

# Lecture 9:

## Estimating eigenvalues: An application of QFT

Luís Soares Barbosa



Universidade do Minho



**Mestrado em Engenharia Física**

Universidade do Minho, 2025-26

## The problem: Eigenvalue estimation

### The eigenvalue estimation problem

Let  $(|\psi\rangle, e^{i2\pi\phi})$ , with  $0 \leq \phi < 1$ , be an eigenvector, eigenvalue pair for a unitary  $U$ . Determine  $\phi$ .

Note that eigenvalues of unitary operators are always of this form. Why?

## The strategy

- Use a **controlled version of  $U$**  to prepare a state from which  $\phi$  can be found.
- Then, resort to the **inverse of the QFT** to find it.
- The **accuracy** of the estimation increases with the **number of qubits available** for the recovery state

Thus, the problem reduces to the already discussed

**phase estimation problem**

## The general case

A **multi-controlled** version of  $U$  is required:

$$\left[ \begin{array}{c} \text{---}^n\text{---} \\ | \\ \text{---}^m\text{---} \end{array} \right] U \left[ \begin{array}{c} \text{---}^n\text{---} \\ | \\ \text{---}^m\text{---} \end{array} \right] = |x\rangle |y\rangle \mapsto |x\rangle U^\times |y\rangle$$

↓  
decimal representation of x

Intuitively it applies  $U$  to  $|y\rangle$  a number of times equal to  $x$

## Examples

$$|10\rangle |y\rangle \mapsto |10\rangle (UU|y\rangle) \text{ and } |00\rangle |y\rangle \mapsto |00\rangle |y\rangle$$

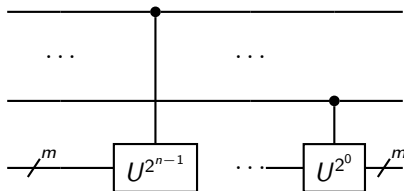
Note that  $|\psi\rangle$  is also an eigenvector of  $U^x$ , with eigenvalue  $e^{i2\pi x\phi}$ , for any integer  $x$ .

## Multi-controlled operations

Recall that a binary number  $x_1 \dots x_n$  corresponds to the natural number

$$2^{n-1}x_1 + \dots + 2^0x_n$$

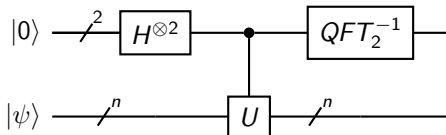
We use this to build the previous multi-controlled operation in terms of  $n$  'simply'-controlled rotations  $U^{2^i}$



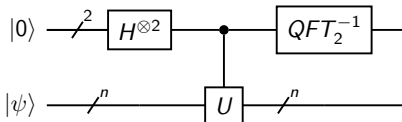
## An Example

Take a unitary  $U$  with an eigenvector  $|\psi\rangle$  whose eigenvalue is  $e^{i2\pi\phi}$   
 $\phi$  is equal to one of the following values  $\{0 \cdot \frac{1}{4}, 1 \cdot \frac{1}{4}, 2 \cdot \frac{1}{4}, 3 \cdot \frac{1}{4}\}$

The following circuit discovers  $\phi$



## Another Example



$$|0\rangle |0\rangle$$

$$H^{\otimes 2} \mapsto \frac{1}{\sqrt{2^2}} (|00\rangle + |01\rangle + |10\rangle + |11\rangle)$$

$$\text{ctrl. } U \mapsto \frac{1}{\sqrt{2^2}} (|00\rangle + e^{i2\pi\phi} |01\rangle + e^{i2\pi\phi \cdot 2} |10\rangle + e^{i2\pi\phi \cdot 3} |11\rangle)$$

$$= \frac{1}{\sqrt{2^2}} (|00\rangle + e^{i2\pi x \cdot \frac{1}{4}} |01\rangle + e^{i2\pi x \cdot \frac{1}{4} \cdot 2} |10\rangle + e^{i2\pi x \cdot \frac{1}{4} \cdot 3} |11\rangle)$$

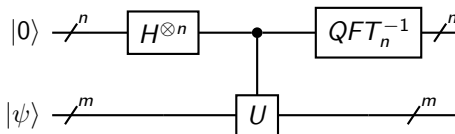
$$= \frac{1}{\sqrt{2^2}} (|00\rangle + \omega_2^x |01\rangle + \omega_2^{x \cdot 2} |10\rangle + \omega_2^{x \cdot 3} |11\rangle)$$

$$QFT_2^{-1} \mapsto |x\rangle$$

## Yet Another Example

Take a unitary  $U$  with eigenvector  $|\psi\rangle$  whose eigenvalue is  $e^{i2\pi\phi}$   
st  $\phi \in \{0 \cdot \frac{1}{2^n}, \dots, 2^n - 1 \cdot \frac{1}{2^n}\}$

The following circuit returns  $x$  such that  $\phi = x \cdot \frac{1}{2^n}$



### Exercise

Prove that indeed the circuit returns  $x$  such that  $\phi = x \cdot \frac{1}{2^n}$

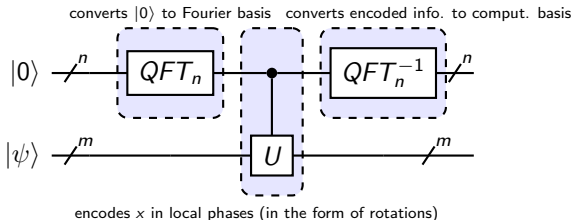


## Yet Another Example

### Exercise

Show that  $QFT_n |0\rangle = H^{\otimes n} |0\rangle$ .

Note that this allows to rewrite the previous circuit in the one below



... but precision is Limited

We assumed  $0 \leq \phi < 1$  takes a value from  $\{0 \cdot \frac{1}{2^n}, \dots, 2^n - 1 \cdot \frac{1}{2^n}\}$   
... an assumption that arose from having only  $n$  qubits to estimate ...

But what to do if  $\phi$  takes none of these values?

Return the  $n$ -bit number  $k$  with  $k \cdot \frac{1}{2^n}$  the value above closest to  $\phi$

Is the circuit above up to this task?

## Setting the stage

Let  $\omega_n = e^{i2\pi \cdot \frac{1}{2^n}}$

and consider the following explicit definition. of  $QFT^{-1}$

$$QFT_n^{-1} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{k=0}^{2^n-1} \omega_n^{-k \cdot x} |k\rangle$$

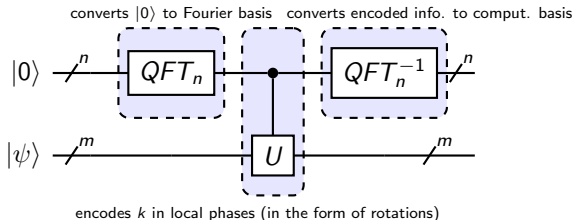
## Setting the stage

Let  $k \cdot \frac{1}{2^n}$  be the value in  $\{0 \cdot \frac{1}{2^n}, \dots, 2^n - 1 \cdot \frac{1}{2^n}\}$  **closest** to  $\phi$ , i.e.

$$\exists_{\epsilon} \cdot 0 \leq |\epsilon| \leq \frac{1}{2^n} \quad \text{and} \quad k \cdot \frac{1}{2^n} + \epsilon = \phi$$

Note that the difference  $\epsilon$  decreases when the number of qubits increases.

Recall the circuit



## Computing the output again

$$|0\rangle$$

$$H^{\otimes n} \mapsto \frac{1}{\sqrt{2^n}} (|0\rangle + |1\rangle + \dots + |2^n - 1\rangle)$$

$$\text{ctrl. } U \mapsto \frac{1}{\sqrt{2^n}} \left( |0\rangle + e^{i2\pi\phi \cdot 1} |1\rangle + \dots + e^{i2\pi\phi \cdot 2^{n-1}} |2^n - 1\rangle \right)$$

$$= \frac{1}{\sqrt{2^n}} \left( |0\rangle + e^{i2\pi(k \cdot \frac{1}{2^n} + \epsilon) \cdot 1} |1\rangle + \dots + e^{i2\pi(k \cdot \frac{1}{2^n} + \epsilon) \cdot 2^{n-1}} |2^n - 1\rangle \right)$$

$$= \frac{1}{\sqrt{2^n}} \sum_{j=0}^{2^n-1} e^{i2\pi(k \cdot \frac{1}{2^n} + \epsilon) \cdot j} |j\rangle$$

$$= \frac{1}{\sqrt{2^n}} \sum_{j=0}^{2^n-1} e^{i2\pi k \cdot \frac{1}{2^n} \cdot j} e^{i2\pi \epsilon \cdot j} |j\rangle$$

$$QFT^{-1} \mapsto \frac{1}{\sqrt{2^n}} \sum_{j=0}^{2^n-1} e^{i2\pi k \cdot \frac{1}{2^n} \cdot j} e^{i2\pi \epsilon \cdot j} \left( \frac{1}{\sqrt{2^n}} \sum_{l=0}^{2^n-1} e^{-i2\pi j \cdot \frac{1}{2^n} \cdot l} |l\rangle \right)$$

$$= \frac{1}{2^n} \sum_{j=0}^{2^n-1} e^{i2\pi k \cdot \frac{1}{2^n} \cdot j} e^{i2\pi \epsilon \cdot j} \left( \sum_{l=0}^{2^n-1} e^{-i2\pi j \cdot \frac{1}{2^n} \cdot l} |l\rangle \right)$$

$$= \frac{1}{2^n} \sum_{j=0}^{2^n-1} \sum_{l=0}^{2^n-1} e^{i2\pi \epsilon \cdot j} e^{i2\pi j \cdot \frac{1}{2^n} \cdot (k-l)} |l\rangle$$

## Looking into the final state

The amplitude of  $|k\rangle$  is

$$\frac{1}{2^n} \sum_{j=0}^{2^n-1} e^{i2\pi\epsilon \cdot j}$$

which is a **finite geometric series**.

Therefore,

$$\frac{1}{2^n} \sum_{j=0}^{2^n-1} e^{i2\pi\epsilon j} = \begin{cases} 1 & \text{if } \epsilon = 0 \\ \frac{1}{2^n} \frac{1-e^{i2\pi\epsilon 2^n}}{1-e^{i2\pi\epsilon}} & \text{if } \epsilon \neq 0 \end{cases}$$

Let us proceed under the assumption  $\epsilon \neq 0$ .

## A geometric detour

$|1 - e^{i\theta}|$  for some angle  $\theta$  is the **Euclidean distance** between 1 and  $e^{i\theta}$  (length of the **straight line segment** between both points)

Consider also **arc length**  $\theta$  between 1 and  $e^{i\theta}$  (distance between the two points by running along the **unit circle**)

### Theorem

Let  $d^E$  and  $d^a$  be respectively the Euclidean distance and arc length between 1 and  $e^{i\theta}$ . Then,

a.  $d^E \leq d^a$

b. if  $0 \leq \theta \leq \pi$  we have  $\frac{d^a}{d^E} \leq \frac{\pi}{2}$

## Finally!

Recall  $\left| \frac{1}{2^n} \frac{1 - e^{i2\pi\epsilon 2^n}}{1 - e^{i2\pi\epsilon}} \right|^2$  is the probability of measuring  $|k\rangle$

$$\begin{aligned} \left| \frac{1}{2^n} \frac{1 - e^{i2\pi\epsilon 2^n}}{1 - e^{i2\pi\epsilon}} \right|^2 &= \left( \frac{1}{2^n} \right)^2 \frac{|1 - e^{i2\pi\epsilon 2^n}|^2}{|1 - e^{i2\pi\epsilon}|^2} \\ &\geq \left( \frac{1}{2^n} \right)^2 \frac{|1 - e^{i2\pi\epsilon 2^n}|^2}{(2\pi\epsilon)^2} && \{\text{Thm a.}\} \\ &\geq \left( \frac{1}{2^n} \right)^2 \frac{\left( \frac{2}{\pi} \cdot 2\pi\epsilon 2^n \right)^2}{(2\pi\epsilon)^2} && \{\text{Thm b.}\} \\ &= \left( \frac{1}{2^n} \right)^2 \frac{(4\epsilon 2^n)^2}{(2\pi\epsilon)^2} \\ &= \left( \frac{1}{2^n} \right)^2 \frac{(2 \cdot 2^n)^2}{\pi^2} = \frac{2^2}{\pi^2} = \frac{4}{\pi^2} \end{aligned}$$



## Working with a superposition of eigenvectors

The algorithm requires an **eigenvector** as input,  
but sometimes is **highly difficult** to build such a vector.

Often it is easier to feed instead a superposition of eigenvectors.

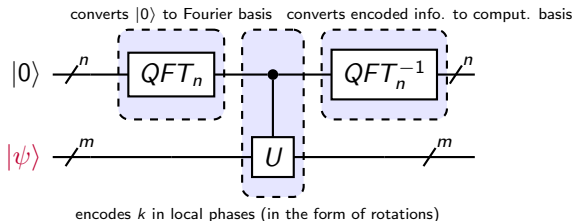
Indeed, by the **spectral theorem** one knows that the eigenvectors  $\{|v_1\rangle, \dots, |v_N\rangle\}$  of  $U$  (with associated eigenvalues  $e^{i2\pi\phi_1}, \dots, e^{i2\pi\phi_N}$ ) form a basis for the  $N(=2^n)$ -dimensional vector space on which  $U$  acts.

Thus, one may define

$$|\psi\rangle = \frac{1}{\sqrt{N}}(|v_1\rangle + \dots + |v_N\rangle)$$

to feed the circuit

## Working with a superposition of eigenvectors



### Exercise

Show that if  $\forall_{i \leq N} \cdot \phi_i \in \{0 \cdot \frac{1}{2^n}, \dots, 2^n - 1 \cdot \frac{1}{2^n}\}$  then the circuit's output is

$$\frac{1}{\sqrt{N}} \left( |x_1\rangle |v_1\rangle + \dots + |x_N\rangle |v_N\rangle \right) \quad \left( \phi_i = x_i \cdot \frac{1}{2^n} \right)$$