## **Quantum Information and Computation**

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## Quantum is trendy ...

Research on quantum technologies is speeding up, and has already created first operational and commercially available applications.

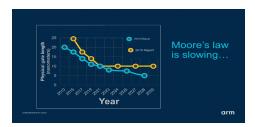
For the first time the viability of quantum computing may be demonstrated in a number of problems and its utility discussed across industries.

Efforts, at national or international levels, to further scale up this research and development are in place.

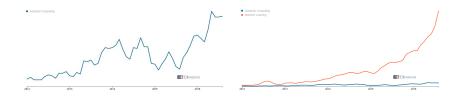
#### ... and full of promises ...

#### Actually,

- Real difficult, complex problems remain out of reach of classical supercomputers
- Classical computer technology is running up against fundamental size limitations (Moore's law),



### ... but the race is just starting



- Clearly, quantum computing will have a substantial impact on societies,
- even if, being a so radically different technology, it is difficult to anticipate its evolution.

#### Two main intelectual achievements of the 20th century met

- Computer Science and Information theory progressed by abstracting from the physical reality. This was the key of its success to an extent that its origin was almost forgotten.
- On the other hand quantum mechanics ubiquitously underlies ICT devices at the implementation level, but had no influence on the computational model itself ...
- until now!

#### Alan Turing (1912 - 1934)



On Computable Numbers, with an Application to the Entscheidungsproblem (1936)

Richard Feynman (1918 - 1988)



Simulating Physics with Computers (1982) (quantum reality as a computational resource)

- C. Bennet and G. Brassard showed how properties of quantum measurements could provide a provably secure mechanism for defining a cryptographic key.
- R. Feynmam recognised that certain quantum phenomena could not be simulated efficiently by a classical computer, and suggested computational simulations may build on quantum phenomena regarded as computational resources.



### Quantum effects as computational resources

#### Superposition

Our perception is that an object — e.g. a bit — exists in a well-defined state, even when we are not looking at it.

However: A quantum state holds information of both possible classical states.

#### Entanglement

Our perception is that objects are directly affected only by nearby objects, i.e. the laws of physics work in a local way.

However: two qubits can be connected, or entangled, st an action performed on one of them can have an immediate effect on the other even at distance.

## Quantum effects as computational resources

#### God plays dice indeed

Our perception is that the laws of Physics are deterministic: there is a unique outcome to every experiment.

However: one can only know the probability of the outcome, for example the probability of a system in a superposition to collapse into a specific state when measured.

#### Uncertainty is a feature, not a bug

Our perception is that with better tools we will be able to measure whatever seems relevant for a problem.

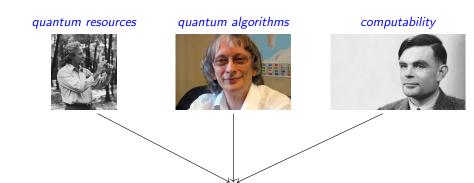
However: there are inherent limitations to the amount of knowledge that one can ascertain about a physical system

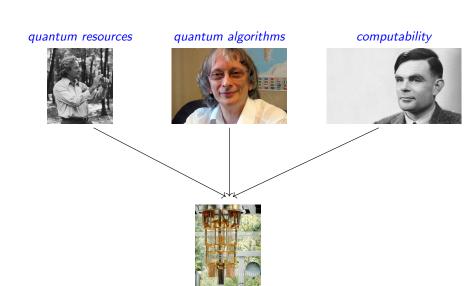
#### Davis Deutsch (1953)

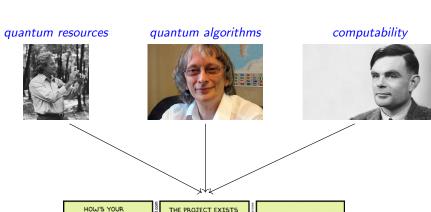


Quantum theory, the Church-Turing principle and the universal quantum computer (1985)

(quantum computability and computational model: first example of a quantum algorithm that is exponentially faster than any possible deterministic classical one)













### Which problems can be addressed?

#### No magic ...

- A huge amount of information can be stored and manipulated in the states of a relatively small number of gubits,
- ... but measurement will pick up just one of the computed solutions and colapse the whole (quantum) state

#### ... but engineering:

To boost the probability of arriving to a solution by canceling out some computational paths and reinforcing others,

depending on the structure of the problem at hands.

## Which problems a Quantum Computer can solve?

- 1994: Peter Shor's factorization algorithm (exponential speed-up),
- 1996: Grover's unstructured search (quadratic speed-up),
- 2018: Advances in hash collision search, i.e finding two items identical in a long list — serious threat to the basic building blocks of secure electronic commerce.
- 2019: Google announced to have achieved quantum supermacy

Availability of proof of concept hardware

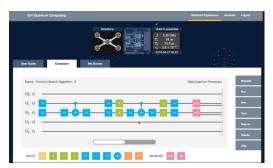
Explosion of emerging applications in several domains: security, finance, optimization, machine learning, ...

## Where exactly do we stand?

#### NISQ - Noisy Intermediate-Scale Quantum Hybrid machines:

- the quantum device as a coprocessor
- typically accessed as a service over the cloud

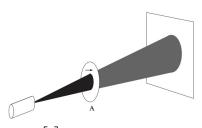




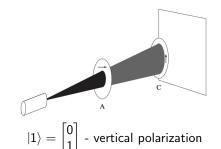
## Still a long way to go ...

- Quantum computations are fragile: noise and decoherence.
- Current methods and tools for quantum software development are still highly fragmentary and fundamentally low-level.
- A lack of reliable approaches to quantum programming will put at risk the expected quantum advantage of the new hardware.

Time to go deeper ...

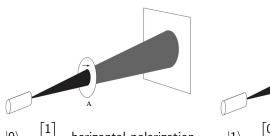


$$|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$
 - horizontal polarization



(from [Reifell & Polak, 2011])

- The probability that a photon passes through the polaroid is the square of the magnitude of the amplitude of its polarization in the direction of the polaroid's preferred axis.
- On passing it becomes polarized in the direction of that axis.



$$|0\rangle = \begin{bmatrix} 1\\0 \end{bmatrix}$$
 - horizontal polarization

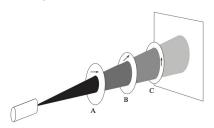
$$|1\rangle = \begin{bmatrix} 0\\1 \end{bmatrix}$$
 - vertical polarization

(from [Reifell & Polak, 2011])

If the photon is polarized as

$$|v\rangle = \alpha |0\rangle + \beta |1\rangle$$

it will go through A with probability  $|\alpha|^2$  and be absorbed with  $|\beta|^2$ .



The polarization of the new polaroid is

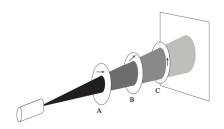
$$|
ightharpoonup 
angle = rac{1}{\sqrt{2}}|1
angle + rac{1}{\sqrt{2}}|0
angle$$

i.e. represented as a superposition of vectors  $|0\rangle$  and  $|1\rangle$ 

Hadamard basis

$$|\rangle\rangle = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$$

$$| \nwarrow \rangle = \frac{1}{\sqrt{2}} | 0 \rangle - \frac{1}{\sqrt{2}} | 1 \rangle$$



Expressing

$$|0\rangle = \frac{1}{\sqrt{2}}|\rangle\rangle + \frac{1}{\sqrt{2}}|\zeta\rangle$$

explains why a visible effect appears when the last polaroid is introduced: the photon goes through C with 50% of probability (i.e.  $|\frac{1}{\sqrt{2}}|^2 = \frac{1}{2}$ ).

Photon's polarization states are represented as unit vectors in a 2-dimensional complex vector space, typically as a

non trivial linear combination  $\equiv$  superposition of vectors in a basis

$$|v\rangle = \alpha |0\rangle + \beta |1\rangle$$

A basis provides an observation (or measurement) tool, e.g.

$$\bigcirc \bigcirc \bigcirc = \{|0\rangle, |1\rangle\} \quad \text{or} \quad \bigcirc \bigcirc \bigcirc = \{|\nearrow\rangle, |\nwarrow\rangle\}$$

Observation of a state

$$|v\rangle = \alpha |u\rangle + \beta |u'\rangle$$

transforms the state into one of the basis vectors in

$$\bigcirc \frown \bigcirc = \{|u\rangle, |u'\rangle\}$$

In other (the quantum mechanics) words: measurement collapses  $|v\rangle$  into a classic, non superimposed state

The probability that observed  $|v\rangle$  collapses into  $|u\rangle$  is the square of the modulus of the amplitude of its component in the direction of  $|u\rangle$ , i.e.

$$|\alpha|^2$$

where, for a complex 
$$\gamma$$
,  $|\gamma| = \sqrt{\overline{\gamma}\gamma}$ 

A subsequent measurement wrt the same basis returns  $|u\rangle$  with probability 1

This observation calls for a restriction to unit vectors, i.e. st

$$|\alpha|^2 + |\beta|^2 = 1$$

to represent quantum states

The notion of superposition is basis-dependent: all states are superpositions with respect to some bases and not with respect to others.

But it is not a probabilistic mixture: it is not true that the state is really either  $|u\rangle$  or  $|u'\rangle$  and we just do not happen to know which.

State  $|u\rangle$  is a definite state, which, when measured in certain bases, gives deterministic results, while in others it gives random results:

The photon with polarization

$$|
ightharpoonup 
angle = rac{1}{\sqrt{2}}|1
angle + rac{1}{\sqrt{2}}|0
angle$$

behaves deterministically when measured with respect to the Hadamard basis but non deterministically with respect to the standard basis

In a sense  $|u\rangle$  can be thought as being simultaneously in both states, but be careful: states that are combinations of basis vectors in similar proportions but with different amplitudes, e.g.

$$\frac{1}{\sqrt{2}}(|u\rangle+|u'\rangle)$$
 and  $\frac{1}{\sqrt{2}}(|u\rangle-|u'\rangle)$ 

are distinct and behave differently in many situations.

Amplitudes are not real (e.g. probabilities) that can only increase when added, but complex so that they can cancel each other or lower their probability, thus capturing another fundamental quantum resource:

interference

#### Qubits

The space of possible polarization states of a photon, as any other quantum system (e.g. photon polarization, electron spin, and the ground state together with an excited state of an atom) that can be modelled by a two-dimensional complex vector space, forms a

quantum bit (qubit)

which has a continuum of possible values.

In practice it is not yet clear which two-state systems will be most suitable for physical realizations of qubits: it is likely that a variety of physical representation will be used.

#### Qubits

#### A qubit has ... a continuum of possible values

- potentially, it can store lots of classical data
- but the amount of information that can be extracted from a qubit by measurement is severely restricted: a single measurement yields at most a single classical bit of information;
- as measurement changes the state, one cannot make two measurements on the original state of a qubit.
- as an unknown quantum state cannot be cloned, it is not possible to measure a qubit's state in two ways, even indirectly by copying its state and measuring the copy.

### The state space of a qubit

Representation redundancy:

qubit state space  $\neq$  complex vector space used for representation

#### Global phase

Unit vectors equivalent up to multiplication by a complex number of modulus one, i.e. a phase  $e^{i\theta}$ , represent the same state.

Let

$$|v\rangle = \alpha |u\rangle + \beta |u'\rangle$$

$$|e^{i\theta}\alpha|^2 = (\overline{e^{i\theta}\alpha})(e^{i\theta}\alpha) = (e^{-i\theta}\overline{\alpha})(e^{i\theta}\alpha) = \overline{\alpha}\alpha = |\alpha|^2$$

and similarly for  $\beta$ .

As the probabilities  $|\alpha|^2$  and  $|\beta|^2$  are the only measurable quantities, the global phase has no physical meaning.

#### The state space of a qubit

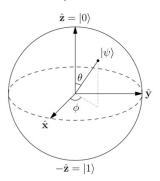
#### Relative phase

Is a measure of the angle between the two complex numbers  $\alpha$  and  $\beta,$  cf

$$\frac{1}{\sqrt{2}}(|u\rangle+|u'\rangle) \quad \frac{1}{\sqrt{2}}(|u\rangle-|u'\rangle) \quad \frac{1}{\sqrt{2}}(e^{i\theta}|u\rangle+|u'\rangle)$$

... cannot be discarded!

There is a bijective correspondence between the state space of a qubit and the complex projective space of dimension 1, which can be represented through the . Bloch sphere:



$$|\psi\rangle=\cos\frac{\theta}{2}|0\rangle+e^{i\varphi}\sin\frac{\theta}{2}|1\rangle$$

where  $0 \leq \theta \leq \pi$  ,  $0 \leq \varphi \leq 2\pi$ 



- The poles represent the classical bits. In general, orthogonal states correspond to antipodal points and every diameter to a basis for the single-qubit state space.
- Once measured a qubit collapses to one of the two poles. Which pole depends exactly on the arrow direction: The angle  $\theta$  measures that probability: If the arrow points at the equator, there is 50-50 chance to collapse to any of the two poles.
- Rotating a vector wrt the z-axis results into a phase change (φ), and does not affect which state the arrow will collapse to, when measured.

Representing 
$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle$$

Express  $|\psi\rangle$  in polar form

$$|\psi\rangle=\rho_1e^{i\varphi_1}|0\rangle+\rho_2e^{i\varphi_2}|1\rangle$$

and eliminate one of the four real parameters multiplying by  $e^{-i\varphi_1}$ 

$$|\psi\rangle=\rho_1|0\rangle+\rho_2e^{i(\varphi_2-\varphi_1)}|1\rangle=\rho_1|0\rangle+\rho_2e^{i\varphi}|1\rangle$$

making  $\phi = \phi_2 - \phi_1$ .

Switch back the coefficient of  $|1\rangle$  to Cartesian coordinates and compute the normalization constraint

$$|\rho_1|^2 + |a+ib|^2 = |\rho_1|^2 + (a-ib)(a+ib) = |\rho_1|^2 + a^2 + b^2 = 1$$

which is the equation of a unit sphere in Real 3-dim space with Cartesian coordinates:  $(a, b, \rho_1)$ .

Back to polar,

$$x = \rho \sin \theta \cos \phi$$
$$y = \rho \sin \theta \sin \phi$$
$$z = \rho \cos \theta$$

So, recalling that  $\rho = 1$ ,

$$\begin{aligned} |\psi\rangle &= z|0\rangle + (a+ib)|1\rangle \\ &= \cos\theta|0\rangle + \sin\theta(\cos\varphi - i\sin\varphi)|1\rangle \\ &= \cos\theta|0\rangle + e^{i\varphi}\sin\theta|1\rangle \end{aligned}$$

which, with two parameters, defines a point in the sphere's surface.

Actually, one may just focus on the upper hemisphere  $(0 \le \theta' \le \frac{\pi}{2})$  as opposite points in the lower one differ only by a phase factor of -1:

Let  $|\psi'\rangle$  be the opposite point on the sphere with polar coordinates  $(1,\pi-\theta',\varphi+\pi)$ 

$$\begin{split} |\psi'\rangle &= \cos{(\pi-\theta')}|0\rangle + e^{i(\phi+\pi)}\sin{(\pi-\theta')}|1\rangle \\ &= -\cos{\theta'}|0\rangle + e^{i\phi}e^{i\pi}\sin{\theta'}|1\rangle \\ &= -\cos{\theta'}|0\rangle + e^{i\phi}\sin{\theta'}|1\rangle \\ &= -|\psi\rangle \end{split}$$

$$|\psi\rangle=\cos\frac{\theta}{2}|0\rangle+e^{i\varphi}\sin\frac{\theta}{2}|1\rangle$$

where  $0 < \theta < \pi$ ,  $0 < \phi < 2\pi$ 

## Complex, inner-product vector space

A set U of vectors generates a complex vector space whose elements can be written as linear combinations of vectors in U:

$$|v\rangle = a_1|u_1\rangle + a_2|u_2\rangle + \cdots + a_n|u_n\rangle$$

i.e.

- Abelian group (V, +, -1, 0)
- with scalar multiplication  $(c \cdot | v)$  distributing over +, often represented by juxtaposition)

• A inner product  $\langle -|-\rangle : V \times V \longrightarrow \mathbb{C}$  such that

$$(1) \quad \langle v | \sum_{i} \lambda_{i} \cdot | w_{i} \rangle \rangle = \sum_{i} \lambda_{i} \langle v | w_{i} \rangle$$

(2) 
$$\langle v|w\rangle = \overline{\langle w|v\rangle}$$

(3) 
$$\langle v|v\rangle \geq 0$$
 (with equality iff  $|v\rangle = 0$ )

Note:  $\langle -|-\rangle$  is conjugate linear in the first argument:

$$\langle \sum_{i} \lambda_{i} \cdot |w_{i}\rangle |v\rangle = \sum_{i} \overline{\lambda_{i}} \langle w_{i}|v\rangle$$

Notation: 
$$\langle v|w\rangle \equiv \langle v,w\rangle \equiv (|v\rangle,|w\rangle)$$

#### Old friends

- $|v\rangle$  and  $|w\rangle$  are orthogonal if  $\langle v|w\rangle=0$
- norm:  $||v\rangle| = \sqrt{\langle v|v\rangle}$
- normalization:  $\frac{|v\rangle}{||v\rangle|}$
- $|v\rangle$  is a unit vector if  $||v\rangle| = 1$
- A set of vectors  $\{|i\rangle,|j\rangle,\cdots,\}$  is orthonormal if each  $|i\rangle$  is a unit vector and

$$\langle i|j\rangle = \delta_{i,j} = \begin{cases} i=j & \Rightarrow 1 \\ \text{otherwise} & \Rightarrow 0 \end{cases}$$

#### Note

A basis for V (set of linearly independent elements of V spanning V) will usually be taken as orthonormal.

 $e^n$ 

The inner product in  $\mathbb{C}^n$  of two vectors over the same orthonormal basis boils down to vector multiplication:

$$\langle v|w\rangle = \langle \sum_{i} v_{i}|i\rangle | \sum_{j} w_{j}|j\rangle \rangle$$

$$= \sum_{i,j} \overline{v_{i}}w_{j}\delta_{i,j}$$

$$= \sum_{i} \overline{v_{i}}w_{i}$$

$$= \left[\overline{v_{1}}\cdots\overline{v_{n}}\right] \begin{bmatrix} w_{1} \\ \vdots \\ w_{n} \end{bmatrix}$$

## Matrices as linear maps

Any  $m \times n$  matrix M can be seen as a linear operator mapping vectors in  $\mathbb{C}^n$  to vectors in  $\mathbb{C}^m$ . Linearity means that

$$M\left(\sum_{j} \alpha_{j} |v_{j}\rangle\right) = \sum_{j} \alpha_{j} M |v_{j}\rangle$$

holds, where the action of M in a m-dimensional vector corresponds to multiplication.

Examples: The Pauli matrices

$$I = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \quad X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad Y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \quad Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$$

## Linear maps as matrices

Let V and W be vector spaces with basis, respectively,

$$B_V = \{|v_1\rangle, \cdots, |v_n\rangle\}$$
 and  $B_W = \{|w_1\rangle, \cdots, |w_m\rangle\}$ 

A linear operator, i.e. a map  $M: V \longrightarrow W$  st

$$M\left(\sum_{j} \alpha_{j} |v_{j}\rangle\right) = \sum_{j} \alpha_{j} M(|v_{j}\rangle)$$

can be represented by a  $m \times n$  matrix st, for each  $j \in 1..n$ ,

$$M(|v_j\rangle) = \sum_i M_{i,j} |w_i\rangle$$

Composition of linear operators amounts to multiplication of the corresponding matrices.

This representation is, of course, basis dependent.



#### Hilbert spaces

Complete, complex, inner-product vector space, complete meaning that any Cauchy sequence

$$|v_1\rangle, |v_2\rangle, \cdots$$

converges

$$\forall_{\epsilon>0} \exists_N \forall_{m,n>0} ||v_m\rangle, |v_n\rangle| \leq \epsilon$$

This completeness condition is trivial in finite dimensional vector spaces

## Classical systems

State spaces in a classical system combine through direct sum:

*n* 2-dimensional vector  $\rightsquigarrow$  a vector in 2*n*-dimensional vector space

#### Direct sum $V \oplus W$

- $B_{V \oplus W} = B_V \cup B_W$  and  $dim(V \oplus W) = dim(V) + dim(W)$
- Vector addition and scalar multiplication are performed in each component and the results added
- $\langle (|u_2\rangle \oplus |z_2\rangle)|(|u_1\rangle \oplus |z_1\rangle)\rangle = \langle u_2|u_1\rangle + \langle z_2|z_1\rangle$
- V and W embed canonically in  $V \oplus W$  and the images are orthogonal under the standard inner product

#### Example

$$\begin{bmatrix} a \\ b \end{bmatrix} \oplus \begin{bmatrix} c \\ d \end{bmatrix} = \begin{bmatrix} a \\ b \\ c \\ d \end{bmatrix}$$

# Quantum systems

State spaces in a classical system combine through tensor:

*n* 2-dimensional vector  $\rightsquigarrow$  a vector in  $2^n$ -dimensional vector space

i.e. the state space of a quantum system grows exponentially with the number of particles: cf, Feyman's original motivation

#### Tensor $V \otimes W$

- $B_{V \otimes W}$  is a set of elements of the form  $|v_i\rangle \otimes |w_j\rangle$ , for each  $|v_i\rangle \in B_V$ ,  $|w_i\rangle \in B_W$  and  $dim(V \otimes W) = dim(V) \times dim(W)$
- $(|u_1\rangle + |u_2\rangle) \otimes |z\rangle = |u_1\rangle \otimes |z\rangle + |u_2\rangle \otimes |z\rangle$
- $|z\rangle \otimes (|u_1\rangle + |u_2\rangle) = |z\rangle \otimes |u_1\rangle + |z\rangle \otimes |u_2\rangle$
- $(\alpha|u\rangle)\otimes|z\rangle = |u\rangle\otimes(\alpha|z\rangle) = \alpha(|u\rangle\otimes|z\rangle)$
- $\langle (|u_2\rangle \otimes |z_2\rangle)|(|u_1\rangle \otimes |z_1\rangle)\rangle = \langle u_2|u_1\rangle \langle z_2|z_1\rangle$

Clearly, every element of  $V \otimes W$  can be written as

$$\alpha_1(|v_1\rangle\otimes|w_1\rangle)+\alpha_2(|v_2\rangle\otimes|w_1\rangle)+\dots+\alpha_{nm}(|v_n\rangle\otimes|w_m\rangle)$$

#### Example

The basis of  $V \otimes W$ , for V, W qubits with the standard basis is

$$\{|0\rangle\otimes|0\rangle,|0\rangle\otimes|1\rangle,|1\rangle\otimes|0\rangle,|1\rangle\otimes|1\rangle\}$$

Thus, the tensor of  $\alpha_1|0\rangle+\beta_1|1\rangle$  and  $\alpha_2|0\rangle+\beta_2|1\rangle$ 

$$\alpha_1\alpha_2|0\rangle\otimes|0\rangle \ + \ \alpha_1\beta_2|0\rangle\otimes|1\rangle \ + \ \alpha_2\beta_1|1\rangle\otimes|0\rangle \ + \ \alpha_2\beta_2|1\rangle\otimes|1\rangle$$

In a simplified notation

$$\alpha_1 \alpha_2 |00\rangle + \alpha_1 \beta_2 |01\rangle + \alpha_2 \beta_1 |10\rangle + \alpha_2 \beta_2 |11\rangle$$

#### Notation

Writing in a more familiar matrix notation requires fixing an ordering for the basis of the tensor product space; typically the lexicographic ordering

#### Example

Let 
$$|u\rangle=\frac{1}{\sqrt{5}}\left[1,-2\right]^T$$
 and  $|z\rangle=\frac{1}{\sqrt{10}}\left[-1,3\right]^T$ . Then 
$$|u\rangle\otimes|z\rangle\ =\ \frac{1}{5\sqrt{2}}\left[-1,3,2,-6\right]^T$$

#### Other basis

... besides the standard one:

#### Bell basis

$$\begin{split} |\Phi^{+}\rangle &= \; \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) \\ |\Phi^{-}\rangle &= \; \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle) \\ |\Psi^{+}\rangle &= \; \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle) \\ |\Psi^{-}\rangle &= \; \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle) \end{split}$$

#### Representation

- As before, vectors that differ only in a global phase represent the same quantum state
- but also the same phase factor in different qubits of a tensor product represent the same state:

$$|u\rangle\otimes(e^{i\varphi}|z\rangle) = e^{i\varphi}(|u\rangle\otimes|z\rangle) = (e^{i\varphi}|u\rangle)\otimes|z\rangle$$

Actually, phase factors in qubits of a single term of a superposition can always be factored out into a coefficient for that term, i.e. phase factors distribute over tensors.

## Representation

• Relative phases still matter (of course!)

$$\frac{1}{\sqrt{2}}(|00\rangle+|11\rangle) \ \ \text{differs from} \ \ \frac{1}{\sqrt{2}}(e^{i\varphi}|00\rangle+|11\rangle)$$

even if

$$\frac{1}{\sqrt{2}}(|00\rangle+|11\rangle) = \frac{1}{\sqrt{2}}(e^{i\phi}|00\rangle+e^{i\phi}|11\rangle) = \frac{e^{i\phi}}{\sqrt{2}}(|00\rangle+|11\rangle$$

- Redundancy: the quantum state space of a n-qubit system has  $2^{n-1}$  complex dimensions
- The complex projective space of dimension 1 (depicted in the Block sphere) generalises to higher dimensions, although in practice linearity makes vector spaces easier to use.

## Entanglement

#### Most states in $V \otimes W$ cannot be written as $|u\rangle \otimes |z\rangle$

- A single-qubit state can be specified by a single complex number so any tensor product of n qubit states can be specified by n complex numbers. But it takes  $2^n 1$  complex numbers to describe states of an n qubit system.
- Since 2<sup>n</sup> >> n, the vast majority of n-qubit states cannot be described in terms of the state of n separate qubits.
- Such states, that cannot be written as the tensor product of *n* single-qubit states, are entangled states.

## Entanglement

#### Example

The Bell state  $|\Phi^+\rangle=\frac{1}{\sqrt{2}}(|00\rangle+|11\rangle)$  is entangled Actually, to make  $|\Phi^+\rangle$  equal to

$$(\alpha_1|0\rangle+\beta_1|1\rangle)\otimes(\alpha_2|0\rangle+\beta_2|1\rangle)\ =\ \alpha_1\alpha_2|00\rangle+\alpha_1\beta_2|01\rangle+\beta_1\alpha_2|10\rangle+\beta_1\beta_2|11\rangle$$

would require that  $\alpha_1\beta_2=\beta_1\alpha_2=0$  which implies that either  $\alpha_1\alpha_2=0$  or  $\beta_1\beta_2=0$ .

#### Note

Entanglement can also be observed in simpler structures, e.g. relations:

$$\{(a,a),(b,b)\}\subseteq A\times A$$

cannot be separated, i.e. written as a Cartesian product of subsets of A.

## Entanglement

The notion of entanglement

- is not basis dependent
- but depends on the tensor decomposition used

Example.

$$u = \frac{1}{2}(|0000\rangle + |0101\rangle + |1010\rangle + |1111\rangle)$$

is entangled wrt the decomposition into single qubits, since it cannot be expressed as the tensor product of four single-qubit states, but it is not for a decomposition consisting of a subsystem of the first and third qubit and another with the second and fourth qubit:

$$u = \frac{1}{\sqrt{2}}(|0_10_3\rangle + |1_11_3\rangle) \otimes \frac{1}{\sqrt{2}}(|0_20_4\rangle + |1_21_4\rangle)$$

# Measuring composed states

## Recalling the single-qubit case

Every measuring tool has an associated orthonormal basis  $\{|v_1\rangle, |v_2\rangle\}$  for the vector space V associated with the single-qubit system.

Each basis vector  $|v_i\rangle$  generates a one-dimensional subspace  $S_i$  consisting of all multiples  $\alpha|v_i\rangle$ , where  $\alpha$  is a complex number, and  $V=S_1\oplus S_2$ , the direct sum decomposition of V.

#### Example

A measuring tool for a qubit in the standard basis has  $V=S_1\oplus S_2$  as the associated direct sum decomposition, where  $S_1$  is generated by  $|0\rangle$  and  $S_2$  by  $|1\rangle$ .

State  $|u\rangle = \alpha |0\rangle + \beta |1\rangle$  will be  $|0\rangle$  with probability  $|\alpha|^2$ , the amplitude of  $|u\rangle$  in the subspace  $S_1$ , and  $|1\rangle$  with probability  $|\beta|^2$ .

# Measuring composed states

## The *n*-qubit case

To every measuring tool corresponds a direct sum decomposition

$$V = S_1 \oplus S_2 \oplus \cdots \oplus S_k$$

of the  $2^n$  dimensional vector space V, for some  $k \le 2^n$  standing for the maximum number of outcomes for a states measured with that toll

# Measuring composed states

#### Example: First qubit of a 2-qubit system with SB

$$V = S_1 \oplus S_2$$

- $S_1 = |0\rangle \otimes V_2$ , the 2-dim subspace spanned by  $\{|00\rangle, |01\rangle\}$
- $S_2=|1\rangle\otimes V_2$ , the 2-dim subspace spanned by  $\{|10\rangle,|11\rangle\}$

#### To measure

$$\begin{split} |u\rangle \; &=\; \alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle \\ \\ |u\rangle \; &=\; \gamma_1|s_1\rangle + \gamma_2|s_2\rangle \\ \\ \gamma_1 = \sqrt{|\alpha_{00}|^2 + |\alpha_{01}|^2} \qquad |s_1\rangle = \frac{1}{\gamma_1}(\alpha_{00}|00\rangle + \alpha_{01}|01\rangle) \\ \\ \gamma_2 = \sqrt{|\alpha_{10}|^2 + |\alpha_{11}|^2} \qquad |s_1\rangle = \frac{1}{\gamma_2}(\alpha_{10}|10\rangle + \alpha_{11}|11\rangle) \end{split}$$

Dirac's bra/ket notation is a handy way to represent elements and constructions on an Hilbert space, amenable to calculations and with direct correspondence to diagrammatic (categorial) representations of process theories

- $|u\rangle$  A ket stands for a vector in an Hilbert space V. In  $\mathbb{C}^n$ , a column vector of complex entries. The identity for + (the zero vector) is just written 0.
- $\langle u|$  A bra is a vector in the dual space  $V^{\dagger}$ , i.e. scalar-valued linear maps in V a row vector in  $\mathbb{C}^n$ .

There is a bijective correspondence between |u
angle and  $\langle u|$ 

$$|u\rangle = \begin{bmatrix} u_1 \\ \vdots \\ u_n \end{bmatrix} \Leftrightarrow [\overline{u}_1 \cdots \overline{u}_n] = \langle u|$$

A tradition going back to Penrose in the 1970's.



Dirac's bra/ket notation provides a convenient way of specifying linear transformations on quantum states:

#### outer product

$$|w\rangle\langle u|(|z\rangle) \widehat{=} |w\rangle\langle u|z\rangle = |w\rangle\langle u|z\rangle = \langle u|z\rangle|w\rangle$$

 matrix multiplication (composition of linear maps) is associative and scalars (zero objects in the corresponding universe) commute with everything

#### Example: $|0\rangle\langle 1|$

$$|0\rangle\langle 1|$$
 maps  $|1\rangle\mapsto|0\rangle$  and  $|0\rangle\mapsto0$ 

$$|0\rangle\langle 1||1\rangle = |0\rangle\langle 1|1\rangle = |0\rangle 1 = |0\rangle |0\rangle\langle 1||0\rangle = |0\rangle\langle 1|0\rangle = |0\rangle 0 = 0$$

Using matrices:

$$|0\rangle\langle 1| \; = \; \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 0 & 1 \end{bmatrix} \; = \; \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix}$$

Example:  $X = |0\rangle\langle 1| + |1\rangle\langle 0|$ 

$$\begin{aligned} |0\rangle\langle 1| + |1\rangle\langle 0| (|0\rangle) &= |0\rangle\langle 1| (|0\rangle) + |1\rangle\langle 0| (|0\rangle) &= 0 + |1\rangle &= |1\rangle \\ |0\rangle\langle 1| + |1\rangle\langle 0| (|1\rangle) &= |0\rangle\langle 1| (|1\rangle) + |1\rangle\langle 0| (|1\rangle) &= |0\rangle + 0 &= |0\rangle \end{aligned}$$

represented by the following matrix in the standard basis:

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

Example:  $|10\rangle\langle00|+|00\rangle\langle10|+|11\rangle\langle11|+|01\rangle\langle01|$  Maps  $|00\rangle\mapsto|11\rangle$  and  $|11\rangle\mapsto|00\rangle$  Clearly,

$$\begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

An operator on an n-qubit system that maps the basis vector  $|j\rangle$  to  $|i\rangle$  and all other standard basis elements to 0 can be expressed in the standard basis as

$$O = |i\rangle\langle j|$$

Matrix for O has a single non-zero entry 1 in the i, j place.

A general operator A with entries  $a_{ij}$  in the standard basis can be written

$$A = \sum_{i} \sum_{j} a_{ij} |i\rangle\langle j|$$

Conversely, the i, j entry of the matrix for A in the standard basis is given by

$$\langle i|A|j\rangle$$

#### Example

Let  $|s\rangle = \sum_{k} \beta_{k} |k\rangle$ .

$$A|s\rangle = \left(\sum_{i} \sum_{j} a_{ij} |i\rangle\langle j|\right) \left(\sum_{k} \beta_{k} |k\rangle\right)$$
$$= \sum_{i} \sum_{j} \sum_{k} a_{ij} \beta_{k} |i\rangle\langle j| |k\rangle$$
$$= \sum_{i} \sum_{j} a_{ij} \beta_{j} |i\rangle$$

In general, given a basis  $B_V = \{|\beta_i\rangle\}$  for a N-dimensional Hilbert space V, an operator

$$A:V\longrightarrow V$$

can be written as

$$\sum_{i} \sum_{j} b_{ij} |\beta_{i}\rangle\langle\beta_{j}|$$

wrt this basis. The matrix entries are  $b_{ij}$ , as expected.

The Dirac's notation is

- independent of the basis and the order of the basis elements
- more compact
- and builds up intuitions ...

## **Projectors**

$$V = S \oplus S^{\dagger}$$

Any vector  $|v\rangle$  can be written uniquely as the sum of a vector  $\vec{s_1}$  from  $S_1$  and  $\vec{s_2}$  from  $S_2$  (not unit vectors in the general case)

## Projector

$$P_S: V \longrightarrow S$$
 st  $|v\rangle = \vec{s_1} + \vec{s_2} \mapsto \vec{s_1}$ 

**Example**  $|u\rangle\langle u|$  is the projector onto the subspace spanned by  $|u\rangle$ .

A measuring tool with associated decomposition

$$V = \bigoplus_{i} S_{i}$$

into ortogonal subspaces  $S_i$ , acting over a state  $|v\rangle$  produces, with probability  $|P_i|v\rangle|^2$ , a state

$$\frac{P_i|v\rangle}{|P_i|v\rangle|}$$



## **Projectors**

**Example** Let  $|v\rangle = \alpha |0\rangle + \beta |1\rangle$ . Projector  $|0\rangle\langle 0|$  obtains its component in the subspace generated by  $|0\rangle$ , i.e.

$$|0\rangle\langle 0|(|v\rangle) = \alpha|0\rangle\langle 0||0\rangle + \beta|0\rangle\langle 0||1\rangle = \alpha|0\rangle$$

Similarly, projector  $|10\rangle\langle10|$  acts on a two-qubit state

$$v = \alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle$$

yielding

$$|10\rangle\langle10|(|v\rangle) = \alpha_{10}|10\rangle$$

and

$$|00\rangle\langle00| + |10\rangle\langle10|(|v\rangle) = \alpha_{00}|00\rangle + \alpha_{10}|10\rangle$$



# Projectors are self-adjoint

## Adjoint operator

Operator  $O^{\dagger}: U \longrightarrow V$  is adjoint to  $O: V \longrightarrow U$  if, for any vectors from V and U, the inner product between  $O^{\dagger}(\vec{u})$  and  $\vec{v}$  coincides with the inner product between  $\vec{u}$  and  $O(\vec{v})$ . In Dirac's notation,

$$(\langle u|O)|v\rangle = \langle u|(O|v\rangle) = \langle u|O|v\rangle$$

recalling that  $(O|v\rangle)^{\dagger} = \langle v|O^{\dagger}$ .

Clearly, the matrix representation of  ${\it O}^{\dagger}$  is the conjugate transpose of that of  ${\it O}$ 

Clearly, PP = P (why?), which combined with  $P^{\dagger} = P$ , yields

$$|P|v\rangle|^2 = (\langle v|P^{\dagger})(P|v\rangle) = \langle v|P|v\rangle$$

## **Projectors**

#### Example

Let  $|v\rangle = \alpha |0\rangle + \beta |1\rangle$ .

Applying projector  $P_0 = |0\rangle\langle 0|$  to  $|v\rangle$  results in the state

$$\frac{P_0|v\rangle}{|P_0|v\rangle|^2} = \frac{\alpha|0\rangle}{|\alpha|} \sim |0\rangle$$

where

$$P_0|v\rangle = (|0\rangle\langle 0|)|v\rangle = |0\rangle\langle 0|v\rangle = \alpha|0\rangle$$

with probability

$$|P_0|v\rangle|^2 = \langle v|P_0|v\rangle = \langle v||0\rangle\langle 0||v\rangle = \langle v|0\rangle\langle 0|v\rangle = \overline{\alpha}\alpha = |\alpha|^2$$

## **Projectors**

## Example: measuring up to (bit equality)

$$V = S_e \oplus S_n$$

with  $S_e$  the subspace generated by  $\{|00\rangle, |11\rangle\}$  in which the two bits are equal, and  $S_n$  its complement.

When measuring

$$v = \alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle$$

with this device, yields a state in which the two bit values are equal with probability

$$\langle v|P_{\rm e}|v\rangle = (\sqrt{|\alpha_{00}|^2 + |\alpha_{11}|^2})^2 = |\alpha_{00}|^2 + |\alpha_{11}|^2$$

Of course, the measurement does not determine the value of the two bits, only whether the two bits are equal

## Can the explicit decomposition be avoided?

#### Hermitian operators

- define a unique orthogonal subspace decomposition, their eigenspace decomposition, and
- for every such decomposition, there exists a corresponding Hermitian operator whose eigenspace decomposition coincides with it

## Hermitian operators

 $O: V \longrightarrow V$  is Hermitian if

$$O^{\dagger} = O$$

The relevant property is that, for every eigenvalue  $\lambda$  with eigenvector  $|I\rangle$ ,  $\lambda = \overline{\lambda}$ , and thus all eigenvalues of a Hermitian operator are real, because

$$\lambda \langle I | I \rangle = \langle I | \lambda | I \rangle = \langle I | (O|I) \rangle = (\langle I | O^{\dagger}) | I \rangle = (O|I)^{\dagger} | I \rangle = (\lambda |I|)^{\dagger} |I| \rangle = \overline{\lambda} \langle I | I \rangle$$

## Orthogonality

For any O, two distinct eigenvalues have disjoint eigenspaces, because, for any unit vector  $|v\rangle$ ,

$$O|v\rangle = \lambda |v\rangle$$
 and  $O|v\rangle = \lambda' |v\rangle$  and  $(\lambda - \lambda')|v\rangle = 0$ 

and thus  $\lambda = \lambda'$ .

For any Hermitian O, the eigenvectors for distinct eigenvalues must be orthogonal, because

$$\lambda \langle v | w \rangle = (\langle v | O^{\dagger}) | w \rangle = \langle v | (O | w \rangle) = \mu \langle v | w \rangle$$

for any pairs  $(\lambda, |\nu\rangle), (\mu, |w\rangle)$  with  $\lambda \neq \mu$ .

Thus,  $\langle v|w\rangle=0$ , because  $\lambda\neq\mu$ , and the corresponding subspaces are orthogonal.

## Eigenspace decomposition of *V* for *O*

Any Hermitian O determines a unique decomposition for V

$$V = \bigoplus_{\lambda_i} S_{\lambda_i}$$

and any decomposition  $V=\oplus_{i=1}^k S_i$  can be realized as the eigenspace decomposition of a Hermitian operator

$$O = \sum_{i} \lambda_{i} P_{i}$$

where each  $P_i$  is the projector onto  $S_i$  and  $L = \{\lambda_1, \dots, \lambda_k\}$  is a set of arbitrary, real k values

Thus, in a measurement, a subspace decomposition can be specified by a Hermitian operator

Note that the values in L are irrelevant — they are just labels for the corresponding subspaces, i.e. labels for the measurement outcomes.

### The measurement postulate

- Any measurement is specified by a Hermitian operator O
- The possible outcomes of measuring a state  $|v\rangle$  with O are labeled by the eigenvalues of O
- The probability of obtaining the outcome labelled by  $\lambda_i$  is

$$|P_i|v\rangle|^2$$

The state after measurement is the normalized projection

$$\frac{P_i|v\rangle}{|P_i|v\rangle}$$

onto the  $\lambda_i$ -eigenspace  $S_i$ . Thus, the state after measurement is a unit length eigenvector of O with eigenvalue  $\lambda_i$ 

#### Notes

- A measurement is not modelled by the action of a Hermitian operator on a state, but of the corresponding projectors.
- Actually, Hermitian operators are only a bookeeping trick
- A Hermitian operator uniquely specifies a subspace decomposition
- For a given subspace decomposition there are many Hermitian operators whose eigenspace decomposition is that decomposition.

### Example: Measuring a single qubit in the Hadamard basis

• Projectors:

$$P_{+} = |+\rangle\langle+| = \frac{1}{2}(|0\rangle\langle0| + |0\rangle\langle1| + |1\rangle\langle0| + |1\rangle\langle1|)$$

$$P_{-} = |-\rangle\langle-| = \frac{1}{2}(|0\rangle\langle0| - |0\rangle\langle1| - |1\rangle\langle0| + |1\rangle\langle1|)$$

• Hermitian:

$$X = |0\rangle\langle 1| + |1\rangle\langle 0| = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

for an arbitrary choice of  $\lambda_+=1$  and  $\lambda_-=-1$ 

### Example: Measuring of the first qubit in the standard basis

$$EB = |00\rangle\langle00| + |01\rangle\langle01| + \pi|10\rangle\langle10| + |11\rangle\langle11| = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & \pi & 0 \\ 0 & 0 & 0 & \pi \end{bmatrix}$$

specifies measurement of a two-qubit system with respect to the decomposition

$$V = \{|00\rangle, |01\rangle\} \oplus \{|10\rangle, |11\rangle\}$$

**Exercise**: What is the Hermitian for measuring bit equality?

- $O_1 \otimes O_2$  is an Hermitian operator over space  $V_1 \otimes V_2$  if each  $O_i$  is such over  $V_i$ .
- Its eigenvalues are the product of eigenvalues of the original operators, in multiple ways.
- However, most Hermitian operators O on  $V_1 \otimes V_2$  cannot be written as a tensor product of two Hermitian operators acting separately in each space.

### but only if

each subspace in the subspace decomposition described by O can be written as  $S = S_1 \otimes S_2$ , for  $S_i$  the subspace decomposition associated to  $O_i$ 

### Example

$$Z \otimes Z = |00\rangle\langle 00| - |01\rangle\langle 01| - |10\rangle\langle 10| + |11\rangle\langle 11|$$

specifies the measurement for bit equality.

Not all measurements are tensor products of single-qubit measurements **Example** 

O determines whether both bits are set to one. The result of a measurement with O is a state in the subspace spanned by

$$\{|11\rangle\}$$
 or by  $\{|00\rangle, |01\rangle, |10\rangle\}$ 

Measuring with O is quite different from measuring both qubits in the standard basis and composing the results: e.g. state

$$|v\rangle = \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle)$$

is unchanged when measured by O.

**Exercise**: but what results from measuring both qubits?

### Measurement

A Hermitian operator of the form

$$1 \otimes \cdots \otimes O \otimes \cdots \otimes I$$

on a *n*-qubit system forms a single-qubit measurement of that system

Measurement operators in the standard basis, when combined with transformations, are sufficient to perform arbitrary quantum measurements.

In particular, all possible subspace decompositions of the state space can be obtained by starting with a subspace decomposition in which all of the subspaces are generated by standard basis vectors and transforming (because there are quantum operations taking any basis to any other)

**Exercise**: How many classical bits does a single measurement of an *n*-qubit system reveal?

# Closed systems

... transformations that map the state space of the quantum system to itself

**Exercise**: Is measurement one of these transformations?

- All quantum transformations on n-qubit quantum systems can be expressed as a sequence of transformations on 1-qubit and 2-qubit subsystems.
- Efficiency of a quantum transform (quantified in terms of the number of 1- or 2-qubit gates used) will not be addressed here.

# Unitary transformations

All transformations are linear:

$$U(\alpha_1|v_1\rangle + \cdots + \alpha_k|v_k\rangle) = \alpha_1 U|v_1\rangle + \cdots + \alpha_2 U|v_k\rangle$$

 Unit length vectors map to unit length vectors, thus orthogonal subspaces map to orthogonal subspaces.

These properties hold iff *U* preserves inner product:

$$\langle v|U^{\dagger}U|w\rangle = \langle v|w\rangle$$

which entails

$$U^{\dagger}U = I$$
 U is unitary

# Unitary transformations

- Unitary operators map orthonormal bases to orthonormal bases, since they preserve the inner product
- Moreover, any linear transformation that maps an orthonormal basis to an orthonormal basis is unitary
- If given in matrix form, being unitary means that the set of columns of its matrix representation are orthonormal (because the *i*th column is the image of  $U|i\rangle$ ).
- equivalently, rows are orthonormal (why?)

Unitary transformations are reversible

# Unitary transformations

New transformations from old Both  $U_1U_1$  and  $U_1 \otimes U_2$  are unitary.

But linear combinations of unitary operators, however, are not in general unitary.

## The no-cloning theorem

### Linearity implies that quantum states cannot be cloned

Let  $U(|a\rangle|0\rangle)=|a\rangle|a\rangle$  and consider state  $|c\rangle=\frac{1}{\sqrt{2}}(|a\rangle+|b\rangle)$  for  $|a\rangle$  and  $|b\rangle$  orthogonal. Then

$$U(|c\rangle|0\rangle) = \frac{1}{\sqrt{2}}(U(|a\rangle|0\rangle) + U(|b\rangle|0\rangle))$$

$$= \frac{1}{\sqrt{2}}(|a\rangle|a\rangle + |b\rangle|b\rangle)$$

$$\neq \frac{1}{\sqrt{2}}(|a\rangle|a\rangle + |a\rangle|b\rangle + |b\rangle|a\rangle + |b\rangle|b\rangle)$$

$$= |c\rangle|c\rangle$$

$$= U(|c\rangle|0\rangle)$$

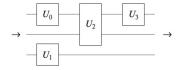
This result, however, does not preclude the construction of a known quantum state from a known quantum state.



## Quantum gates

A gate is a transformation that acts on only a small number of qubits Differently from the classical case, they do not necessarily correspond to physical objects

#### Notation



### Is there a complete set?

In general no: there are uncountably many quantum transformations, and a finite set of generators can only generate countably many elements.

However, it is possible for finite sets of gates to generate arbitrarily close approximations to all unitary transformations.



### Quantum gates

### Pauli gates

$$\begin{split} I &= |0\rangle\langle 0| + |1\rangle\langle 1| = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \quad X &= |1\rangle\langle 0| + |0\rangle\langle 1| = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \\ Z &= |0\rangle\langle 0| - |1\rangle\langle 1| = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \quad Y &= ZX &= -|1\rangle\langle 0| + |0\rangle\langle 1| = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix} \end{split}$$

### Hadamard gate

$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1\\ 1 & -1 \end{bmatrix}$$

$$H|0\rangle = |+\rangle = \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle)$$

$$H|1\rangle = |-\rangle = \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$



### The CNOT gate

Acts on the standard basis for a 2-qubit system, flipping the second bit if the first bit is 1 and leaving it unchanged otherwise.

$$\begin{array}{lll} \textit{CNOT} &=& |0\rangle\langle 0| \otimes \textit{I} \; + \; |1\rangle\langle 1| \otimes \textit{X} \\ &=& |0\rangle\langle 0| \otimes (|0\rangle\langle 0| + |1\rangle\langle 1|) \; + \; |1\rangle\langle 1| \otimes (|1\rangle\langle 0| + |0\rangle\langle 1|) \\ &=& |00\rangle\langle 00| + |01\rangle\langle 01| + |11\rangle\langle 10| + |10\rangle\langle 11| \\ &=& \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \\ \end{array}$$

CNOT is unitary and is its own inverse, and cannot be decomposed into a tensor product of two 1-qubit transformations

### The CNOT gate

The importance of CNOT is its ability to change the entanglement between two qubits, e.g.

$$\begin{array}{ll} \textit{CNOT} \, \left( \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \otimes |0\rangle \right) \; = \; \textit{CNOT} \, \left( \frac{1}{\sqrt{2}} (|00\rangle + |10\rangle) \right) \\ & = \; \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle) \end{array}$$

Since it is its own inverse, it can take an entangled state to an unentangled one.

Note that entanglement is not a local property in the sense that transformations that act separately on two or more subsystems cannot affect the entanglement between those subsystems:

$$(U \otimes V) |v\rangle$$
 is entangled iff  $|v\rangle$  is

### Generalising the CNOT gate



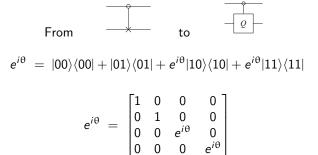
$$C_Q = |0\rangle\langle 0| \otimes I + |1\rangle\langle 1| \otimes Q$$

In the standard basis

$$C_Q = \begin{bmatrix} 1 & 0 \\ 0 & Q \end{bmatrix}$$

# Controlled phase shift gate

Changes the phase of the second bit iff the control bit is 1:

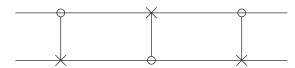


Transforming a global into a local phase

$$rac{1}{\sqrt{2}}(|00
angle+|11
angle \longrightarrow rac{1}{\sqrt{2}}(|00
angle+e^{i heta}|11
angle$$

### Exercise

#### Discuss

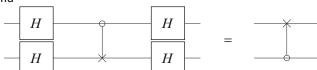


### Notes

- A unitary transformation on the complex vector space is completely determined by its action on a basis, but not by specifying what states the states corresponding to basis states are sent to. Example:  $e^{i\theta}$  takes the four quantum states to themselves (because  $|10\rangle$  and  $e^{i\theta}|10\rangle$  represent the same state, but a global phase can be transformed into a local one, as above).
- The notions of control/target bit depends on the basis.
   Example: Apply CNOT in the Hadamard basis to get

$$|++\rangle \mapsto |++\rangle \ |+-\rangle \mapsto |--\rangle \ |-+\rangle \mapsto |-+\rangle \ |--\rangle \mapsto |+-\rangle$$

and



## Dense coding

Aim: encode and transmit two classical bits with one qubit and a shared EPR pair.

This result is surprising, since only one bit can be extracted from a qubit

The idea is that, since entangled states can be distributed ahead of time, only one qubit needs to be physically transmitted to communicate two bits of information.

Let Alice (Bob) be sent and operate the first (second) qubit of pair

$$|r\rangle = \frac{1}{\sqrt{2}} (|0\rangle|0\rangle + |1\rangle|1\rangle)$$

### EPR pairs

... are entangled states

named after Einstein, Podolsky, and Rosen, from the *hidden-variable* controversy

## Dense coding

#### Alice

wishes to transmit the state of two classical bits encoding one of the numbers 0 through 3. Depending on this number, Alice performs one of the Pauli transformations on her qubit of the entangled pair  $|r\rangle$ , and sends her qubit to Bob.

	Transformation	New state
0	$ r\rangle = (I \times I) r\rangle$	$\frac{1}{\sqrt{2}}( 00\rangle +  11\rangle$
1	$ r_1\rangle = (X \times I) r\rangle$	$\frac{\sqrt{1}}{\sqrt{2}}( 10\rangle +  01\rangle$
2	$ r_3\rangle = (Z \times I) r\rangle$	$\frac{\sqrt{1}}{\sqrt{2}}( 00\rangle -  11\rangle$
3	$ r_3\rangle = (Y \times I) r\rangle$	$\begin{array}{c} \frac{1}{\sqrt{2}}( 00\rangle +  11\rangle \\ \frac{1}{\sqrt{2}}( 10\rangle +  01\rangle \\ \frac{1}{\sqrt{2}}( 00\rangle -  11\rangle \\ \frac{1}{\sqrt{2}}(- 10\rangle +  01\rangle \end{array}$

### Dense coding

### Bob

to decode the information, applies a *CNOT* to the two qubits of the entangled pair and then H to the first qubit:

$$CNOT \longrightarrow \begin{bmatrix} \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) \\ \frac{1}{\sqrt{2}}(|11\rangle + |01\rangle) \\ \frac{1}{\sqrt{2}}(|00\rangle - |10\rangle) \\ \frac{1}{\sqrt{2}}(-|11\rangle + |01\rangle) \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \otimes |0\rangle \\ \frac{1}{\sqrt{2}}(|1\rangle + |0\rangle) \otimes |1\rangle \\ \frac{1}{\sqrt{2}}(-|1\rangle + |0\rangle) \otimes |1\rangle \end{bmatrix}$$

$$H \otimes I \longrightarrow \begin{bmatrix} |00\rangle \\ |01\rangle \\ |10\rangle \\ |11\rangle \end{bmatrix}$$

Bob then measures the two qubits in the standard basis to obtain the 2-bit binary encoding of the number Alice wished to send

Aim: to transmit, using two classical bits, the state of a single qubit.

#### Surprisingly,

- shows that two classical bits suffice to communicate a qubit state (which has an infinite number of configurations)
- provides a mechanism for the transmission of an unknown quantum state (in spite of the no-cloning theorem)

Note that the original state cannot be preserved (precisely because of the no-cloning result), which motivates the name of the protocol ...

#### Alice

... has a qubit whose state  $|v\rangle=\alpha|0\rangle+\beta|1\rangle$  she does not know, but wants to send to Bob through classical channels.

The starting point is the 3-qubit state whose first 2 qubits are controlled by Alice and the last by Bob:

$$|v\rangle \otimes |r\rangle = \frac{1}{\sqrt{2}} (\alpha|0\rangle \otimes (|00\rangle + |11\rangle) + \beta|1\rangle \otimes (|00\rangle + |11\rangle))$$
$$= \frac{1}{\sqrt{2}} (\alpha|000\rangle + \alpha|011\rangle + \beta|100\rangle + \beta|111\rangle)$$

#### Alice

... then she applies  $CNOT \otimes I$  and  $H \otimes I \otimes I$  to obtain

$$\begin{split} (H \otimes I \otimes I)(CNOT \otimes I)(|\nu\rangle \otimes |r\rangle) \\ &= (H \otimes I \otimes I) \frac{1}{\sqrt{2}} (\alpha|000\rangle + \alpha|011\rangle + \beta|110\rangle + \beta|101\rangle) \\ &= \frac{1}{2} (\alpha(|000\rangle + |011\rangle + |100\rangle + |111\rangle) + \beta(|010\rangle + |001\rangle - |110\rangle - |101\rangle)) \\ &= \frac{1}{2} (|00\rangle (\alpha|0\rangle + \beta|1\rangle) + |01\rangle (\alpha|1\rangle + \beta|0\rangle) + \\ &+ |10\rangle (\alpha|0\rangle - \beta|1\rangle) + |11\rangle (\alpha|1\rangle - \beta|0\rangle)) \end{split}$$

#### Alice

Alice measures the first two qubits and obtains one of the four standard basis states,  $|00\rangle, |01\rangle, |10\rangle, |11\rangle$ , with equal probability.

Depending on the result of her measurement, the state of Bob's qubit is projected to

$$\alpha|0\rangle + \beta|1\rangle, \ \alpha|1\rangle + \beta|0\rangle, \ \alpha|0\rangle - \beta|1\rangle, \ \alpha|1\rangle - \beta|0\rangle$$

Then, Alice sends the result of her measurement as two classical bits to Bob.

After these transformations, crucial information about the original state  $|v\rangle$  is contained in Bob's qubit, Alice's being destroyed ...

#### Bob

When Bob receives the two bits from Alice, he knows how the state of his half of the entangled pair compares to the original state of Alice's qubit.

Bob can reconstruct the original state of Alice's qubit,  $|v\rangle$ , by applying the appropriate decoding transformation to his qubit, originally part of the entangled pair.

Bits received	Bob's state	Transformation to decode
00	$\alpha 0\rangle + \beta 1\rangle$	1
01	lpha 1 angle+eta 0 angle	X
10	$\alpha 0\rangle - \beta 1\rangle$	Z
10	$\alpha 1\rangle - \beta 1\rangle$	Y

After decoding, Bob's qubit will be in the state Alice's qubit started.

Teleportation and dense coding are in some sense inverse protocols (why?)



# A probabilistic machine

States: Given a set of possible configurations, states are vectors of probabilities in  $\mathbb{R}^n$  which express indeterminacy about the exact physical configuration, e.g.  $[p_0 \cdots p_n]^T$  st  $\sum_i p_i = 1$ 

Operator: double stochastic matrix (must come (go) from (to) somewhere), where  $M_{i,j}$  specifies the probability of evolution from configuration i to i

Evolution: computed through matrix multiplication with a vector  $|u\rangle$  of current probabilities

- $M|u\rangle$  (next state)
- $|u\rangle^T M^T$  (previous state)

Measurement: the system is always in some configuration — if found in i, the new state will be a vector  $|t\rangle$  st  $t_i = \delta_{i,i}$ 

# A probabilistic machine

### Composition:

$$p\otimes q \ = \ egin{bmatrix} p_1 \ 1-p_1 \end{bmatrix} \otimes egin{bmatrix} q_1 \ 1-q_1 \end{bmatrix} \ = \ egin{bmatrix} p_1 q_1 \ p_1 (1-q_1) \ (1-p_1) q_1 \ (1-p_1) (1-q_1) \end{bmatrix}$$

• correlated states: cannot be expressed as  $p \otimes q$ , e.g.

Operators are also composed by ⊗ (Kronecker product):

$$M \otimes N = \begin{bmatrix} M_{1,1}N & \cdots & M_{1,n}N \\ \vdots & & \vdots \\ M_{m,1}N & \cdots & M_{m,n}N \end{bmatrix}$$

## A quantum machine

States: given a set of possible configurations, states are unit vectors of (complex) amplitudes in  $\mathbb{C}^n$ 

Operator: unitary matrix  $(M^{\dagger}M = I)$ . The norm squared of a unitary matrix forms a double stochastic one.

Evolution: computed through matrix multiplication with a vector  $|u\rangle$  of current amplitudes (wave function)

- $M|u\rangle$  (next state)
- $|u\rangle^T M^T$  (previous state)

Measurement: configuration i is observed with probability  $|\alpha_i|^2$  if found in i, the new state will be a vector  $|t\rangle$  st  $t_j = \delta_{j,i}$  Composition: also by a tensor on the complex vector space; may exist

entangled states

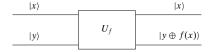
# A quantum machine

### Structure of a quantum algorithm

- 1. State preparation (fix initial setting): typically the qubits in the initial classical state are put into a superposition of many states;
- 2. Transform, through unitary operators applied to the superposed state;
- Measure, i.e. projection onto a basis vector associated with a measurement tool.

Is  $f: \mathbf{2} \longrightarrow \mathbf{2}$  constant, with a unique evaluation?

### Oracle



where  $\oplus$  stands for exclusive disjunction.

- The oracle takes input  $|x,y\rangle$  to  $|x,y\oplus f(x)\rangle$
- for y = 0 the output is  $|x, f(x)\rangle$

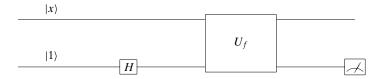
Is  $f: \mathbf{2} \longrightarrow \mathbf{2}$  constant, with a unique evaluation?

### Oracle

• The oracle is a unitary, i.e. reversible gate

$$|x,(y \oplus f(x)) \oplus f(x)\rangle = |x,y \oplus (f(x) \oplus f(x))\rangle = |x,y \oplus 0\rangle = |x,y\rangle$$

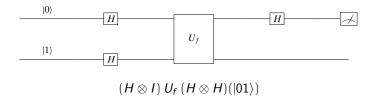
Idea: Avoid double evaluation by superposition



The circuit computes:

output 
$$=|x\rangle \frac{|0 \oplus f(x)\rangle - |1 \oplus f(x)\rangle}{\sqrt{2}}$$
  
 $=\begin{cases} |x\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}} & \Leftarrow f(x) = 0\\ |x\rangle \frac{|1\rangle - |2\rangle}{\sqrt{2}} & \Leftarrow f(x) = 1 \end{cases}$   
 $=(-1)^{f(x)}|x\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}}$ 

### Idea: Avoid double evaluation by superposition



### Input in superposition

$$|\sigma_1\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}} \frac{|0\rangle - |1\rangle}{\sqrt{2}} = \frac{|00\rangle - |01\rangle + |10\rangle - |11\rangle}{2}$$

$$\begin{array}{ll} |\sigma_2\rangle \;=\; \left(\frac{(-1)^{f(0)}|0\rangle+(-1)^{f(1)}|1\rangle}{\sqrt{2}}\right) \, \left(\frac{|0\rangle+|1\rangle}{\sqrt{2}}\right) \\ &=\; \begin{cases} (\underline{+}1) \left(\frac{|0\rangle+|1\rangle}{\sqrt{2}}\right) \, \left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right) & \Leftarrow f \text{ constant} \\ (\underline{+}1) \left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right) \, \left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right) & \Leftarrow f \text{ not constant} \end{cases} \end{array}$$

$$\begin{array}{ll} |\sigma_3\rangle \ = \ H|\sigma_2\rangle \\ \\ = \ \begin{cases} (\underline{+}1)\,|0\rangle\,\left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right) & \Leftarrow f \text{ constant} \\ (\underline{+}1)\,|1\rangle\,\left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right) & \Leftarrow f \text{ not constant} \end{cases}$$

To answer the original problem is now enough to measure the first qubit: if it is in state  $|0\rangle$ , then f is constant.