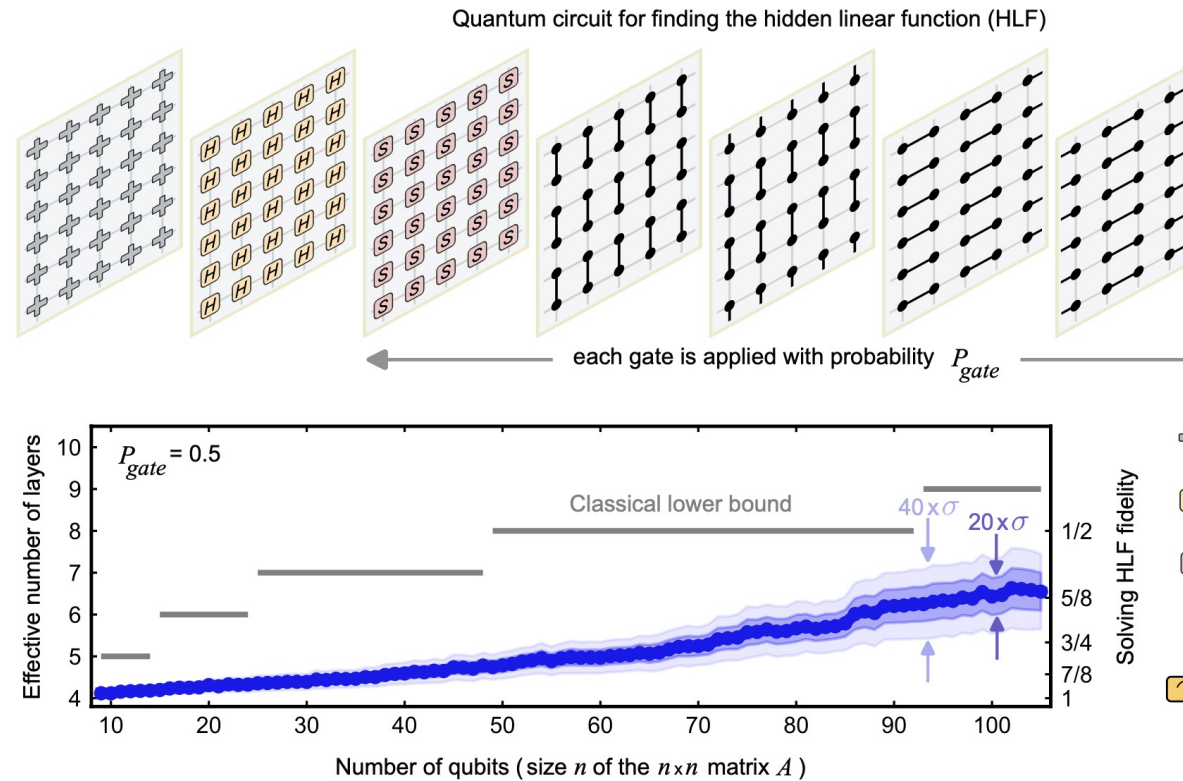
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- This advantage is **robust to noise** and operates on **qudit quantum devices**.

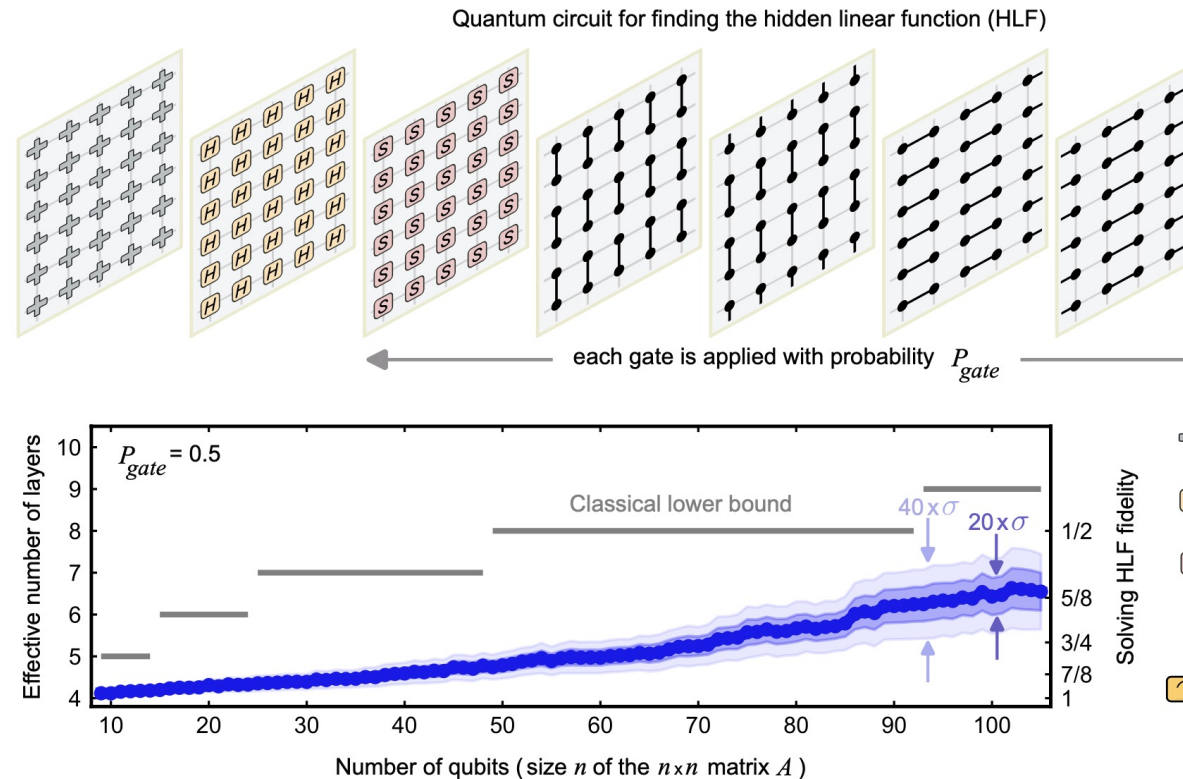
Solving the Hidden Linear shift problem on quantum devices

This week, the Google team released a paper on the realization of HLF on their Willow chip.



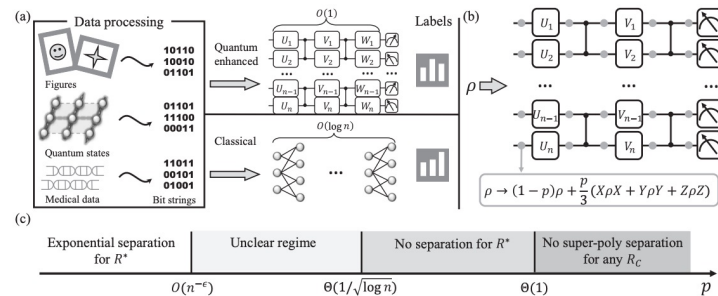
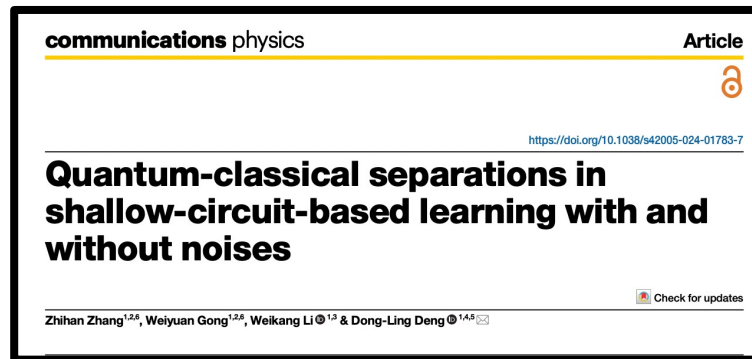
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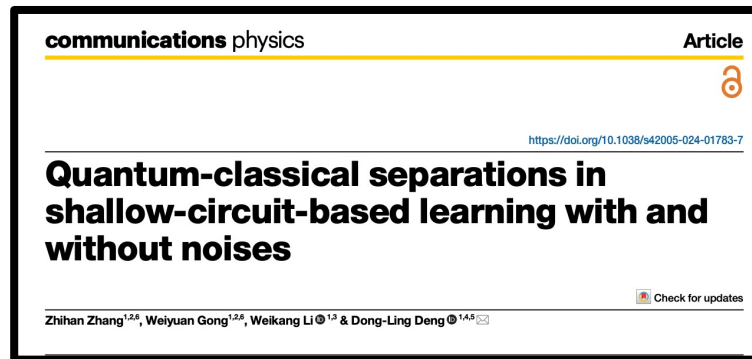


*If they could scale the number of qubits, they could realize our new quantum advantages.

New Practically Motivated Quantum-Classical unconditional separations



New Practically Motivated Quantum-Classical unconditional separations



An unconditional distribution learning advantage with shallow quantum circuits

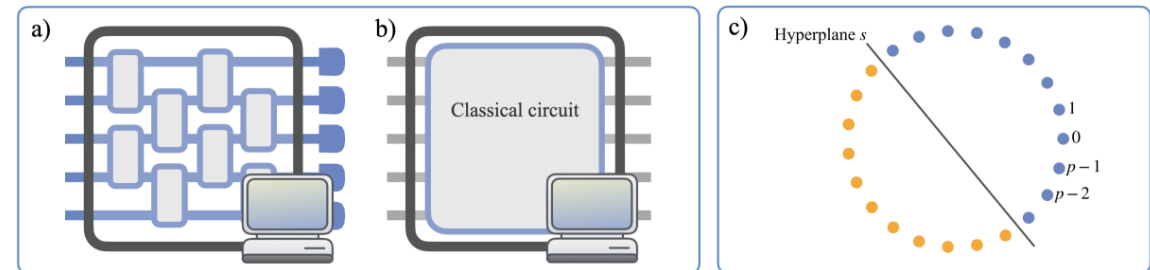
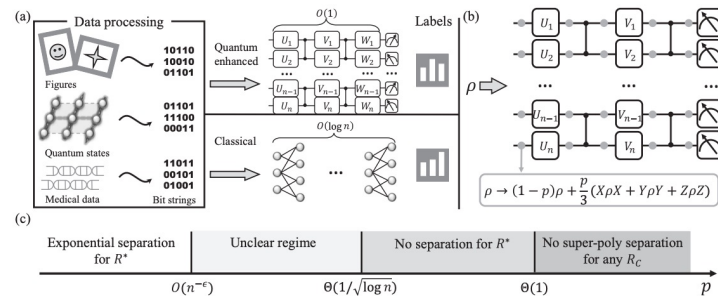
N. Pirnay,¹ S. Jerbi,² J.-P. Seifert,¹ and J. Eisert^{2,3,4}

¹Electrical Engineering and Computer Science, Technische Universität Berlin, Berlin, 10587, Germany

²Dahlem Center for Complex Quantum Systems, Freie Universität Berlin, 14195 Berlin, Germany

³Fraunhofer Heinrich Hertz Institute, 10587 Berlin, Germany

⁴Helmholtz-Zentrum Berlin für Materialien und Energie, 14109 Berlin, Germany



The Quantum Benchmarking problem

Quantum Benchmarks

Quantum Benchmarks which test quantum devices should be:

- Algorithm agnostic

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Quantum Benchmarks

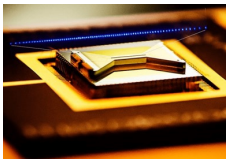
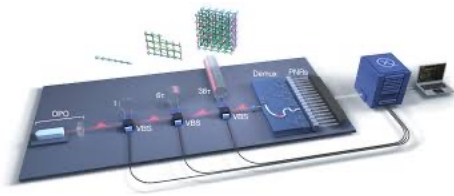
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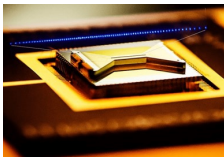
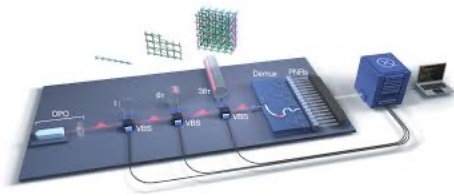
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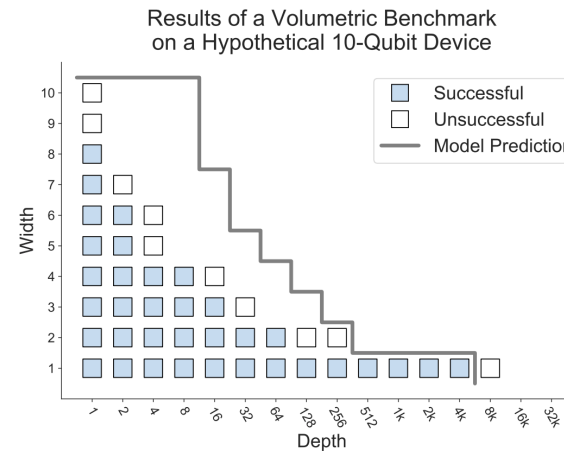
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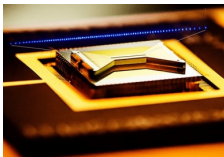
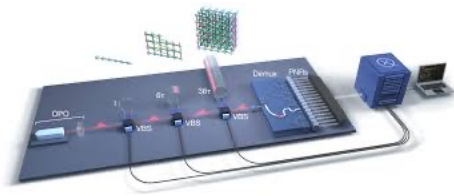
Quantum Volume



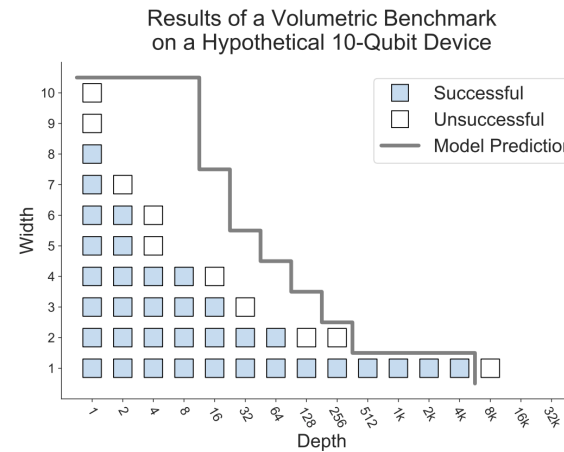
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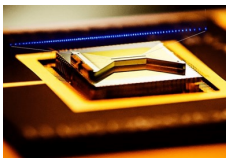
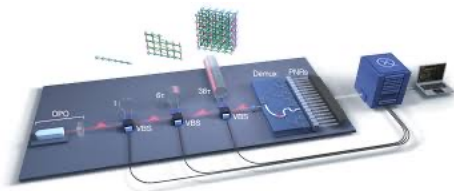
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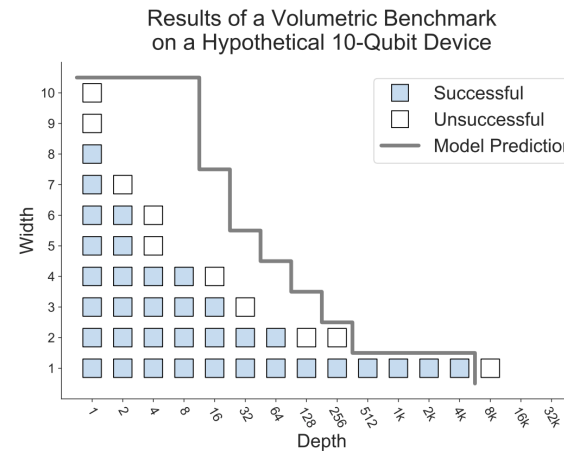
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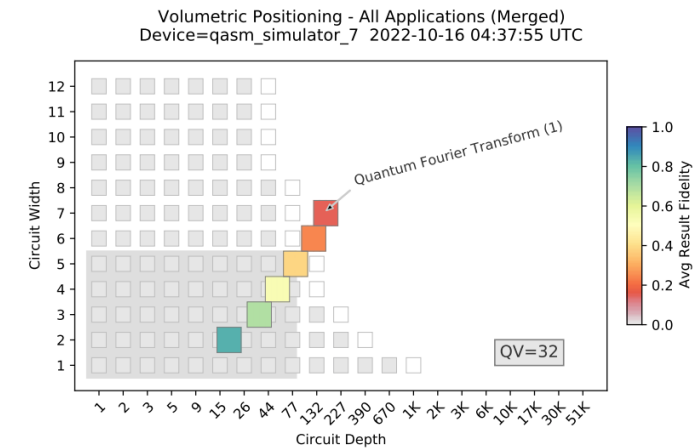


Quantum Volume



Misses to be predictive and scalable

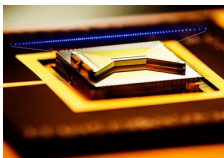
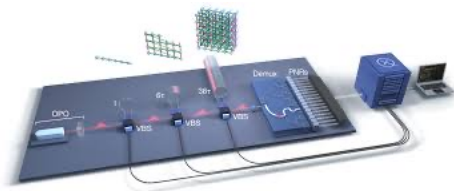
Application-oriented benchmarks



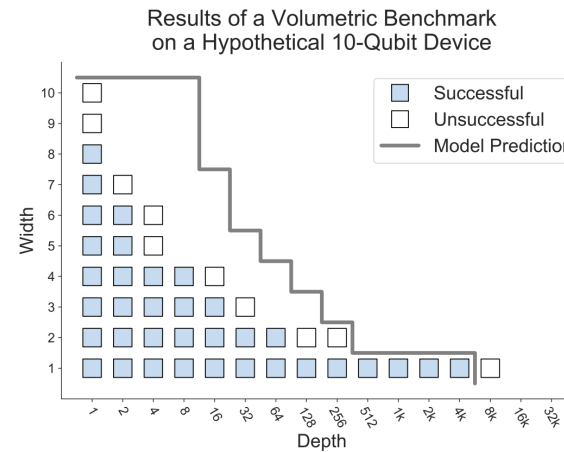
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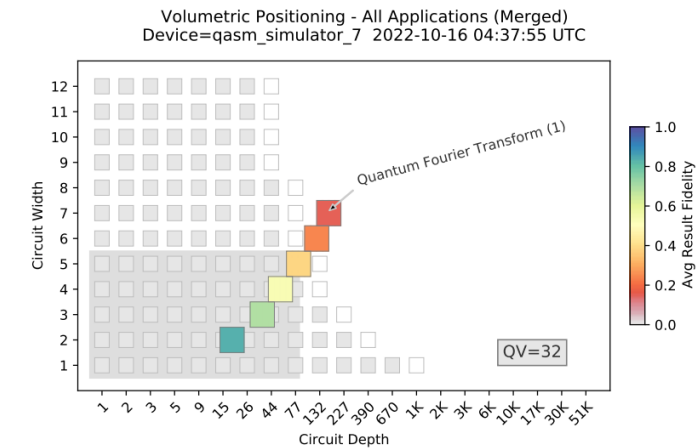


Quantum Volume



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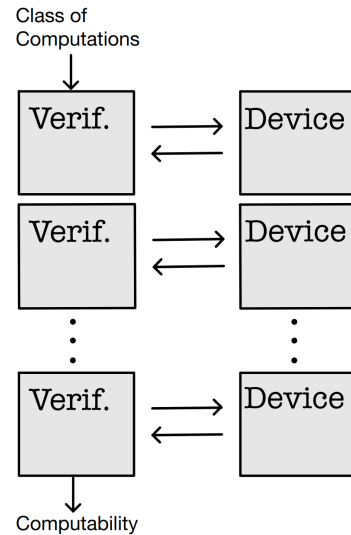
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Misses to algorithm agnostic and predictive

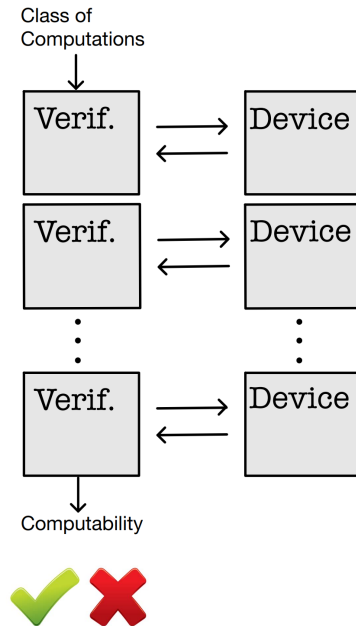
Certifying the accuracy of quantum computers

- Developed a benchmark to formally certify quantum devices for executing broad classes of circuits.



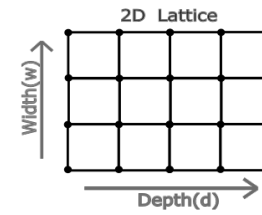
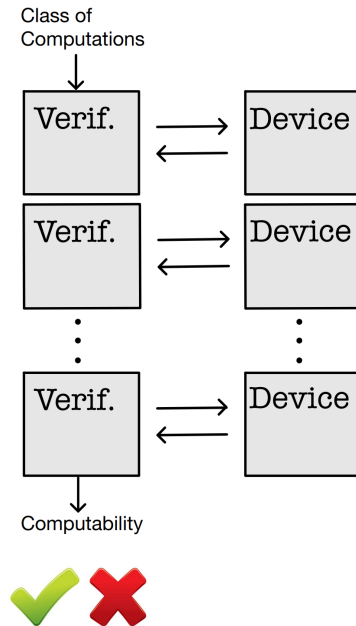
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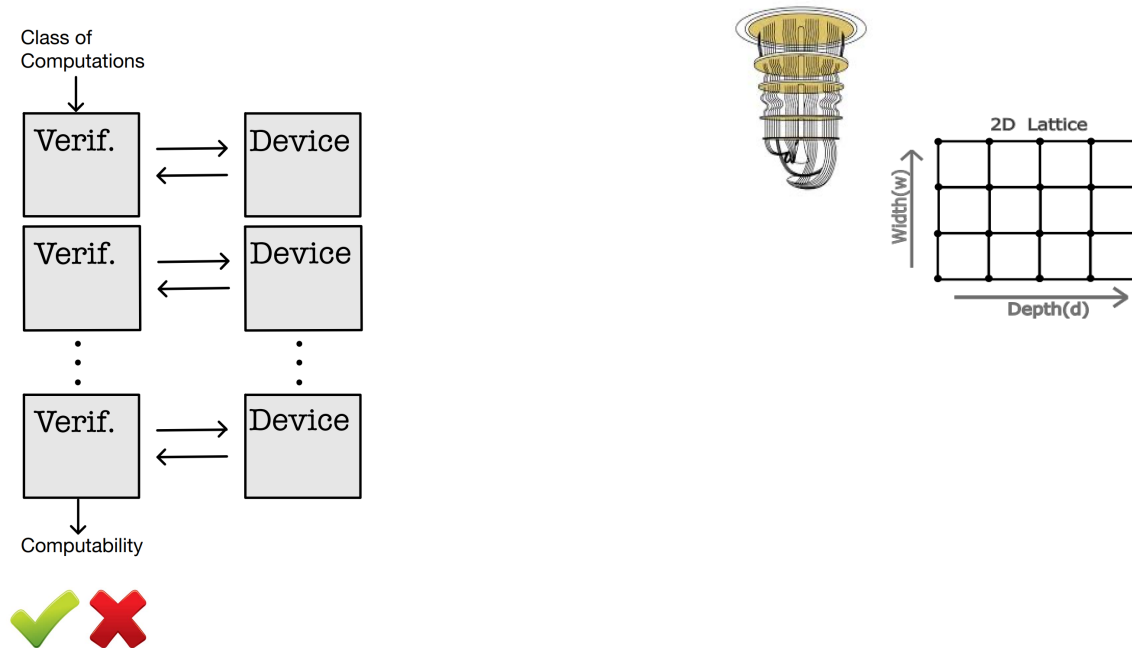
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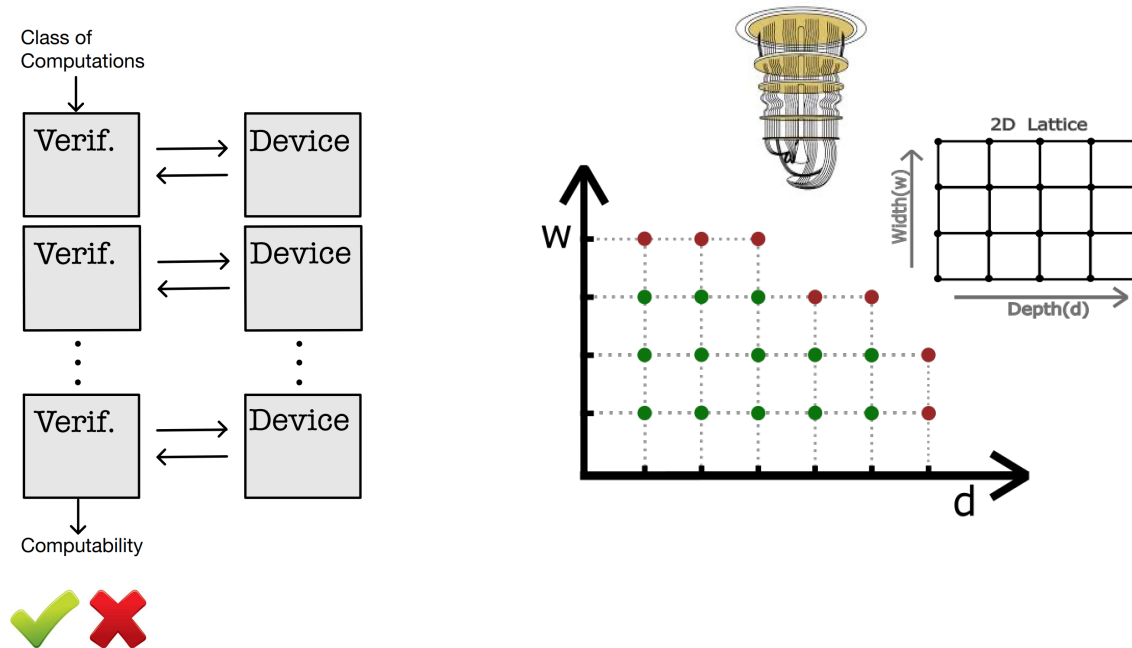
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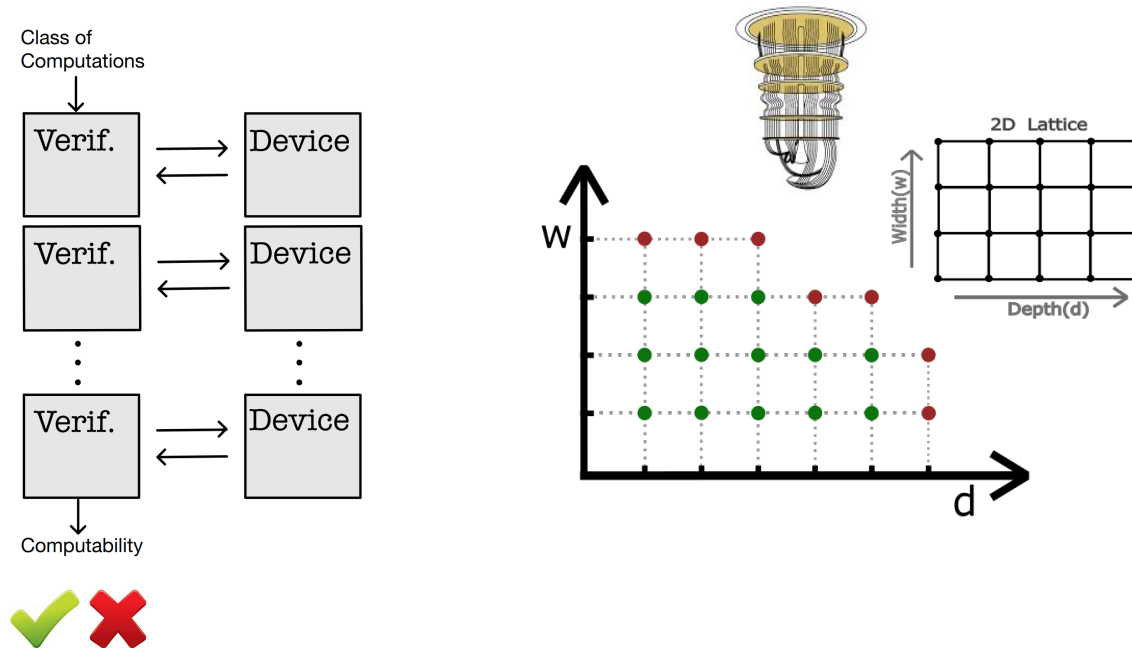
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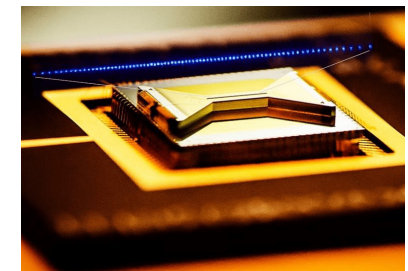


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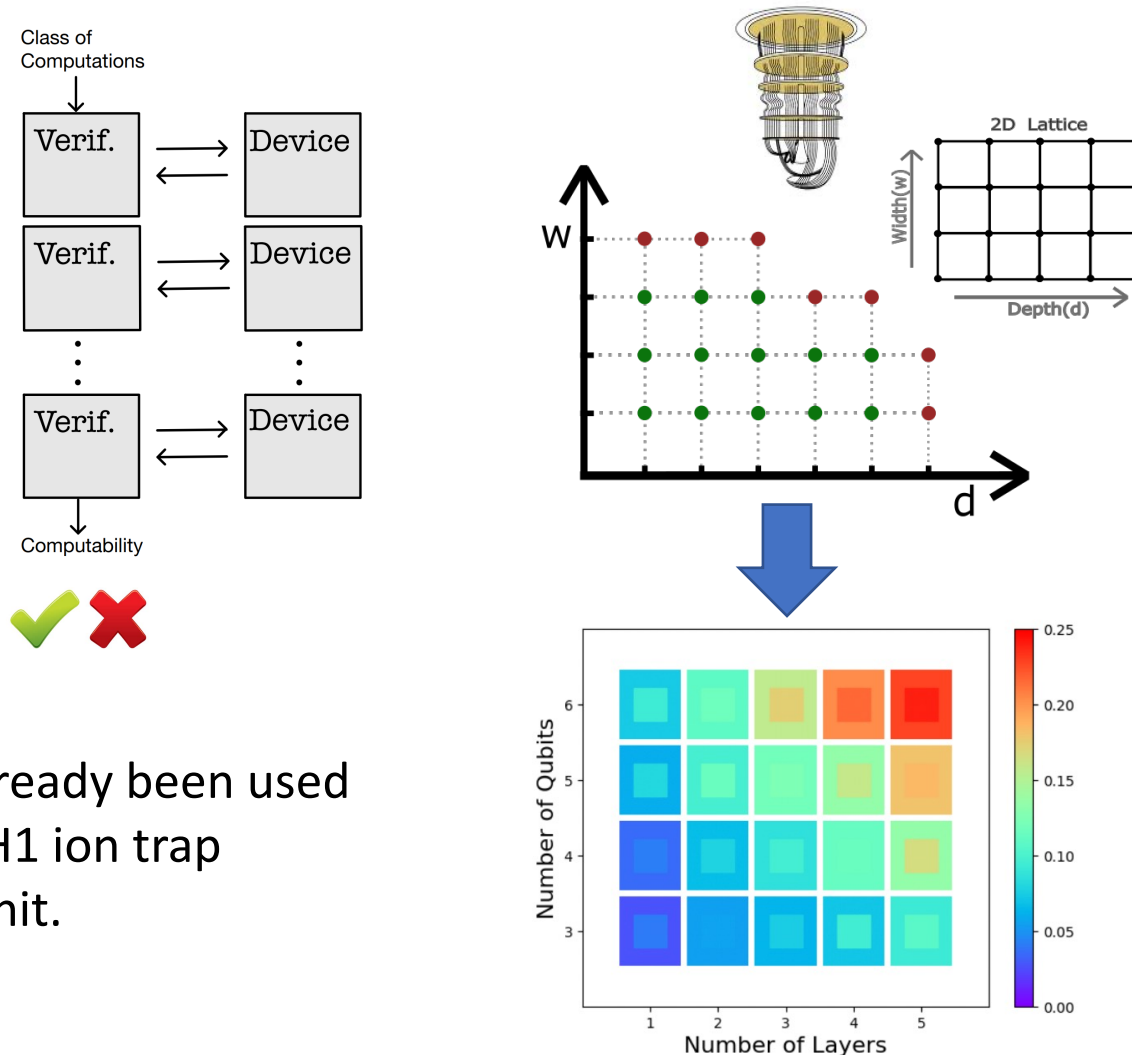


- This benchmark has already been used to test Quantinuum's H1 ion trap quantum processing unit.

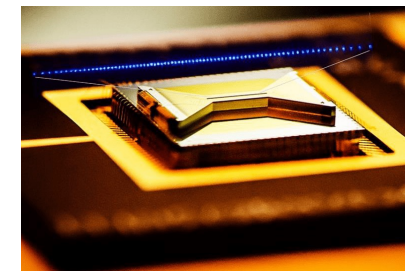


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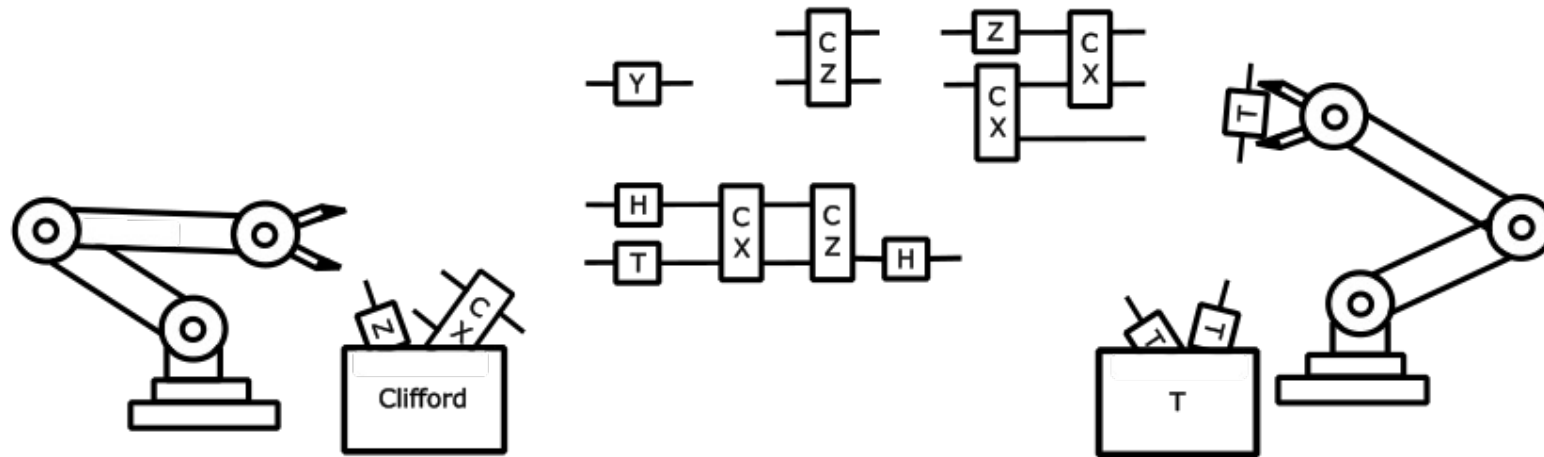
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The circuit synthesis problem



The quest for quantum advantage

This work and its results gave origin to,



My PhD thesis

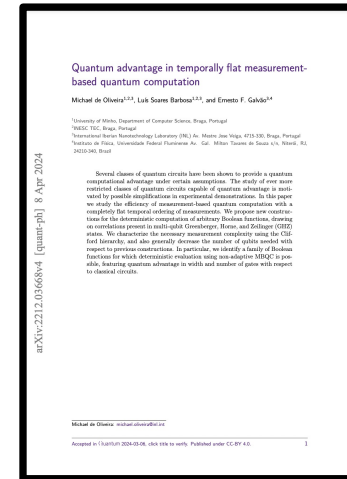
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New synthesis techniques for
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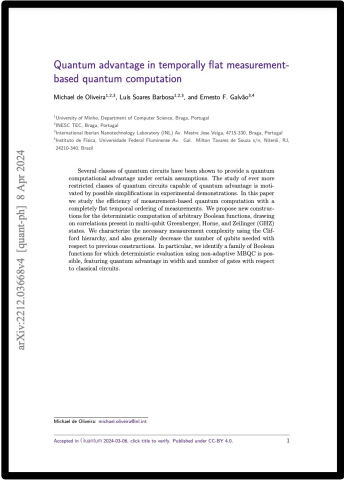
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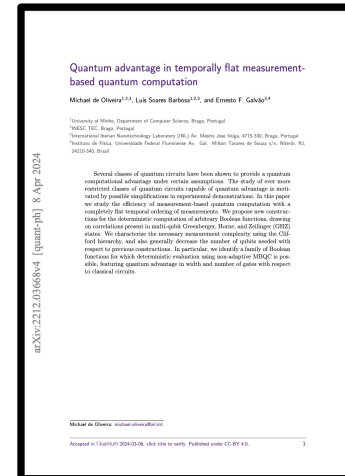


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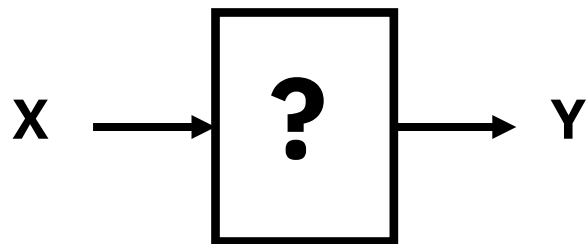
Unconditional quantum advantages with shallow depth circuits



Thank you!

@IBM - A New Quantum Learning Algorithm

Understand how hard it is to learn an unknown process:



$$? : \{0, 1\}^n \rightarrow \{0, 1\}$$

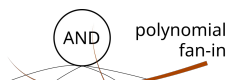
Focus:

- Sample complexity
- Accuracy of the learned hypothesis
- Run-time

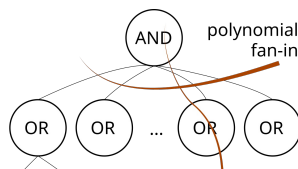
Learning arbitrary Boolean functions is an extremely hard problem!

?

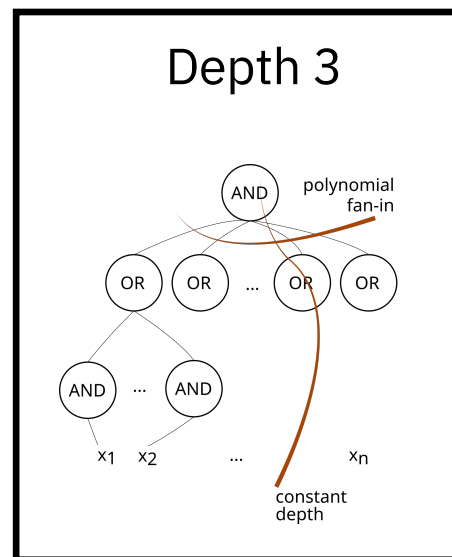
Depth 1



Depth 2



Depth 3



We present a novel, state-of-the-art quantum learning algorithm for PAC learning depth-3 circuits.

